

Symmetric Crypto Block Cipher

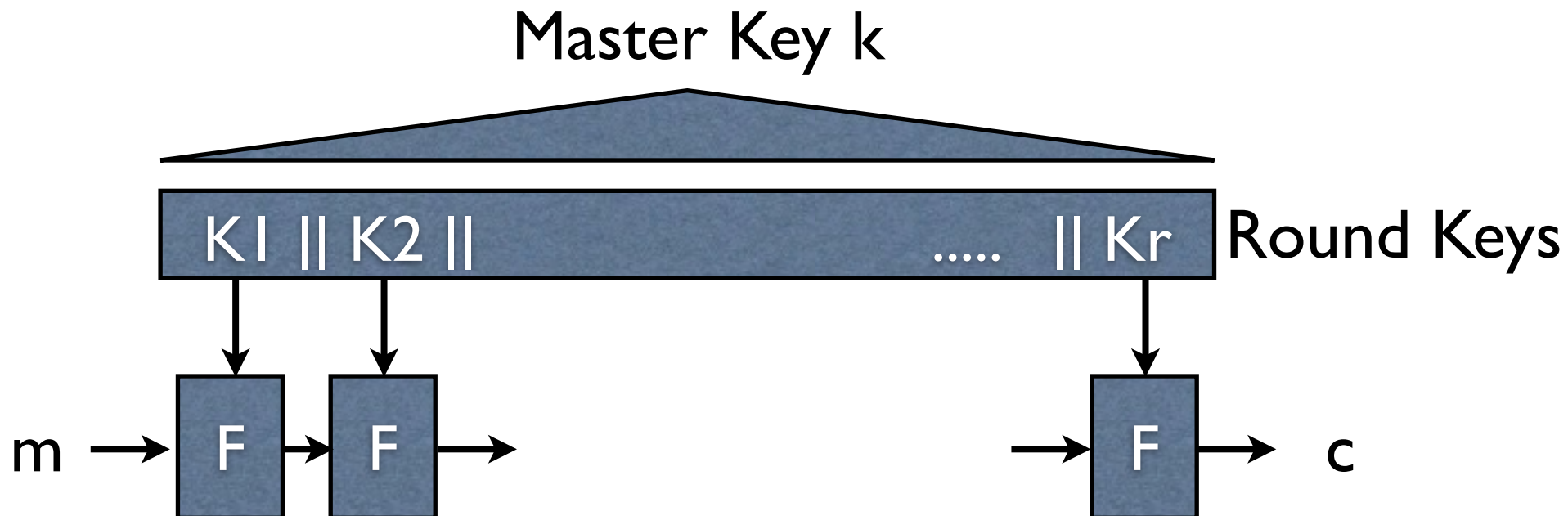
Pierre-Alain Fouque

Family of permutations

- Block cipher = family of 2^k substitutions on an alphabet of size 2^n indistinguishable from randomly chosen substitutions
- Constructions:
 - Feistel scheme (DES)
 - SPN: Substitution-Permutation Network (AES)

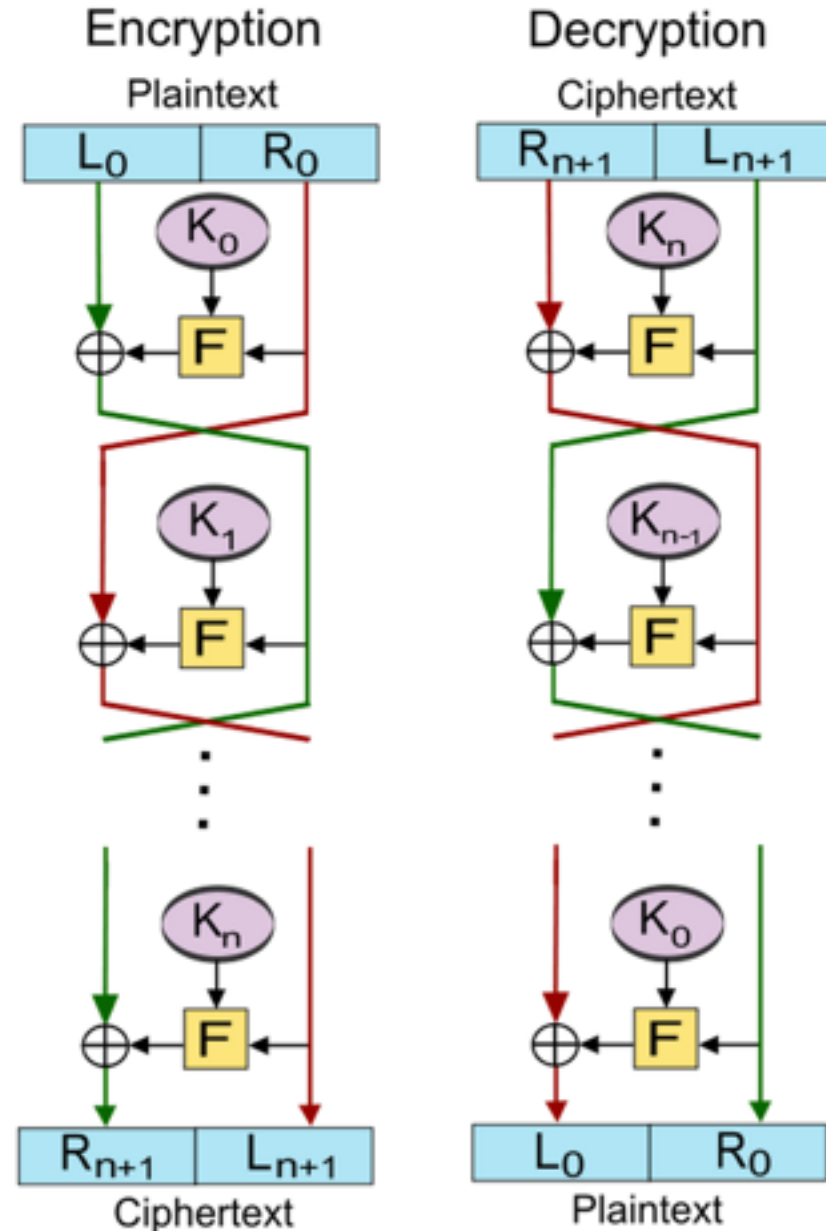
Block cipher

- Cipher (E,D) «eff. algs» such that $D(k,E(k,m))=m$
- Main drawback of stream cipher: lacks of theory to construct secure PRG
- Iterate many times a «small» round function F



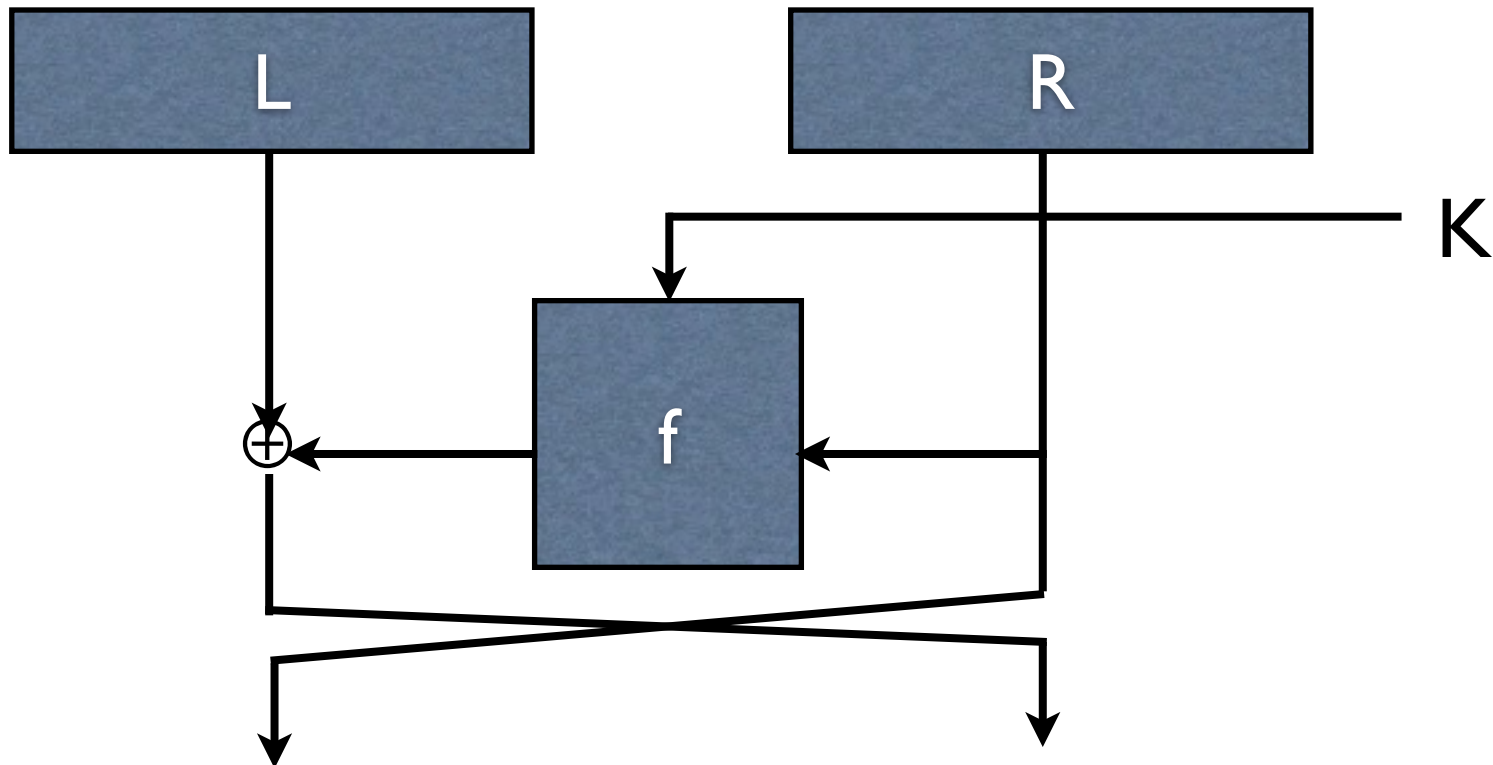
Data Encryption Standard

- DES (IBM 1973) and NBS (1977)
- Key Length: 56 bits
- Block Length: 64 bits
- 16 rounds with 48-bit round keys



Feistel scheme

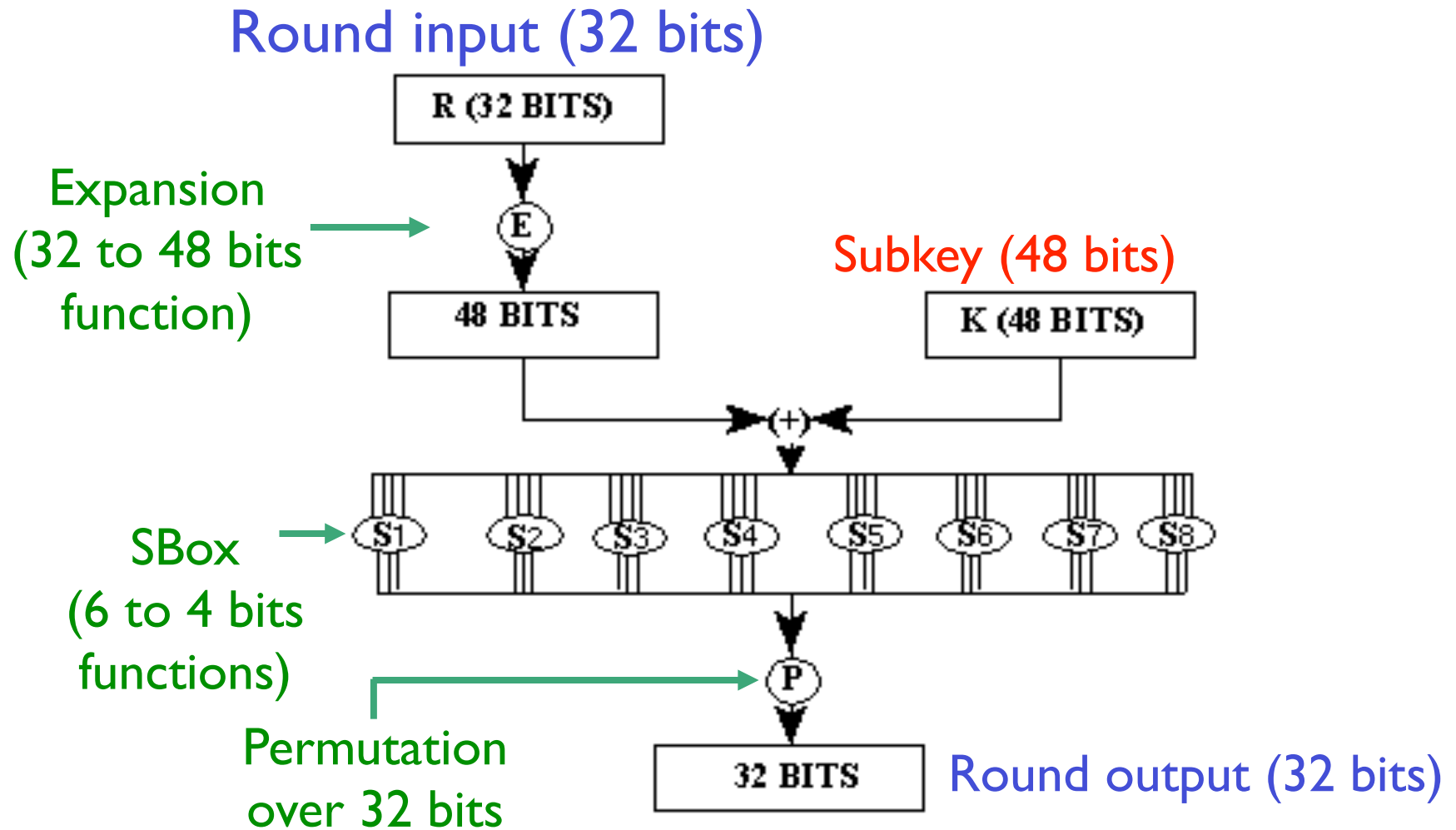
- Designed by Horst Feistel at IBM
- Transform random function to random permutation



Feistel security

- Could you distinguish one-round Feistel ?
- Could you distinguish two-round Feistel ?
- Could you distinguish three-round Feistel ?

f function



Attacks against DES

- Before 1990: attacks against round reduced version (less than 16 rounds)
- 1990-92: **Differential** cryptanalysis
- 1993-94: **Linear** cryptanalysis
- other attacks: Davies-Murphy, side-channel
- In **practice, the most efficient attack is the exhaustive search (EFF, copacabana)**

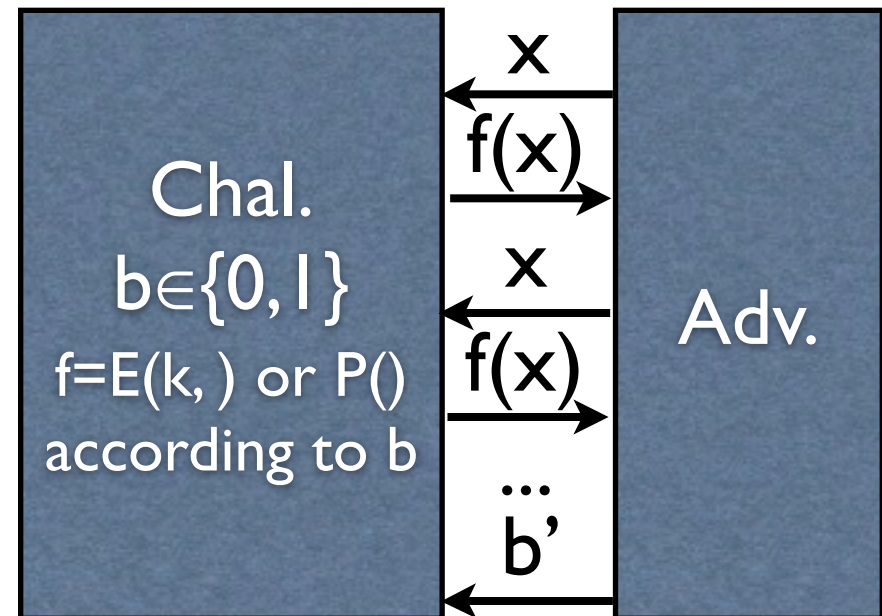
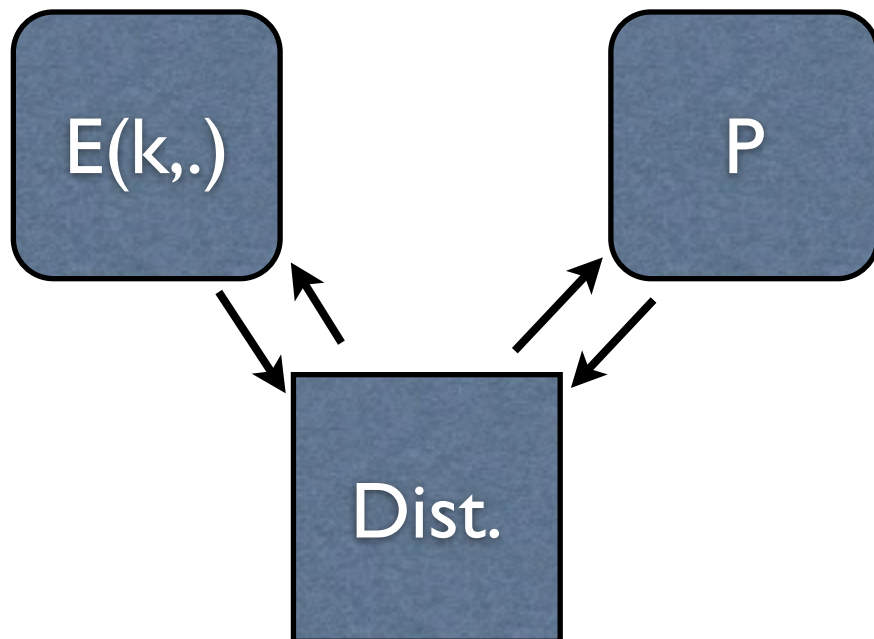


Main drawback of DES

- Exhaustive key search in 2^{56} (3DES)
- Block size (collision for 2^{32} blocks)
- Differential / Linear Cryptanalysis
- DES: well-designed and withstands successfully 30 years of cryptanalysis

Security game

- Block cipher must be indistinguishable from a random permutation
- for all k , $E(k,x)$ is a permutation which looks random provided the key is not known



$$\text{Adv}(A) = |\Pr[b = b'] - 1/2|$$

Birthday Paradox

- Randomly and independently chose k balls among N balls, whenever $k \approx \sqrt{N}$, $\Pr[\text{coll}] > 1/2$
- If $N=365$, the probability that two people same birthday is $1/2$ when $k \approx 23$
- $\Pr(\text{no coll with } k \text{ balls}) = (1-1/N)(1-2/N)\dots(1-(k-1)/N) < \exp(\sum_i i/N) = \exp(k(k-1)/2N)$
- if $k \approx \sqrt{N}$, $\Pr(\text{coll}) > 1/2$
- Two sets N and M of random elements in a set D :
Number of expected collisions is $|N| \times |M| / |D|$
(Birthday paradox with boys and girls)

2DES → 3DES

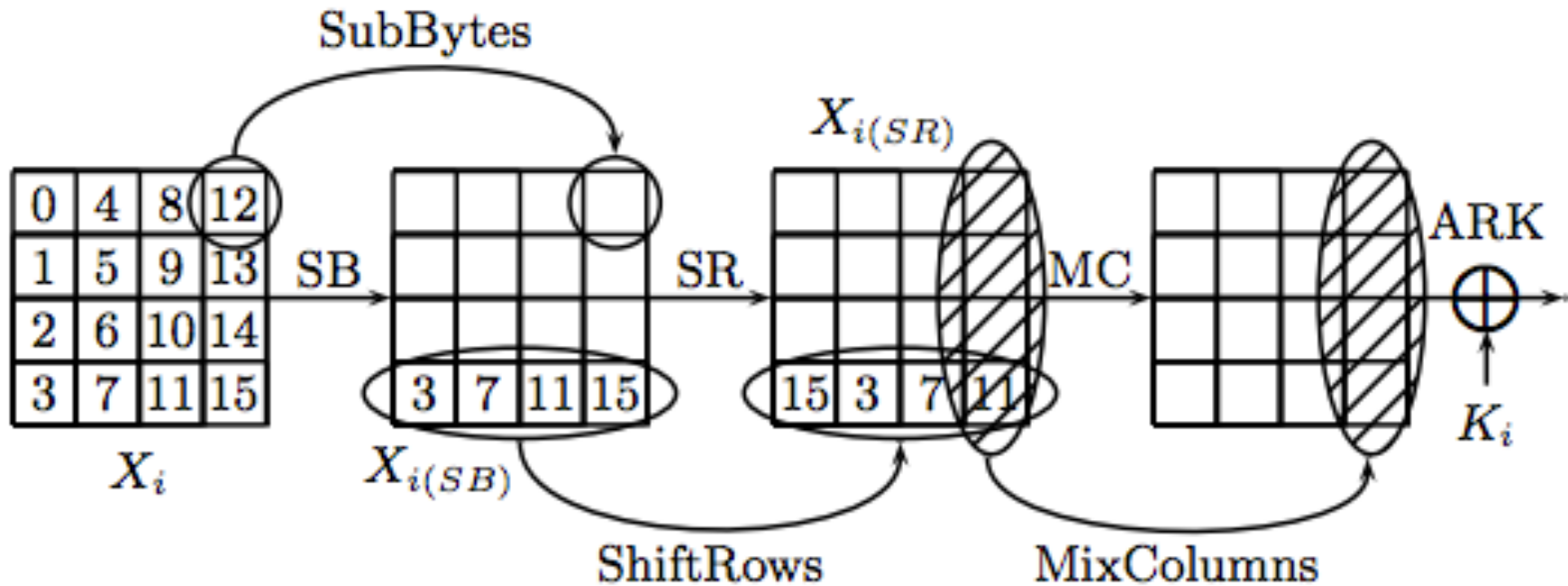
Differential Attack

- Consider an Even-Mansour Scheme

Advanced Encryption Standard

- Substitution / Permutation Network
- Key Length: 128 / 192 / 256 bits
- Round numbers: 10 / 12 / 14 (last round w/o MixColumns)
- Block Length: 128 bits
- Designed by Daemen and Rijmen
- Standardized by NIST in 2000

AES



Diffusion: 2 tours

Attaque sur 6 tours en 2^{48}

Attaque sur 7 tours en 2^{100}

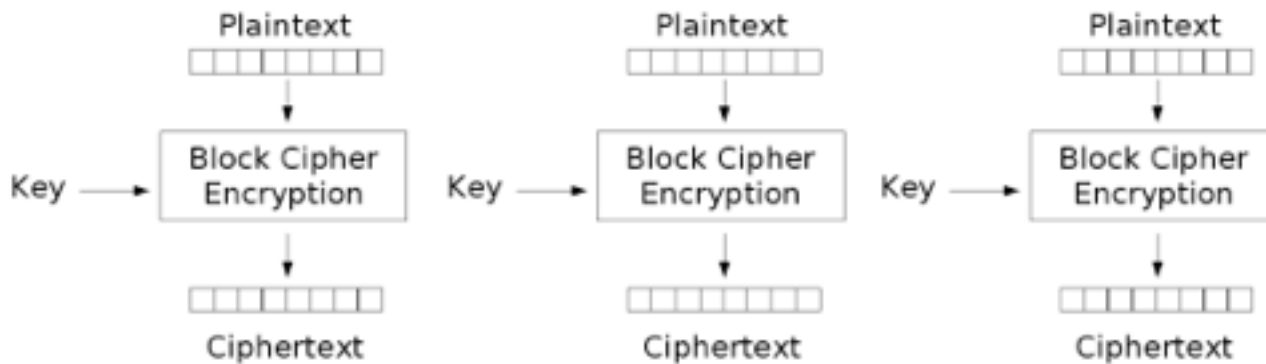
Pas d'attaque au-delà mais en attaque par canaux auxiliaires

Security

- MDS Property of the MixColumn
- Diffusion: In 2 rounds, any difference affect the whole cipher
- Security against Differential / Linear Cryptanalysis: good properties of the SBox
- Square Attack
- Side-channel Attack (timing, Power Analysis) are the most efficient attacks against AES

Modes of operation

- How to encipher larger messages ?
 - ECB, CBC, CTR, OFB, CFB



Electronic Codebook (ECB) mode encryption

Drawbacks:

- deterministic

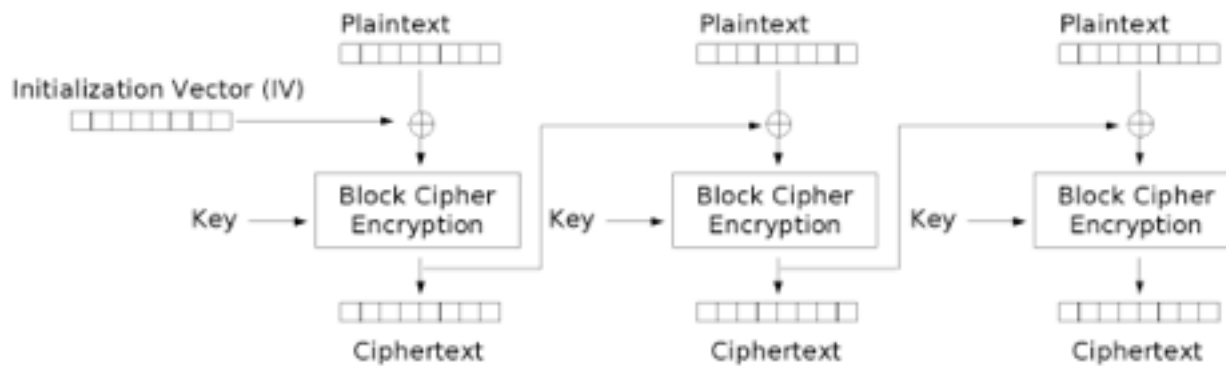
Advantages:

- parallelisable

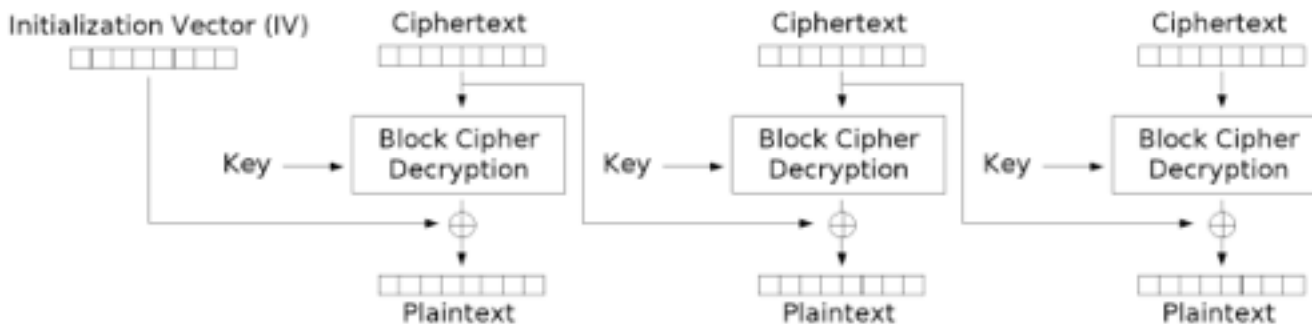


Ciphertext Block Chaining (CBC)

- Encrypting: $C_0=IV, \dots, C_i=E(k, C_{i-1} \oplus M_i)$
- Decrypting: $M_i=D(k, C_i) \oplus C_{i-1}$



Cipher Block Chaining (CBC) mode encryption



Cipher Block Chaining (CBC) mode decryption

Drawbacks:

- sequential

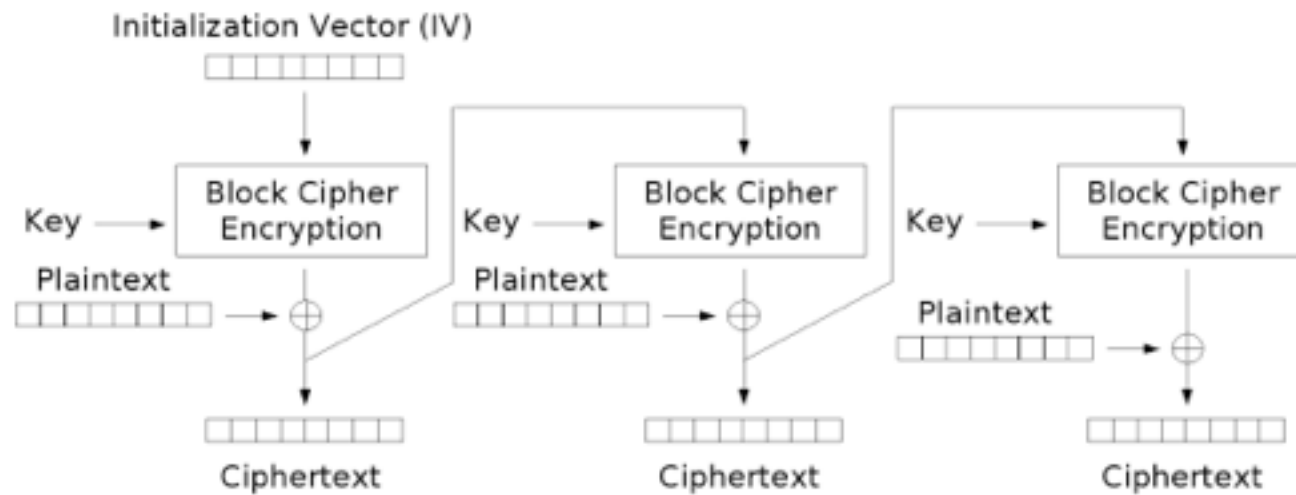
Advantages:

- randomized

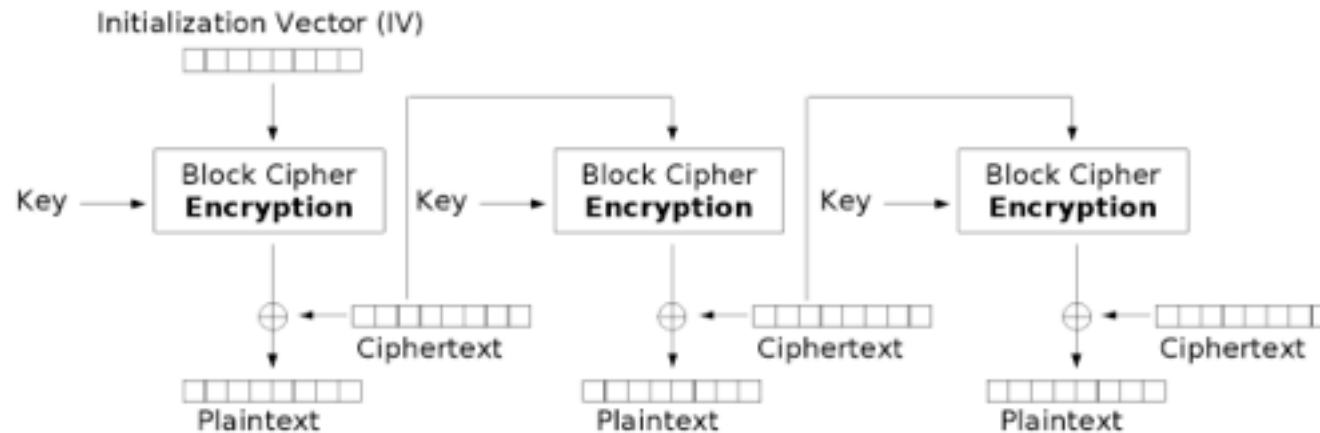
- propagation of error in decryption

Ciphertext FeedBack (CFB)

- How to use a block cipher as a stream cipher ?



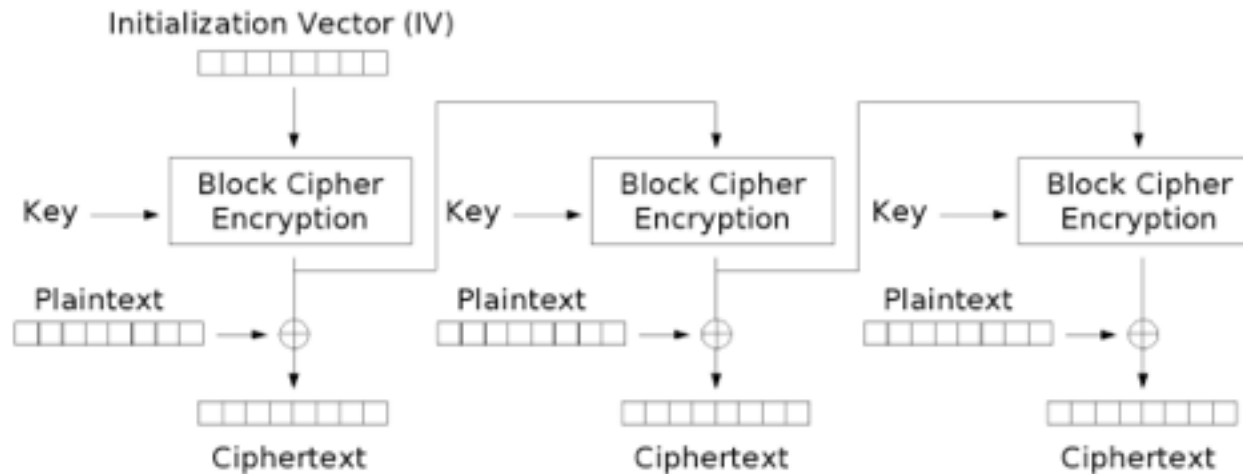
Cipher Feedback (CFB) mode encryption



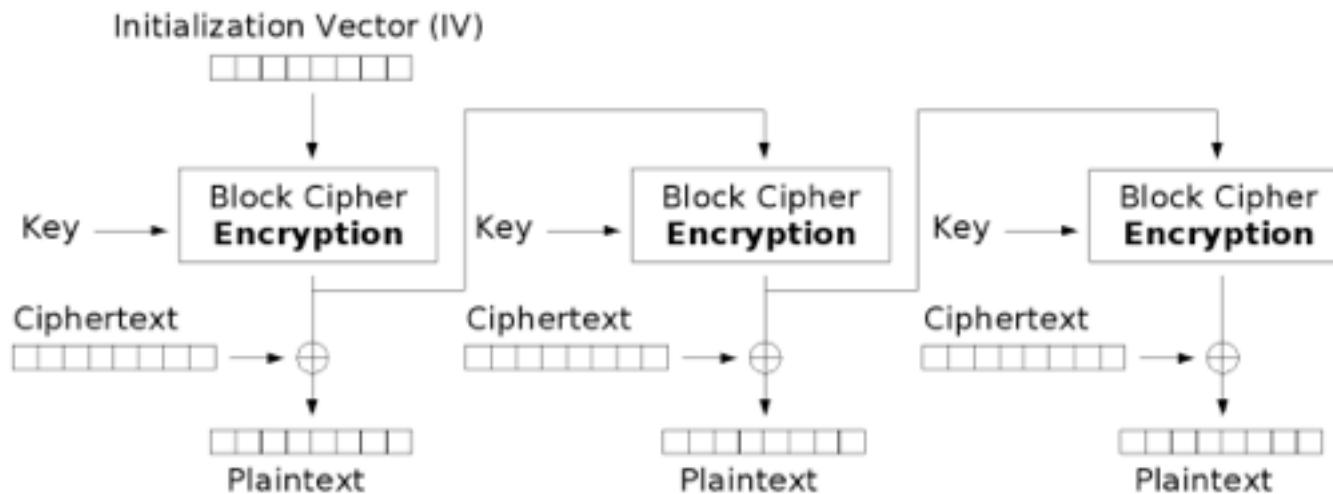
Cipher Feedback (CFB) mode decryption

Output FeedBack (OFB)

- How to use a block cipher as a stream cipher ?



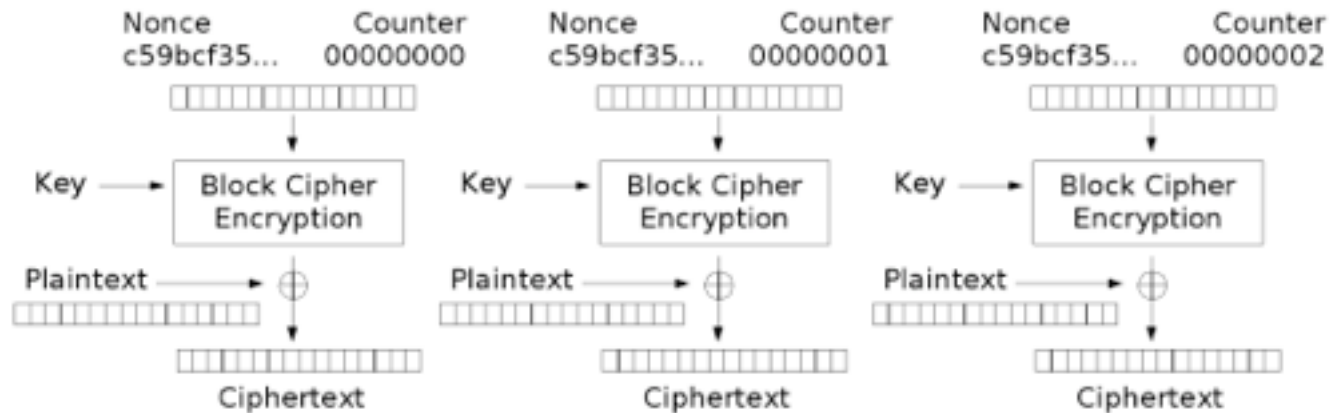
Output FeedBack (OFB) mode encryption



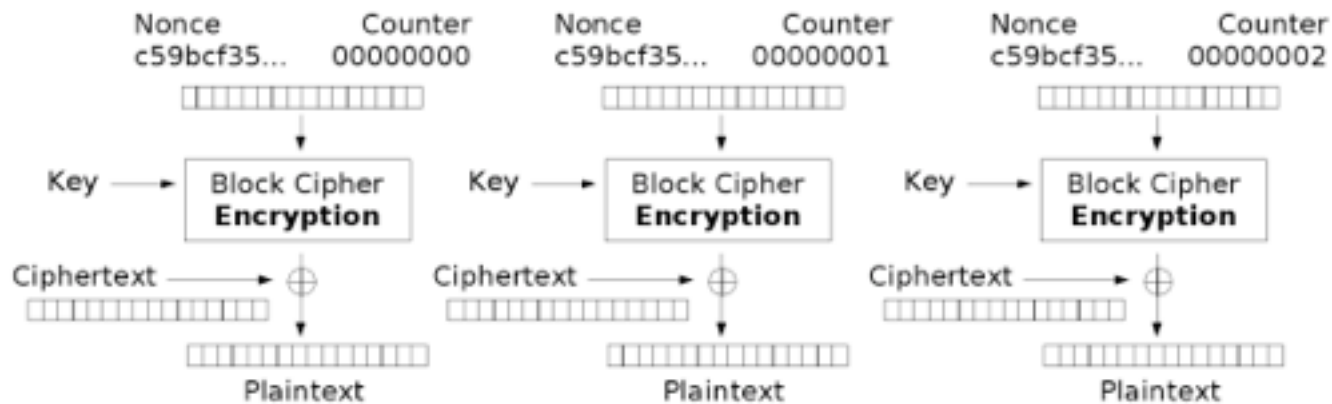
Output FeedBack (OFB) mode decryption

Counter Mode (CTR)

- Better solution



Counter (CTR) mode encryption



Counter (CTR) mode decryption

Security

- Confidentiality is ensure by the mode of operation
- Integrity: first block of CBC ?
- Main idea: the ciphertext must be indistinguishable from random for **polynomial-time adversaries**
- **Security Game:**
- **Example on CBC:**