

The Coq Proof Assistant

Semantics and applications to verification

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April 1st. 2022

What is a proof assistant ?

A tool to **formalize** and **verify** proofs

The key word is assistant: it *assists* the user in

- defining the proof goals formally;
- setting up the structure of the proofs;
- making the proof steps;
- checking the overall consistency of the proof, at the end.

Some steps are more assisted than others:

- formalization is done with respect to the knowledge of the user, it is **error prone**
- key structural arguments (induction hypotheses and such) are very hard to get right in general
- checking a series of proof steps is easier to mechanize...

Purpose of Coq and principle

Coq is a programming language

- We can **define data-types** and **write programs** in Coq
- Language similar to a **pure functional language**
- **Very expressive** type system (more on this later)

- Programs can be ran inside Coq
- Programming language of the year ACM Award in 2014...

Coq is a proof assistant

- It allows to **express theorems** and **proofs**
- It can **verify** a proof
- It can also **infer some proofs** or **proof steps**

- Proof search is usually mostly manual and takes most of the time

Main proof assistants

Coq: the topic of this lecture

Isabelle / HOL: a higher order logic framework

- syntax is closer to the logics
- proof term underneath...

ACL2: A Computational Logic for Applicative Common Lisp

- a framework for automated reasoning
- based on functional common lisp

PVS: Prototype Verification System

- kernel extends Church types
- less emphasis on the notion of proof term, more emphasis on automation

Overall workflow

- 1 **Define the objects** properties need be proved about
Data-structures, base types, programs written in the Coq (or vernacular) language
- 2 Write and prove **intermediate lemmas**
 - ▶ a theorem is defined by a formula in the Coq language.
 - ▶ a proof requires a sequence of **tactics applications**
tactics are described as part of a separate language.
 - ▶ at the end of the proof, a **proof term** is constructed and verified.
- 3 Write and prove the **main theorems**
- 4 If needed, **extract** programs

Two languages: one for **definitions/theorems/proofs**, one for **tactics**

In Coq, everything is a term

- The **core of Coq** is defined by a language of **terms**
- **Commands** are used in order to manipulate terms

Examples of terms:

- **base values:** 0, 1, true...
- **types:** nat, bool, but also Prop...
- **functions:** `fun (n: nat) => n + 1`
- **function applications:** `(fun (n: nat) => n + 1) 8`
- **logical formulas:**
`exists p: nat, 8 = 2 * p,`
`forall a b: Prop, a /\ b -> a`
- **complex functions** (more on this one later):
`fun (a b : Prop) (H : a /\ b) =>`
`and_ind (fun (H0 : a) (_ : b) => H0) H`

In Coq, every term has a type

- As observed, **types are terms**
- Every term also **has a type**, denoted by `term: type`

- `0: nat`
- `nat: Set`
- `Set: Type`
- `Type: Type` (*caveat: not quite the same instance*)
- `(fun (n: nat) => n + 1): nat -> nat`
- more complex types get interesting:

```
fun (a b : Prop) (H : a /\ b) =>
  and_ind (fun (H0 : a) (_ : b) => H0) H
: forall a b: Prop, a /\ b -> a
```

Curry-Howard correspondence

The core principle of Coq

- A proof of P can be viewed a **term of type P**
- A proof of $P \implies Q$ can be viewed a **function** transforming a proof of P into a proof of Q , hence, a **function of type $P \rightarrow Q$** ...

Similarity between **typing** rules and **proof** rules:

$$\frac{\frac{\Gamma, x : P \vdash u : Q}{\Gamma \vdash \lambda x. u : P \rightarrow Q} \text{ fun}}{\Gamma \vdash u : P \rightarrow Q \quad \Gamma \vdash v : P} \text{ app}$$

$$\frac{\frac{\Gamma, P \vdash Q}{\Gamma \vdash P \implies Q} \text{ implic}}{\Gamma \vdash P \implies Q \quad \Gamma \vdash P} \text{ mp}$$

Correspondance:

program	proof
type	theorem

Searching a proof of P
 \equiv **searching u of type P**

Defining a term

Two ways:

- 1 **Define it fully**, with **its type** and **its definition**

```
Definition zero: nat := 0.
```

```
Definition incr (n: nat): nat := n + 1.
```

- 2 Provide **only its type** and **search for a proof of it**

```
Lemma lzero: nat.
```

```
  exact 0.
```

```
Save.
```

```
Definition lincr: forall n: nat, nat.
```

```
  intro. exact (n + 1).
```

```
Save.
```

- **Definition:** Definition name `u: t := def.`
- **Proof:** Definition name `u: t.` or Lemma name `u: t.`

Inductive definition

- A **very powerful** mechanism
- In Coq, **almost everything** is actually an inductive definition
... examples: **integers, booleans, equality, conjunction...**

- **Syntax:**

```
Inductive tree : Set :=  
  | leaf: tree  
  | node: tree -> tree -> tree.
```

- **Induction principles** automatically provided by Coq, and to use in induction proofs:

```
tree_ind: forall P : tree -> Prop,  
  P leaf  
-> (forall t : tree, P t -> forall t0 : tree, P t0  
    -> P (node t t0))  
-> forall t : tree, P t
```

Recursive functions

- Very natural to work with inductive definitions
- **Caveat: must provably terminate**
this is usually checked with a **strict sub-term condition**

- **Syntax:**

```
Fixpoint size (t: tree) : nat :=  
  match t with  
  | leaf => 0  
  | node t0 t1 => 1 + (size t0) + (size t1)  
end.
```

- **Ill formed definition, rejected by the system (termination issue):**

```
Fixpoint f (t: tree): nat :=  
  match t with  
  | leaf | node leaf leaf => 0  
  | node _ _ => f (node leaf leaf)  
end.
```

Proving a term

View in proof mode:

```
a : Prop
b : Prop
H : a /\ b
H0 : a
H1 : b
=====
a
```

- above the bar: **current assumptions**
- below the bar: **current subgoal** (there may be several goals)
- **at the end:** displays No more subgoals.
- command `Save.` stores the term.

Progression towards a finished proof:

- based on commands called **tactics**
- in the background, Coq **constructs the proof term**

A few tactics, and their effect

- Each tactic performs a basic operation on the current goal
 - In the background, Coq **constructs the proof term**
 - At the end, the term is **independantly checked** (very reliable !)
- **Introduction of an assumption** (proof tree and term):

$$\frac{\Gamma, P \vdash Q}{\Gamma \vdash P \implies Q}$$

$$\frac{\Gamma, x : P \vdash u : Q}{\Gamma \vdash \lambda x. u : P \longrightarrow Q}$$

- **Application of an implication:**

$$\frac{\Gamma \vdash P \implies Q \quad \Gamma \vdash P}{\Gamma \vdash Q}$$

$$\frac{\Gamma \vdash u : P \longrightarrow Q \quad \Gamma \vdash v : P}{\Gamma \vdash u v : Q}$$

- **Immediate conclusion of a subgoal:**

$$\overline{\Gamma, P \vdash P}$$

$$\overline{\Gamma, x : P \vdash x : P}$$

Automation in Coq

So far, we have considered fairly manual tactics...

There are also **automated tactics**, that typically call an external program to try to solve a goal, and then constructs a proof term:

- either verify the proof term afterwards...
- ... or call a function proved once and for all to build it

Examples:

- **Tauto**: decides propositional logic
- **Omega**: solves a class of numeric (in)-equalities (see manual)

Language of tactics:

more advanced users can combine tactics to build their own

Proof by reflection: prove decision procedures, and invoke them...

A glimpse at the tactic language

Most common tactics:

Tactic	Effect
<code>intro.</code>	Introduce one assumption
<code>intros.</code>	Introduce as many assumptions as possible
<code>apply H.</code>	Applies assumption H (should be of the form $A \rightarrow B$)
<code>elim H.</code>	Decomposes assumption H
<code>exact t.</code>	Provides a proof term for current sub-goal
<code>trivial.</code>	Conclude immediately very simple proofs.
<code>induction t.</code>	Perform induction proof over term t
<code>rewrite H.</code>	Rewrite assumption H (should be of the form $t_0 = t_1$)
<code>tauto.</code>	Decision procedure in propositional logic

Do not hesitate to look at the online manual !

A glimpse at the command language

Most common tactics (should be enough for a TD):

Command	Meaning
Check <code>t</code> .	Prints the type of term <code>t</code>
Print <code>t</code> .	Prints the type and definition of term <code>t</code>
Definition <code>u</code> : <code>t := [term]</code> .	Full definition of term <code>u</code>
Lemma <code>t</code> . Theorem <code>t</code> . Definition <code>t</code> .	Start a proof of term <code>t</code>
Save.	Exit proof mode and save proof term