An Introduction to Lustre

Marc Pouzet

École normale supérieure
Marc.Pouzet@ens.fr

April 30, 2020
The language Lustre

- Invented by Paul Caspi and Nicolas Halbwachs around 1984, in Grenoble (France).
- A program is a set of equations. An equation defines an infinite sequence of values.
- Boolean operators applied point-wise, a unit-delay, and sampling operators.
- Equivalent graphical representation by block diagrams.
- Feedback loops must cross a unit delay.
- Time is synchronous: at every tick of a global clock, every operation does a step.
- Code generation to sequential code and formal verification techniques.
- An industrial success: SCADE (Esterel-Technologies company) is used for programming critical control software (e.g., planes, nuclear plants).
Lustre

Program by writing stream equations.

|   | 1 | 2 | 1 | 4 | 5 | 6 | ...
|---|---|---|---|---|---|---|---|
| X | 1 | 2 | 1 | 4 | 5 | 6 | ...
| Y | 2 | 4 | 2 | 1 | 1 | 2 | ...
| 1 | 1 | 1 | 1 | 1 | 1 | 1 | ...
| X + Y | 3 | 6 | 3 | 5 | 6 | 8 | ...
| X + 1 | 2 | 3 | 2 | 5 | 6 | 7 | ...

Equation $Z = X + Y$ means that at any instant $n \in \mathbb{N}$, $Z_n = X_n + Y_n$.

Time is logical: inputs $X$ and $Y$ arrive at the same time; the output $Z$ is produced at the same time.

Synchrony means that at instant $n$, all streams take their $n$-th value.

In practice, check that the current output is produced before the next input arrives.
Example: 1-bit adder

node full_add(a, b, c:bool) returns (s, co:bool);
let
   s = (a xor b) xor c;
   co = (a and b) or (b and c) or (a and c);
   tel;

or:

node full_add(a, b, c:bool) returns (s, co:bool);
let
   co = if a then b or c else b and c;
   s = (a xor b) xor c;
   tel;
Full Adder

Compose two “half adder”

```plaintext
node half_add(a, b: bool)
returns (s, co: bool);
    let s = a xor b;
    co = a and b;
    tel;
```

Instanciate it twice:

```plaintext
node full_add_h(a, b, c: bool)
returns (s, co: bool);
    var s1, c1, c2: bool;
    let
        (s1, c1) = half_add(a, b);
        (s, c2) = half_add(c, s1);
        co = c1 or c2;
    tel;
```
Verify properties

How to be sure that `full_add` and `full_add_h` are equivalent?

\[ \forall a, b, c : \text{bool}. \, \text{full}_\text{add}(a, b, c) = \text{full}_\text{add}_\text{h}(a, b, c) \]

Write the following program and prove that it returns true at every instant!

```plaintext
-- file prog.lus
node equivalence(a,b,c:bool) returns (ok:bool);
    var o1, c1, o2, c2: bool;
    let
        (o1, c1) = full_add(a,b,c);
        (o2, c2) = full_add_h(a,b,c);
        ok = (o1 = o2) and (c1 = c2);
    tel;

Then, use the model-checking tool lesar:

% lesar prog.lus equivalence
--Pollux Version 2.2

TRUE PROPERTY
```
The Unit Delay

One can refer to the value of an input at the “previous” step.

<table>
<thead>
<tr>
<th>$X$</th>
<th>1</th>
<th>2</th>
<th>1</th>
<th>4</th>
<th>5</th>
<th>6</th>
<th>...</th>
</tr>
</thead>
<tbody>
<tr>
<td>$pre\ X$</td>
<td>nil</td>
<td>1</td>
<td>2</td>
<td>1</td>
<td>4</td>
<td>5</td>
<td>...</td>
</tr>
<tr>
<td>$Y$</td>
<td>2</td>
<td>4</td>
<td>2</td>
<td>1</td>
<td>1</td>
<td>2</td>
<td>...</td>
</tr>
<tr>
<td>$Y -&gt; pre\ X$</td>
<td>2</td>
<td>1</td>
<td>2</td>
<td>1</td>
<td>4</td>
<td>5</td>
<td>...</td>
</tr>
<tr>
<td>$S$</td>
<td>1</td>
<td>3</td>
<td>4</td>
<td>8</td>
<td>13</td>
<td>19</td>
<td>...</td>
</tr>
</tbody>
</table>

The stream $(S_n)_{n \in \mathbb{N}}$ with $S_0 = X_0$ and $S_n = S_{n-1} + X_n$, for $n > 0$ is written:

$$S = X \rightarrow \text{pre}\ S + X$$

Introducing intermediate equations does not change the meaning of programs:

$$S = X \rightarrow I; I = \text{pre}\ S + X$$
Example: convolution

Define the sequence:

\[
Y_0 = X_0 / 2 \quad \land \quad \forall n > 0. Y_n = (X_n + X_{n-1}) / 2
\]

node convolution(X:real) returns (Y:real);
let Y = (X + (0 -> pre X)) / 2.0;
tel;

or:

node convolution(X:real) returns (Y:real);
var pY:int;
let Y = (X + pY) / 2;
   pY = 0 -> pre X;
tel;
Linear filters

FIR (Finite Impulse Response)

\[ y(n) = \sum_{m=0}^{L-1} x(n - m)b(m) \]

IIR (Infinite Impulse Response) or recursive filter

\[ y(n) = \sum_{m=0}^{L-1} x(n - m)b(m) + \sum_{m=1}^{M-1} y(n - m)a(m) \]
Build a block-diagram with three operators: a gain (multiplication by a constant), a sum and a unit delay (register).

Previous example

\[ \forall n \geq 0. y(n) = \frac{1}{2} (x(n) + x(n - 1)) \]
Example: follow $x$ with a 20% gain.

$$\forall n \geq 0. y(n) = 0.2(x(n) - y(n - 1)) + y(n - 1)$$

node filter(x: real) returns (y:real);
let y = 0.0 -> 0.2 * (x - pre y) + pre y; tel;

Retiming:
Optimise by moving unit delays around combinatorial operators.

**DEMO**: type luciole filter.lus filter
Counting events

Count the number of instants where the input signal tick is true between two top.

node counter(tick, top:bool) returns (cpt:int);
let
    cpt = if top then 0
         else if tick then (0 -> pre cpt) + 1 else pre cpt;
tel;

Is this program well defined? Is it deterministic? No: initialization issue.

<table>
<thead>
<tr>
<th>tick</th>
<th>f</th>
<th>f</th>
<th>t</th>
<th>t</th>
<th>t</th>
<th>t</th>
<th>t</th>
<th>t</th>
<th>t</th>
<th>...</th>
</tr>
</thead>
<tbody>
<tr>
<td>top</td>
<td>f</td>
<td>f</td>
<td>f</td>
<td>f</td>
<td>f</td>
<td>t</td>
<td>f</td>
<td>f</td>
<td>f</td>
<td>...</td>
</tr>
<tr>
<td>cpt</td>
<td>nil</td>
<td>nil</td>
<td>nil</td>
<td>nil</td>
<td>nil</td>
<td>0</td>
<td>1</td>
<td>2</td>
<td>...</td>
<td></td>
</tr>
</tbody>
</table>
An explicit Euler integrator

node integrator(const step: real; x0, x’:real) returns (x:real);
let
    x = x0 -> pre(x) + x’ * step;
tel;

step is a constant stream computed at compile-time.

Sinus/cosine functions

node sinus_cosinus(theta:real)
returns (sin,cos:real);
let sin = theta * integrator(0.01, 0.0, cos);
    cos = -. theta * integrator(0.01, 1.0, pre sin);
tel;
Initial Value Problem (IVP)

$f$ is a combinatorial function with $y$ of type $ty$. $t$ is the current time. $x(t)$ be defined by the IVP:

\[
\dot{x} = f(y, t, x) \quad \text{with} \quad x(0) = x_0
\]

```plaintext
node ivp(const step: real; y: ty; read) returns (x: real)
    var t: real;
    let
        x = integr(step, x0, f(y, t, x));
        t = 0.0 -> pre t + step;
    tel;

Exercice

- Program a classical explicit Runge Kutta method (e.g., order 4).
- More difficult: program a variable step Runge Kutta method (RK45). Hint: use a control bit `error_too_large` to shrink the step dynamically.
```
Counting Beacons

Counting beacons and seconds to decide whether a train is on time.

Use an **hysteresis** with a low and high threshold to reduce oscillations.

```plaintext
node beacon(sec, bea: bool) returns (ontime, late, early: bool);
var diff, pdiff: int; pontime: bool;
let
    pdiff = 0 -> pre diff;
    diff = pdiff + (if bea then 1 else 0) +
        (if sec then -1 else 0);
    early = pontime and (diff > 3) or
        (false -> pre early) and (diff > 1);
    late = pontime and (diff < -3) or
        (false -> pre late) and (diff < -1);
    ontime = not (early or late);
    pontime = true -> pre ontime;
end;
```

---

1 This example is due to Pascal Raymond
Two types of properties

Safety property
“Something wrong never happen”, i.e., a property is invariant and true in any accessible state. E.g.:
- “The train is never both early and late”, it is either on time, late or early;
- “The train never passes immediately from late to early”; “It is impossible to stay late only a single instant”.

Liveness property
“Something good with eventually happen.”, i.e., any execution will reach a state verifying the property. E.g., “If the trains stop, it will eventually be late.”

Remark:
“If the train is on time and stops for ten seconds, it will be eventually late” is a safety property. Safety properties are critical ones in practice.
Formal verification and modeling of systems

A safety property (“something bad will never happen”) is a boolean proved to be true at every instant.

Example: the alternating bit protocol
A transmitter $A$; a receiver $B$. Two unreliable lines $A2B$ and $B2A$ that may lose messages.

- $A$ asks for one input. It re-emits the data with $bit = true$ until it receives $ack = true$.
- It asks for another input and emits the data with $bit = false$ until it receives $ack = false$.
- $B$ sends $ack = false$ until it receives $bit = true$; it sends $ack = true$ until it receives $bit = false$.
- initialization: send anything with $bit = true$. The first message arriving with $bit = false$ is valid.
Objective:

Model and prove the protocol is correct, i.e., the network is the identity function (input sequence = output sequence) with two unreliable lines.

Model the asynchronous communication by adding a “presence” bit to every data: a pair \((data, enable)\) is meaningful when \(enable = true\).
The Sender

- A asks for one input. It re-emits the data with \( bit = true \) until it receives \( ack = true \).
- It asks for another input and emits the data with \( bit = false \) until it receives \( ack = false \).

```
node A(dataIn: int; recB: bool; ack: bool)
returns (reqData: bool; send: bool; data: int; bit: bool);

var
    buff: int; chstate : bool;

let
    buff = if reqData then dataIn else (0 -> pre buff);
    chstate = recB and (bit = ack);
    reqData, send, bit =
        (false, true, true) ->
        pre (if chstate then (true, true, not bit)
            else (false, send, bit));

    data = buff;

tel
```
The Receiver

- $B$ sends $ack = false$ until it receives $bit = true$; it sends $ack = true$ until it receives $bit = false$;

```plaintext
node B(recA : bool; data: int; bit: bool;)
returns (sendOut: bool; dataOut: int; send2A: bool; ack: bool);

var chstate : bool;

let
  chstate = recA and (ack xor bit);

  sendOut, send2A, ack = 
  (false, true, true) ->
    pre (if chstate then (true, true, not ack)
         else (false, true, ack));
  dataOut = data;

tel
```
Modeling the channel and the main property

node unreliable(loose: bool; presIn: bool) returns (presOut: bool);
let
  presOut = presIn and not loose;
tel

-- The property that two signals [r] and [s] alternate.
node altern(r,s: bool) returns (ok: bool);
var
  s0, s1 : bool;
  ps0, ps1 : bool;
let
  ps0 = true -> pre s0;
  ps1 = false -> pre s1;
  s0 = ps0 and (r = s) or ps1 and s and not r;
  s1 = ps0 and r and not s or ps1 and not r and not s;
  ok = (true -> pre ok) and (s0 or s1);
tel
The main system

node obs(dataIn: int; looseA2B, looseB2A : bool;) returns (ok : bool; reqData: bool; sendOut: bool);

var
dataOut: int;
sendA2B: bool; data: int; bit: bool;
recA2B, recB2A : bool;
sendB2A: bool; ack: bool;

let

ok = altern(reqData, sendOut);

recA2B = unreliable(looseA2B, sendA2B);
recB2A = unreliable(looseB2A, sendB2A);

reqData, sendA2B, data, bit = A(dataIn, recB2A, ack);
sendOut, dataOut, sendB2A, ack = B(recA2B, data, bit);
tel

%aneto.local: lesar ba.lus obs
TRUE PROPERTY
Clocks: mixing slow and fast processes

A slow process is made by sampling its inputs; a fast one by oversampling its inputs.

The operators when, current and merge

<table>
<thead>
<tr>
<th>B</th>
<th>(false)</th>
<th>(true)</th>
<th>(false)</th>
<th>(true)</th>
<th>(false)</th>
<th>(false)</th>
</tr>
</thead>
<tbody>
<tr>
<td>X</td>
<td>(x_0)</td>
<td>(x_1)</td>
<td>(x_2)</td>
<td>(x_3)</td>
<td>(x_4)</td>
<td>(x_5)</td>
</tr>
<tr>
<td>Y</td>
<td>(y_0)</td>
<td>(y_1)</td>
<td>(y_2)</td>
<td>(y_3)</td>
<td>(y_4)</td>
<td>(y_5)</td>
</tr>
<tr>
<td>(Z = X \text{ when } B)</td>
<td></td>
<td>(x_1)</td>
<td></td>
<td>(x_3)</td>
<td></td>
<td></td>
</tr>
<tr>
<td>(K = Y \text{ when } \neg B)</td>
<td>(y_0)</td>
<td></td>
<td>(y_2)</td>
<td>(y_4)</td>
<td>(y_5)</td>
<td></td>
</tr>
<tr>
<td>(T = \text{ current } Z)</td>
<td>(nil)</td>
<td>(x_1)</td>
<td>(x_1)</td>
<td>(x_3)</td>
<td>(x_3)</td>
<td>(x_3)</td>
</tr>
<tr>
<td>(O = \text{ merge } B ; Z ; K)</td>
<td>(y_0)</td>
<td>(x_1)</td>
<td>(y_2)</td>
<td>(x_3)</td>
<td>(y_4)</td>
<td>(y_5)</td>
</tr>
</tbody>
</table>

The operator merge is not part of Lustre. It was introduced later in Lucid Synchrone and SCADE 6.
The Gilbreath trick

The Gilbreath shuffle (from Wikipedia):

- Deal off any number of the cards from the top of a deck onto a new pile.
- Riffle the new pile with the remainder of the deck.

A trick based on the resulting Gilbreath permutations was formalized and verified in Coq by G. Huet.²

Presentation of the magic trick in G. Huet paper:

Why is this a card trick? Our boolean words are card decks, with true for red and false for black. Take an even deck x, arranged alternatively red, black, red, black, etc. Ask a spectator to cut the deck, into sub-decks u and v. Now shuffle u and v into a new deck w. When shuffling, note carefully whether u and v start with opposite colors or not. If they do, the resulting deck is composed of pairs red-black or black-red; otherwise, you get the property by first rotating the deck by one card. The trick is usually played by putting the deck behind your back after the shuffle, to perform “magic”. The magic is either rotating or doing nothing. When showing the pairing property, say loudly “red black red black...” in order to confuse in the spectator’s mind the weak paired property with the strong alternate one.

There is a variant. If the cut is favorable, that is if u and v are opposite, just go ahead showing the pairing, without the “magic part.” If the spectator says that he understands the trick, show him the counter-example in the non-favorable case. Of course now you have to leave him puzzled, and refuse to redo the trick.

Input: two decks of alternating colours (red, black, red, black, ...) whose bottom cards have different colours.

Output: one deck of alternating red/black pairs.

The property is implied by the following one on Boolean streams:

\textit{if } s_1 \textit{ and } s_2 \textit{ be two alternating streams starting with different values; let } o \textit{ be a stream built by “riffle shuffling” } s_1 \textit{ and } s_2, \textit{ then } o \textit{ is such that it is the succession of pairs of different values.}
The Gilbreath trick in Scade 6 [2]

node Gilbreath_stream (clock c:bool)
returns (prop: bool; o:bool);
var
  s1 : bool when c;
  s2 : bool when not c;
  half : bool;
let
  s1 = (false when c) -> not (pre s1);
  s2 = (true  when not c) -> not (pre s2);
  o = merge (c; s1; s2);
  half = false -> (not pre half);
prop = true -> not (half and (o = pre o));
tel;
The Gilbreath trick in Lustre

node Gilbreath_stream (c:bool) returns (OK: bool; o:bool);
var ps1, s1 : bool;
    ps2, s2 : bool;
    half : bool;
let
    s1 = if c then not ps1 else ps1;
    ps1 = false -> pre s1;
    s2 = if not c then not ps2 else ps2;
    ps2 = true  -> pre s2;

    o = if c then s1 else s2;

    half = false -> not (pre half);

    OK = true  -> not (half and (o = pre o));
tel;

Proved automatically using Lesar or Kind 2.
A classical use of clock: the activation condition
Run a process on a slower by sub-sampling its inputs; hold outputs.

\[\text{node sum}(i: \text{int}) \text{ returns } (s: \text{int});\]

\[
\begin{align*}
\text{let} \\
\quad s &= i \rightarrow \text{pre } s + i; \\
\text{tel};
\end{align*}
\]

<table>
<thead>
<tr>
<th>cond</th>
<th>sum(1)</th>
<th>sum(1 when cond)</th>
<th>(sum 1) when cond</th>
</tr>
</thead>
<tbody>
<tr>
<td>true</td>
<td>1</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>false</td>
<td>2</td>
<td>2</td>
<td>3</td>
</tr>
<tr>
<td>true</td>
<td>3</td>
<td>3</td>
<td>4</td>
</tr>
<tr>
<td>false</td>
<td>4</td>
<td>4</td>
<td>5</td>
</tr>
<tr>
<td>true</td>
<td>5</td>
<td>6</td>
<td>6</td>
</tr>
</tbody>
</table>

Sampling inputs vs sampling outputs

- current \((f(x \text{ when } c))\) is called an “activation condition”
- \(f(x \text{ when } c) \neq (f x) \text{ when } c\)
- current\((x \text{ when } c) \neq x\)
Why synchrony?

let half = true -> not (pre half);
  o = x & (x when half);
tel

It defines the sequence: $\forall n \in \mathbb{N}. o_n = x_n \& x_{2n}$

- It cannot be computed in bounded memory.
- Its corresponding Kahn networks has unbounded buffering.
- This is forbidden, a dedicated analysis for that: clock calculus
Clocks must be declared and visible from the interface of a node.

node stables(i:int) ← base clock (true)
returns (s:int; ncond:bool;
   (ns:int) when ncond); ← clock declaration
var cond:bool;
   (l:int) when cond; ← clock declaration
let
   cond = true -> i <> pre i;
   ncond = not cond;
   l = somme(i when cond);
   s = current(l);
   ns = somme(i when ncond);
tel;
Constraints

Rules

- Constants are on the base clock of the node.
- By default, variables are on the base clock of the node.
- Unless a clock is associated to the variable definition.
- $\text{clock}(e_1 \ op \ e_2) = \text{clock}(e_1) = \text{clock}(e_2)$
- $\text{clock}(e \ when \ c) = c$
- $\text{clock}(\text{current}; e) = \text{clock}(\text{clock}(e))$

Implementation choices

- Clocks are declared and verified. No automatic inference.
- Two clocks are equal if expressions that define them are syntactically equal.
One hot coding of Mealy machines

Represent a state by a Boolean variable.

node switch(set,reset:bool) returns (ok :bool);
  var on: bool;
  let
    on = false ->
      if set and not (pre on) then true
      else if reset and (pre on) then false
      else (pre on);
    ok = on;
  tel;

Think in term of an invariant: what is the expression defining the current value of on at every instant?
Verification with assertions

Consider a second version.

```plaintext
node switch2(set, reset:bool) returns (ok:bool);
  var s1, s2: bool;
let
  s1 = true -> if reset and pre s2 then true
               else if pre s1 and set then false else pre s1;
  s2 = false -> if set and pre s1 then true
               else if pre s2 and reset then false else pre s2;
  ok = s2;
let;
```

```plaintext
node compare(set, reset: bool) returns (ok: bool);
  let ok = switch(set, reset) = switch2(set, reset); tel;
```

We get:

```
% lesar prog.lus compare
--Pollux Version 2.2

TRUE PROPERTY
```
Synchronous observers

Comparison is a particular case of a synchronous observer.

• Let \( y = F(x) \), and \( ok = P(x, y) \) for the property relating \( x \) and \( y \)
• \( \text{assert}(H(x, y)) \) is an hypothesis on the environment.

\[
\text{node check}(x:t) \text{ returns } (ok:bool);
\]
\[
\text{let}
\]
\[
\text{assert } H(x,y); \]
\[
y = F(x); \]
\[
ok = P(x,y); \]
\[
\text{tel};
\]

If \( \text{assert} \) is (infinitely) true, then \( \text{ok} \) stay infinitely true
\( (\text{always}(\text{assert})) \Rightarrow (\text{always}(\text{ok})) \).

Any safety temporal property can be expressed as a Lustre program [7, 6]. No temporal logic/language is necessary.

\textbf{Safety temporal properties are regular Lustre programs}
Array and slices

Array are manipulated by slices with implicit point-wise extension of operations. \( t[0..N] \) defines a slice of \( t \) from index 0 to \( N \).

\[
\text{const } N = 10;
\]

\[
\text{node plus(const N: int; a1, a2: int}^N\text{) returns (o: int}^N\text{);}\\
\text{let}\\
\quad o[1..N] = a1[1..N] + a2[1..N];\\
\text{tel;}
\]
-- serial adder

def node add(a1: bool^N; a2: bool^N; carry: bool)
returns (a: bool^N; new_carry: bool);
    var c: bool^N;
    let
        (a[0..N-1], c[0..N-1]) =
            bit_add(a1[0..N-1], a2[0..N-1], ([carry] | c[0..N-2]));
        new_carry = c[N-1];
    tel;

def node add_short(a1: bool^N; a2: bool^N; carry: bool)
returns (a: bool^N; new_carry: bool);
    var c: bool^N;
    let
        (a, c) = bit_add(a1, a2, ([carry] | c[0..N-2]));
        new_carry = c[N-1];
    tel;
Modeling asynchronous communication

Example: “quasi-synchrony”.

node is_qs (const d: int; x, y: bool) returns (ok : bool);
var xav, yav, pxav, pyav : int;
let
  pxav = 0 -> pre xav;
  pyav = 0 -> pre yav;
  xav = if y then 0 else if x then pxav + 1 else pxav;
  yav = if x then 0 else if y then pyav + 1 else pyav;
  ok = (xav <= d) and (yav <= d);
  
tel
A purely boolean implementation

const sz=3;
type count = bool^sz;

const zero = [true, false, false];

node succ(x: count) returns (y: count);
let
    y = ([false] | x[0..1]);
tel

node inf2(x: count) returns (y:bool);
let
    y = x[0] or x[1] or x[2];
tel
A purely boolean implementation

node is_qs2 ( x, y: bool) returns (ok : bool);
var xav, yav, pxav, pyav : count;
let
    pxav = zero -> pre xav;
    pyav = zero -> pre yav;

    xav = if y^3 then zero else if x^3 then succ(pxav) else pxav;
    yav = if x^3 then zero else if y^3 then succ(pyav) else pyav;

    ok = (inf2(xav) and inf2(yav)) and (true -> pre ok);
Conclusion

Compilation

- Static, compile-time checking to ensure the absence of deadlock, that the code behave deterministically.
- Execution in bounded memory and time.
- Code generation into sequential “single loop” code. More advanced methods into automata and/or modular.

Verification by Model-checking

- Synchronous observer: a safety property is a Lustre program
- Avoid to introduce an ad-hoc temporal logic.
- Tool Lesar (BDD technique) by Pascal Raymond (VERIMAG, France).
- KIND and KIND2 (k-induction, PDR based on SMT techniques) by Cesare Tinelli (Iowa State Univ., USA).
- Plug-in (k-induction based on SAT techniques) by Prover-Technologies (associated to SCADE 6).
Related languages and verification tools

Various teams have done their own variant of Lustre.

Language embedding in Haskell

- Copilot (Nasa, USA), an Embedding of Lustre.
- FRAN (images, animation), Functional Reactive Programming (FRP), Hawk (architecture), Lava (synchronous circuits).
- Based on a compilation-by-evaluation technique.

Language extensions, formal verification

- Heptagon: Extended Lustre (automata, arrays) with controller synthesis (Gwenael Delaval, Univ. Grenoble)
- Prelude: Lustre with periodic clocks and a compiler that generates tasks for a real-time OS (Julien Forget, Univ. Lille).
- Lustre compiler at Onera for verification purposes (Pierre-Loic Garoche)
References

Darek Biernacki, Jean-Louis Colaco, Grégoire Hamon, and Marc Pouzet.
Clock-directed Modular Code Generation of Synchronous Data-flow Languages.

Jean-Louis Colaco, Bruno Pagano, and Marc Pouzet.

Jr Edmund M Clarke, Orna Grumberg, and Doron A Peled.
*Model Checking*.

N. Halbwachs.
*Synchronous programming of reactive systems.*

N. Halbwachs, P. Caspi, P. Raymond, and D. Pilaud.
The synchronous dataflow programming language LUSTRE.

N. Halbwachs, F. Lagnier, and C. Ratel.
Programming and verifying real-time systems by means of the synchronous data-flow programming language lustre.

N. Halbwachs, F. Lagnier, and P. Raymond.
Synchronous observers and the verification of reactive systems.

N. Halbwachs, P. Raymond, and C. Ratel.
Generating efficient code from data-flow programs.