

# Overview of Cryptography

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Reference: <http://cacr.uwaterloo.ca/hac/about/chap1.pdf>

# Vernam Ciphers a.k.a. One-Time Pad

- Let  $A=\{0,1\}$ . A binary message  $m_1m_2\dots m_t$  is operated by a binary key  $k_1k_2\dots k_t$  of the same length to produce the ciphertext string  $c_1c_2\dots c_t$ :
- $c_i = m_i \oplus k_i$  for all  $i=1\dots t$ .
- To decrypt, it suffices to compute  $c_i \oplus k_i$  to recover  $m_i$ .
  
- If we encrypt two messages with the same key, we can obtain  $c \oplus c' = m \oplus m'$ , where  $c=m \oplus k$  and  $c'=m' \oplus k$  since  $k \oplus k=0$  and  $k \oplus 0=k$ .

Unbreakable cipher used during the cold war: **unconditional security**

Main drawbacks:

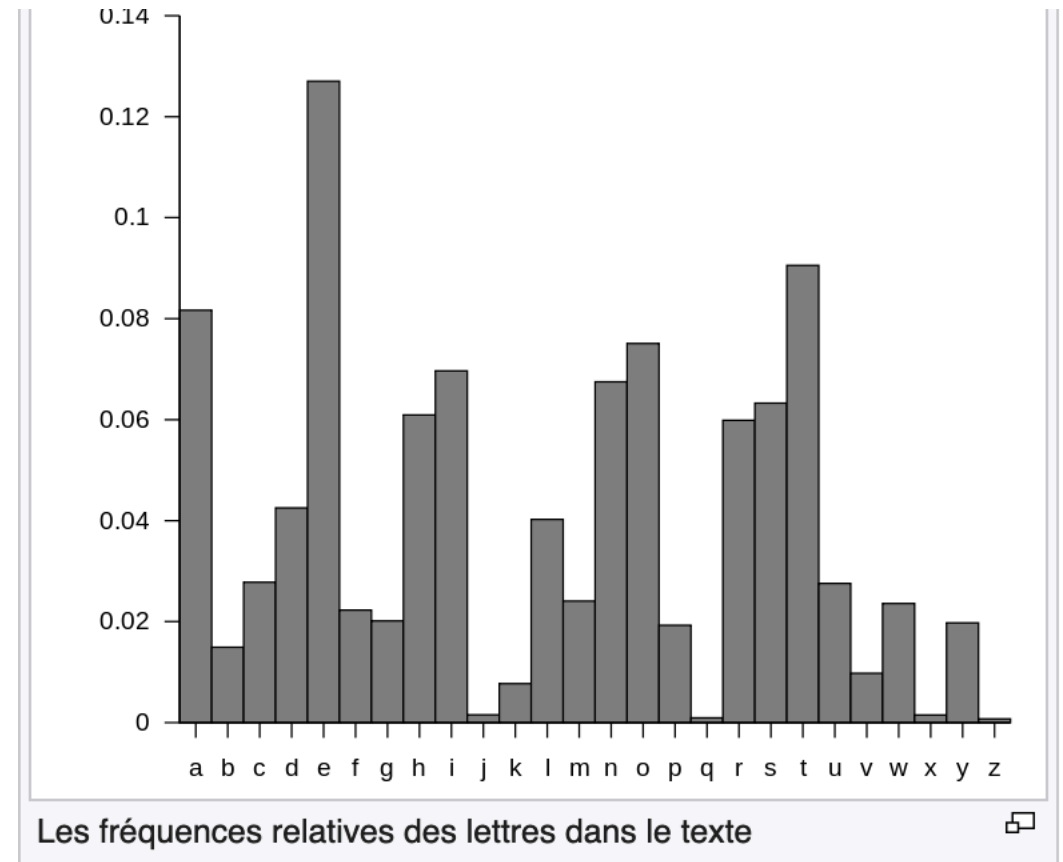
1. the key must be random,
2. key must be as long as the message, and
3. key must be changed for each message.

# Unconditional vs. Computational Security

- A powerful adversary with infinite time cannot obtain information about the plaintext given only the ciphertext
- If a ciphertext  $c$  is obtained, anyone can produce a plaintext  $p$  and a key  $k$  such that  $c=p\oplus k$  for any plaintext  $|p|=|c|$
- The adversary cannot distinguish plaintext with equal length
- However, in practice keys are reused across several ciphertexts
- In a practical point of view, **computational security** is preferred: it is computationally hard to recover the plaintext (but possible for an adversary with infinite time...).
- E.g.: for a block cipher, we can exhaust all keys  $2^{128}$  operations

# Vigenere Cipher

- Vigenere is based on the same idea as Vernam with {a,b,c,...z} alphabet or ascii characters
- Vigenere cipher reuses the same key
- IC (index of coincidence) =  $\sum_{a \in A} p_a$  where  $p_a$  is the probability of character a
- IC is invariant with substitution
- IC is higher when the distribution is **not** uniform



# Cryptanalysis Vigenere Cipher

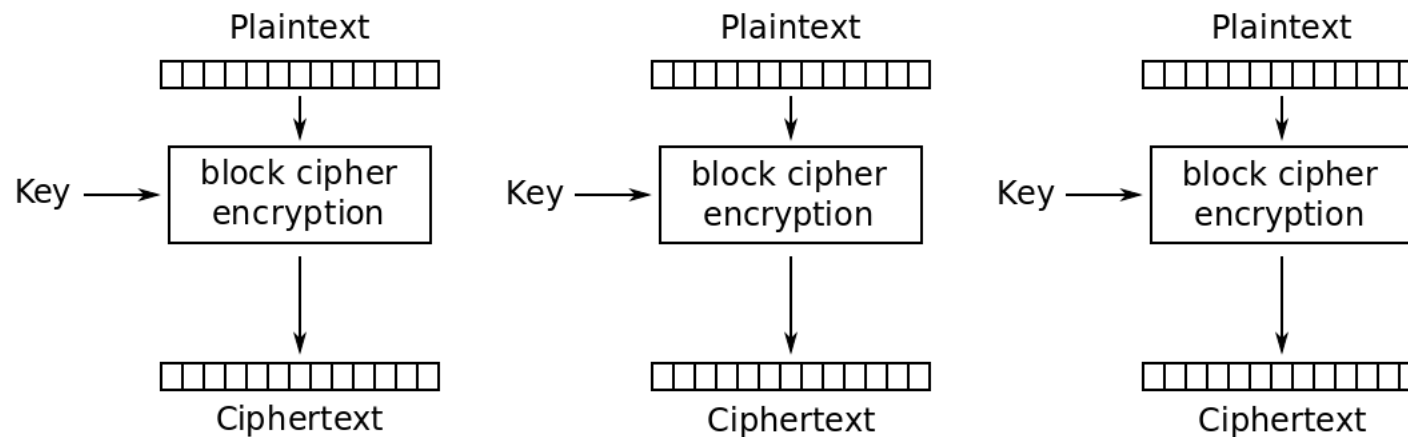
- Assume the key length is known, one can extract substrings encrypted with the same letter
- Such encryption is called a shift encryption since the whole alphabet is shifted
- Easier to break than substitution: once the encryption of one letter is known, we can deduce all the substitution
- Learning the length: Guess all length and compute the IC
- ICM:  $\sum_{a \in A} p_a p'_a$  where  $p_a$  and  $p'_a$  are the probabilities of two strings

# The key space

- The size of the key space is the number of encryption/decryption key pairs available in the cipher system. A **key is a compact way to specify the encryption function** (from the set of all encryption functions).
- E.g. a **substitution of block length  $t$  has  $(2^t)!$  encryption functions**
- A necessary, but usually not sufficient, condition for an encryption scheme to be secure is that the **key space be large enough to preclude exhaustive search**. E.g.  $26! \approx 4 \times 10^{26}$ .

# Encrypting long messages

- Mode of operations: how using a block cipher to encrypt large messages ?



Electronic Codebook (ECB) mode encryption

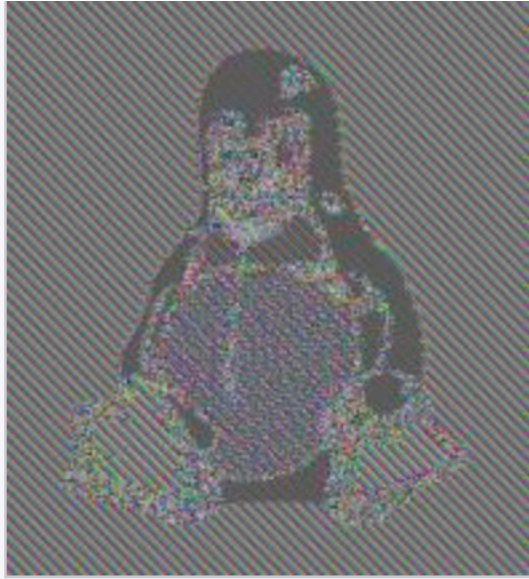
Problems:  
- deterministic

Other modes: CFB, OFB, PCBC, CTS (Ciphertext stealing)

# Electronic Code Book is deterministic ... CBC better

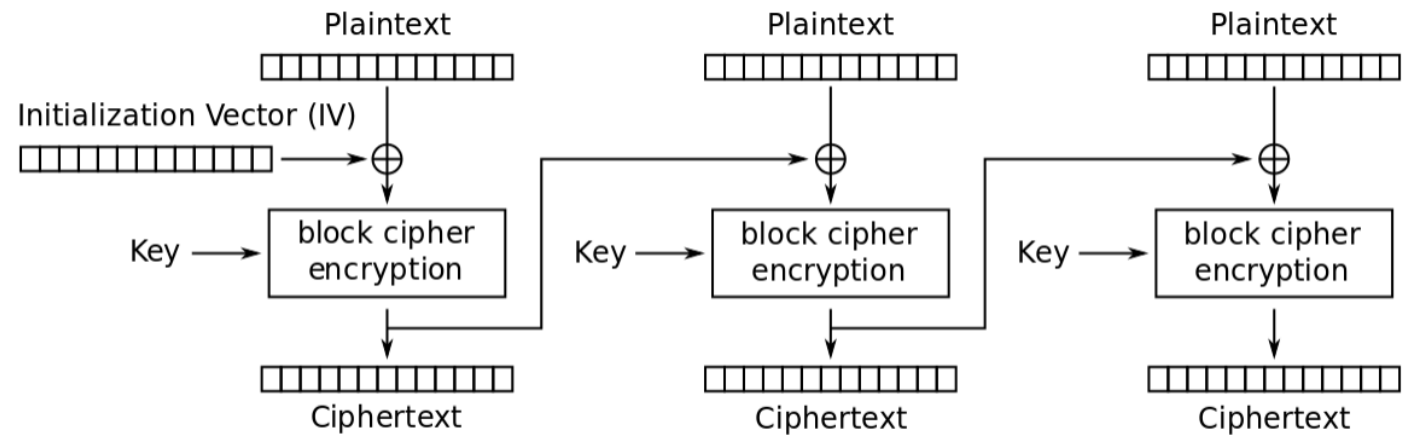


Original image



Encrypted using ECB mode

Randomization is useful ...



Cipher Block Chaining (CBC) mode encryption

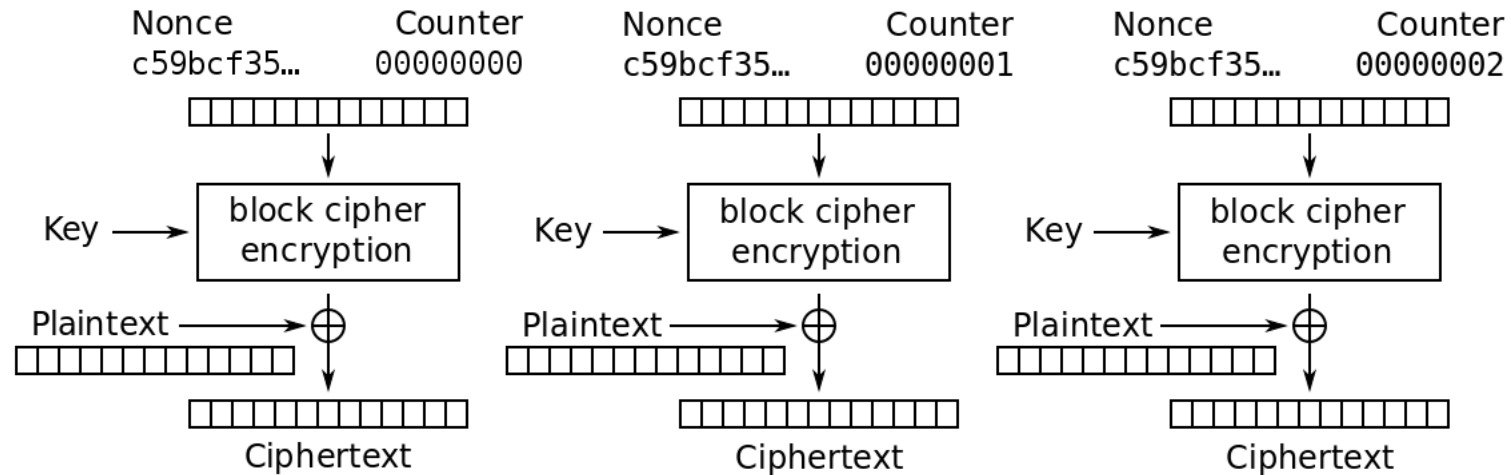


# Stream Cipher and CTR mode of operations

Vernam cipher is unconditionally secure

Main problem: key reuse

- Generate the key with a smaller one using a pseudorandom number generator : output looks random but bitstring generated deterministically with from a secret seed
- Cryptographically secure pseudorandom generator are hard to design: rand from c language is not good
- Block cipher can be used as follows
- If stream cipher are resynchronized, same key is generated (WPA)
- In order to make it stateless, a nonce is usually added to generate different keystream



Counter (CTR) mode encryption