

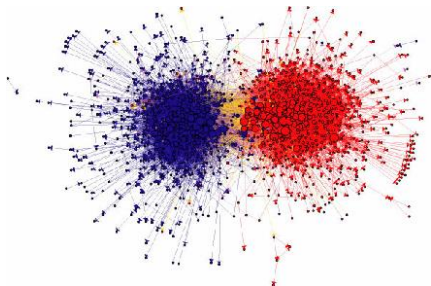
Inference in large random graphs

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Community Detection: cluster nodes $i \in V$ of graph $G = (V, E)$, into subsets V_1, \dots, V_K of “statistically similar nodes”



Applications:

- recommendation (graph between users and products);
- biology (graph of chemical interactions between proteins);
- ...

→ Focus on $G = (V, E)$ drawn from the **Stochastic Block Model** $\mathcal{G}(n, \alpha, P)$:

- ▶ $\alpha = \{\alpha_1, \dots, \alpha_K\}$ probability distribution on $[K]$
- ▶ $P \in [0, 1]^{K \times K}$: Symmetric matrix
- ▶ $\{\sigma_i\}_{i \in [n]}$: i.i.d., $\sim \alpha$
- ▶ Conditionally on $\sigma_{[n]} := \{\sigma_1, \dots, \sigma_n\}$, edge $(i, j) \in E$ with probability P_{σ_i, σ_j}

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→ Generative model that generalizes the Erdős-Rényi random graph $\mathcal{G}(n, p)$

→ Community Detection: inference of node attributes σ_i from observation $G = (V, E)$

→ Spectral methods and phase transitions on feasibility of community detection

Graphical models, a generalization of the Ising model on $\{-1, 1\}^n$:

$$\mathbb{P}(X_{[n]} = x_{[n]}) \propto e^{-h \sum_{i \in [n]} x_i + \sigma \sum_{i,j: i \sim j} x_i x_j}$$

Goal: Infer unobserved variables X_i of nodes i of graph G

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Goal: Infer unobserved variables X_i of nodes i of graph G

→ Belief Propagation Algorithm

→ Emphasis on **Tree Reconstruction** problem: infer characteristics X_a where a : ancestor in genealogical tree, from characteristics of its descendants

→ Phase transitions on feasibility of inference, and links to community detection

Hypothesis tests: does observed graph G have some structure (ex: is drawn from Stochastic Block Model with $K > 1$ blocks) or is it “totally random”, i.e. an Erdős-Rényi graph?

Special Case: the Planted Clique Problem

→ Highlight existence of two kinds of phase transitions:

- ▶ **Informational**, i.e. is there enough information present in the observation
- ▶ **Computational**, i.e. can the information present in the observation be extracted in polynomial time

Outline for today

- ▶ Background results from Linear Algebra
- ▶ Bounds on spectral radius of random matrices
- ▶ Spectral methods for community detection in Stochastic Block Model

Singular Value Decomposition (SVD)

SVD of matrix $X \in \mathbb{C}^{n \times p}$:

$$X = U \Lambda V^* = \sum_{i=1}^{n \wedge p} \sigma_i u_i v_i^*, \text{ where:}$$

$$U = (u_1, \dots, u_n) \in \mathbb{C}^{n \times n}, U^* U = I_n, (u_i: i\text{-th left singular vector)}$$

$$V = (v_1, \dots, v_p) \in \mathbb{C}^{p \times p}, V^* V = I_p, (v_i: i\text{-th right singular vector})$$

$$\Lambda \in \mathbb{R}^{n \times p}, \Lambda = \begin{pmatrix} \sigma_1 & 0 & \dots & 0 \\ 0 & \ddots & & \vdots \\ 0 & 0 & \sigma_{n \wedge p} & 0 \end{pmatrix},$$

$$\sigma_1 \geq \dots \geq \sigma_{n \wedge p} \geq 0, (\sigma_i: i\text{-th singular value})$$

Principal Component Analysis (Karl Pearson “On Lines and Planes of Closest Fit to Systems of Points in Space”, 1901)

Definition

For matrix $X \in \mathbb{C}^{n \times p}$ with SVD $X = \sum_{i=1}^{n \wedge p} \sigma_i u_i v_i^*$,

Operator norm: $\|X\|_{op} := \sup_{u \in \mathbb{C}^p} \frac{\|Xu\|}{\|u\|} = \sigma_1$;

Frobenius norm: $\|X\|_F := \sqrt{\sum_{i,j} X_{ij}^2} = \sqrt{\sum_i \sigma_i^2}$.

For $r < n \wedge p$, let $X_r := \sum_{i=1}^r \sigma_i u_i v_i^*$. Then:

X_r best rank- r approximation of X both for $\|\cdot\|_F$ and $\|\cdot\|_{op}$, with

$$\inf_{Y: rk(Y)=r} \|X - Y\|_{op} = \|X - X_r\|_{op} = \sigma_{r+1},$$

$$\inf_{Y: rk(Y)=r} \|X - Y\|_F = \|X - X_r\|_F = \sqrt{\sum_{i>r} \sigma_i^2}.$$



Perturbation results: Eigenvalues

Definition

For $M \in \mathbb{C}^{n \times n}$,

spectral radius $\rho(M) := \sup\{|\lambda|, \lambda \in \text{Spectrum}(M)\}$.

For Hermitian $M \in \mathbb{C}^{n \times n}$, order its (real) eigenvalues as $\lambda_1(M) \geq \dots \geq \lambda_n(M)$.

Then $\rho(M) = \sigma_1(M) = \max(|\lambda_1(M)|, |\lambda_n(M)|)$.

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Lemma

(Weyl's inequality) For Hermitian H, W in $\mathbb{C}^{n \times n}$, for all $i \in [n]$, $|\lambda_i(H) - \lambda_i(\overline{H+W})| \leq \rho(W)$

Proof: by Courant-Fisher theorem,

$$\lambda_i(H + W) = \sup_{\dim(E)=i} \inf_{x \in E, \|x\|=1} x^T (H + W)x$$

Apply to $E = \text{Vect}\{x_1(H), \dots, x_i(H)\}$ to obtain

$$\begin{aligned} \lambda_i(H + W) &\geq \inf_{x \in E, \|x\|=1} x^T (H + W)x \\ &\geq \inf_{x \in E, \|x\|=1} x^T Hx + \inf_{x \in E, \|x\|=1} x^T Wx \\ &\geq \lambda_i(H) - \rho(W). \end{aligned}$$

Similarly, $\lambda_i(H) \geq \lambda_i(H + W) - \rho(-W)$ hence the result

Cauchy's interlacing theorem

Theorem

For Hermitian $A \in \mathbb{C}^{n \times n}$, $P \in \mathbb{C}^{n \times n}$, $m < n$ such that $P^*P = I_m$,
then $B := P^*AP$ verifies

$$\lambda_i(A) \geq \lambda_i(B) \geq \lambda_{n-m+i}(A)$$

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Proof: Let $x_1(B), \dots, x_n(B)$ eigenvector basis of B ,
 $V_i = \text{Vect}(x_1(B), \dots, x_i(B))$, $W_i = P(V_i)$.

Courant-Fisher:

$$\lambda_i(B) = \inf_{u \in V_i, \|u\|=1} u^* P^* A P u = \inf_{v \in W_i, \|v\|=1} v^* A v \leq \lambda_i(A).$$

Similarly:

$$\lambda_{m-i+1}(-A) \geq \lambda_{m-i+1}(-B); \text{ equivalently:} \\ -\lambda_{n-m+i}(A) \geq -\lambda_i(B).$$

Perturbation of eigenvectors: the Davis-Kahane “sin Θ ” theorem

Theorem

For symmetric, real $H, \hat{H}, W = \hat{H} - H \in \mathbb{R}^{n \times n}$ any $1 \leq r \leq s \leq n$, $d = s - r + 1$, if

$$\delta := \min(\lambda_{r-1}(H) - \lambda_r(H), \lambda_s(H) - \lambda_{s+1}(H)) > 0,$$

then for matrices $V = (v_r, \dots, v_s)$, $\hat{V} = (\hat{v}_r, \hat{v}_s)$ of orthonormal eigenvectors of H resp. of \hat{H} , there exists $O \in \mathbb{R}^{d \times d}$ such that $O^\top O = I_d$ and

$$\|\hat{V}O - V\|_{op} \leq \|\hat{V}O - V\|_F \leq \frac{2^{3/2} \min(d^{1/2} \|W\|_{op}, \|W\|_F)}{\delta}.$$

Perturbation of eigenvectors: a more elementary result

Lemma

For fixed i let $\Delta := \inf_{j:\lambda_j(H) \neq \lambda_i(H)} |\lambda_j(H) - \lambda_i(H)|$. Assume $\rho(W) < \Delta$. Then for any normed eigenvector \hat{x}_i of $\hat{H} = H + W$ associated with $\lambda_i(\hat{H})$ there exists x_i normed eigenvector of H associated with $\lambda_i(H)$ such that $\langle x_i, \hat{x}_i \rangle \geq \sqrt{1 - \left(\frac{\rho(W)}{\Delta - \rho(W)}\right)^2}$

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Proof: Decompose $\hat{x}_i = \sum_j \theta_j x_j$ to get

$$\hat{H}\hat{x}_i = \lambda_i(\hat{H})\hat{x}_i = \sum_j \theta_j \lambda_j(H)x_j + W\hat{x}_i$$

$$\text{hence } W\hat{x}_i = \sum_j (\lambda_i(\hat{H}) - \lambda_j(H))\theta_j x_j$$

By Weyl's inequality, $|\lambda_i(\hat{H}) - \lambda_j(H)| \geq \Delta - \rho(W)$ if $\lambda_i(H) \neq \lambda_j(H)$. Thus

$$\rho(W) \geq (\Delta - \rho(W)) \sqrt{1 - \sum_{k:\lambda_k(H) = \lambda_i(H)} |\theta_k|^2}$$

$$\Rightarrow \sum_{k:\lambda_k(H) = \lambda_i(H)} |\theta_k|^2 \geq 1 - \left(\frac{\rho(W)}{\Delta - \rho(W)}\right)^2$$

Perturbation arguments recap

For $\hat{H} = H + W$, eigenvalue $\lambda_i(\hat{H})$ and associated eigenvector \hat{x}_i of \hat{H} , with $\Delta = \inf_{\lambda_j(H) \neq \lambda_i(H)} |\lambda_j(H) - \lambda_i(H)|$:

$$|\lambda_i(H) - \lambda_i(\hat{H})| \leq \rho(W)$$

$$\exists x_i \leftrightarrow \lambda_i(H) \text{ such that } \langle x_i, \hat{x}_i \rangle = 1 - O\left(\frac{\rho(W)^2}{\Delta^2}\right).$$

Spectral properties of graphs

Non-oriented graph $G = (V, E)$, $V = [n]$, $E = \{\text{edges } (i, j)\}$

Adjacency matrix $A(G)$: $A_{ij} = \mathbb{I}_{(i,j) \in E} = \mathbb{I}_{i \sim j}$, $i, j \in V = [n]$

→ Symmetric for non-oriented graph

Laplacian matrix $L(G)$: $L_{ij} = \begin{cases} -A_{ij} & \text{if } i \neq j, \\ d_i := \sum_{k \neq i} A_{ik} & \text{if } j = i. \end{cases}$

d_i : **degree** of node i .

$x^T L x = \sum_{i < j} A_{ij} (x_i - x_j)^2$ so that:

$0 \preceq L$, where \preceq : semi-definite order on symmetric matrices, and

$\lambda_n(L) = 0$ with associated eigenvector $\{1/\sqrt{n}\}_{i \in [n]}$

L : infinitesimal generator of continuous-time random walk on G ,
with transition rates A_{ij} from i to $j \neq i$

Isoperimetric constant of graph G :

$$I(G) := \min \left\{ \frac{|E(S, \bar{S})|}{|S|}, S \subset V, 0 < |S| \leq \frac{n}{2} \right\}$$

Let $\Delta(G) := \max_{i \in V=[n]} d_i(G)$, largest node degree.

Lemma

(Cheeger's inequality). $I(G) \leq \sqrt{2\Delta(G)\lambda_{n-1}(L(G))}$.

Easier result:

Lemma

$$I(G) \geq \frac{\lambda_{n-1}(L(G))}{2}$$

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Proof: Courant-Fisher: $\lambda_{n-1} = \inf \left\{ \frac{x^T L x}{\|x\|^2}, x : \langle x, e \rangle = 0 \right\}$

For $x : x_i = \mathbb{I}_{i \in S} - \frac{|S|}{n}$, yields

$$\lambda_{n-1} \leq \frac{|E(S, \bar{S})|}{|S|(1-|S|/n)}$$

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Corollary

Graph connected iff $\lambda_{n-1}(L) > 0$

Alon-Boppana theorem

Definition

d -regular graph G is **Ramanujan** if $\lambda_2(A) \leq 2\sqrt{d-1}$

Note: $L = dI_n - A$, hence $\lambda_{n-1}(L) = d - \lambda_2(A)$

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Theorem

For d -regular G with diameter $\geq 2r + 1$,
 $\lambda_2(A) \geq 2\sqrt{d-1} \cos\left(\frac{\pi}{r+2}\right) = 2\sqrt{d-1}(1 - O(r^{-2}))$.

Hence Ramanujan graph G has maximal spectral gap $\lambda_{n-1}(L)$

Proof

Take $i, j \in V : d_G(i, j) \geq 2r + 1$.

G' : subgraph of G induced by nodes of $\mathcal{B}(i, r), \mathcal{B}(j, r)$

$A'_G = P^* A_G P$, where $P \in \mathbb{R}^{n \times m}$ such that $P^* P = I_m$

Cauchy's interlacing theorem

$$\Rightarrow \lambda_2(A_G) \geq \lambda_2(A_{G'}) = \min\{\lambda_1(A_{\mathcal{B}(i,r)}), \lambda_1(A_{\mathcal{B}(j,r)})\}$$

Nb of closed walks of length $2q$ started at i in $\mathcal{B}(i, r)$: $w_{2q}(\mathcal{B}(i, r))$.

$$w_{2q}(\mathcal{B}(i, r)) \geq w_{2q}(\mathcal{T}_{d-1,r}) = (d-1)^q w_{2q}(\mathcal{P}_r)$$

$$\lim_{q \rightarrow \infty} [w_{2q}(\mathcal{B}(i, r))]^{1/2q} = \lambda_1(A_{\mathcal{B}(i,r)}) \geq \dots$$

$$\dots \geq \sqrt{d-1} \lambda_1(\mathcal{P}_r) = \sqrt{d-1} \times 2 \cos \frac{\pi}{r+2}$$

Bounding spectral radius $\rho(W)$ of random matrices W : the Trace method

(method due to [Füredi-Komlos'81]; here a sub-optimal version avoiding sharp combinatorics)

Lemma

Let $W \in \mathbb{R}^{n \times n}$: symmetric matrix with entries independent up to symmetry, such that $|W_{ij}| \leq 1$, $\mathbb{E}(W_{ij}) = 0$, and $\mathbb{E}(W_{ij}^2) \leq O(d/n)$ for some $d \geq 1$. Then for any fixed $\epsilon > 0$, with high probability, $\rho(W) \leq O(\sqrt{dn}^\epsilon)$.

Corollary

If we assume $d \geq n^\delta$ for some $\delta \in]0, 1[$, then with high probability $\rho(W) = o(d)$.

Proof: Take $\epsilon < \delta/2$ to obtain $\rho(W) = O(n^{\delta/2+\epsilon}) = o(n^\delta)$.

Proof

For fixed $k \in \mathbb{N}$, write $\rho^{2k} \leq \sum_{i \in [n]} \lambda_i(W)^{2k} = \text{Trace}(W^{2k})$

Thus $\mathbb{P}(\rho \geq x) \leq x^{-2k} \mathbb{E}(\rho^{2k}) \leq x^{-2k} \mathbb{E} \text{Trace}(W^{2k})$

Combinatorial expression of trace:

$$\text{Trace}(W^{2k}) = \sum_{i_0^{2k} \in [n]^{2k+1}: i_0 = i_{2k}} \prod_{j=1}^{2k} W_{i_{j-1} i_j}$$

Recall W_{ij} : centered and independent

→ Only paths contributing non-zero expectation: traverse each edge at least twice

$$\Rightarrow \mathbb{E} \text{Trace}(W^{2k}) \leq \sum_{e=1}^k \sum_{v=1}^{e+1} C(e, v) n^v O((d/n)^e)$$

Yields $\mathbb{E} \text{Trace}(W^{2k}) = O(nd^k)$.

For $x = \sqrt{dn^\epsilon}$, yields $\mathbb{P}(\rho \geq x) \leq O(n^{1-2k\epsilon})$

Result follows by taking $k > 1/(2\epsilon)$

Method 2: Bernstein's inequality for matrices

Theorem

(Tropp'15) Let X_1, \dots, X_m be independent Hermitian random matrices such that:

$$\mathbb{E}(X_k) = 0, \quad \|X_k\|_{op} \leq L \text{ almost surely, } k \in [m].$$

Let $Y = \sum_{k \in [m]} X_k$, and $v(Y) := \|\mathbb{E}(Y^2)\|_{op} = \|\sum_{k \in [m]} \mathbb{E}X_k^2\|_{op}$.

Then for all $t > 0$, $\mathbb{P}(\lambda_1(Y) \geq t) \leq n \exp\left(\frac{-t^2}{2(v(Y)+Lt/3)}\right)$.

This implies for all $t > 0$: $\mathbb{P}(\|Y\|_{op} \geq t) \leq 2n \exp\left(\frac{-t^2}{2(v(Y)+Lt/3)}\right)$.

Corollary

For W s.t. W_{ij} independent up to symmetry, $\mathbb{E}W_{ij} = 0$, $|W_{ij}| \leq 1$, $\mathbb{E}W_{ij}^2 = O(d/n)$:

If $d \gg \log(n)$ then with high probability $\rho(W) = o(d)$.

Stronger bounds on $\rho = \rho(W)$

Theorem (Feige and Ofek, 2005)

Let $A \in \mathbb{R}^{n \times n}$: symmetric matrix with entries independent up to symmetry, $A_{ij} \in [0, 1]$, and such that $\mathbb{E}(A_{ij}) \leq d/n$, where $d \leq n^{1/5}$.

Then for some (universal) constant $\kappa > 0$, with high probability $\rho(A - \mathbb{E}(A)) \leq \kappa \sqrt{\max(d, \log(n))}$.

Corollary

For $d \gg \sqrt{\max(d, \log(n))}$, i.e. $d \gg \sqrt{\log(n)}$, with high probability $\rho(A - \mathbb{E}A) = o(d)$

Proof of Bernstein matrix inequality

Lemma

For independent Hermitian matrices X_k , $k \in [m]$, and

$$Y = \sum_{k \in [m]} X_k:$$

$$\mathbb{E} \text{Tr} e^{\theta Y} \leq \text{Tr} \exp \left(\sum_{k \in [m]} \ln \mathbb{E} e^{\theta X_k} \right)$$

Lemma

For Hermitian X such that $\mathbb{E}(X) = 0$ and $\|X\| \leq L$ almost surely, then:

$$\forall \theta \in (0, 3/L), \begin{cases} \mathbb{E} e^{\theta X} \preceq \exp \left(\frac{\theta^2/2}{1-\theta L/3} \mathbb{E} X^2 \right), \\ \ln \mathbb{E} e^{\theta X} \preceq \frac{\theta^2/2}{1-\theta L/3} \mathbb{E} X^2, \end{cases}$$

Lemma

For Hermitian A, B , if $A \preceq B$, then $\forall i \in [n], \lambda_i(A) \leq \lambda_i(B)$. Hence for all non-decreasing $f : \mathbb{R} \rightarrow \mathbb{R}$,

$$\text{Tr} f(A) \leq \text{Tr} f(B).$$

Application: Community Detection in the Stochastic Block Model

$\mathcal{G}(n, \{\alpha_i\}_{i \in [K]}, P)$, where $\alpha_i > 0, \sum_{i \in [K]} \alpha_i = 1, P \in [0, 1]^{K \times K}$:
multi-type version of the Erdős-Rényi random graph

- ▶ n vertices partitioned into K communities
- ▶ Type (community) of node $i : \sigma_i \in [K], \sigma_i : \text{i.i.d.}, \sim \alpha$
- ▶ Conditionally on $\sigma_{[n]}$, independently for each pair $i, j \in [n]$:
edge (i, j) present with probability $P_{\sigma(i), \sigma(j)}$.

Strong signal regime: fixed $K, \alpha, B \in \mathbb{R}_+^{K \times K}; P = (d/n)B$, with
 $\lim_{n \rightarrow +\infty} d = +\infty$

Spectral embedding

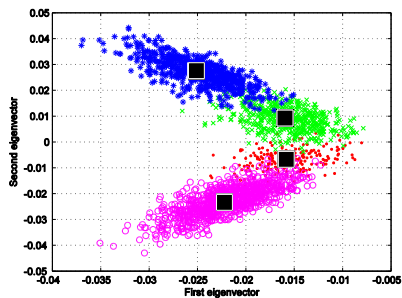
- ▶ Extract top two (more generally top R) eigenvalues λ_1, λ_2 of graph's adjacency matrix $A \in \mathbb{R}^{n \times n}$ (ordered by absolute value: $|\lambda_1| \geq |\lambda_2| \geq \dots$)
- ▶ Let $x_1, x_2 \in \mathbb{R}^n$: corresponding normalized eigenvectors
- ▶ Embed vertex $k \in [n]$ into \mathbb{R}^2 by letting $z_k := \sqrt{n}(x_1(k), x_2(k))$

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→ based on PCA dimensionality reduction of A to dimension R

Example: spectral embedding for SBM



A case with $K = 4$ communities

Spectral embedding seems to reflect community structure

→ Why / when do spectral methods work?

Theorem

Assume communities are **distinguishable**, i.e. for each $k \neq \ell \in [K]$, there exists $m \in [K]$ such that $B_{km} \neq B_{\ell m}$.

Assume $\sqrt{\ln(n)} \ll d \ll n^\delta$ for some fixed $\delta \in]0, 1[$. Let R : rank of matrix B . Then with high probability:

(i) the spectrum of A consists of R eigenvalues of order $\Theta(d)$ and $n - R$ eigenvalues of order $o(d)$.

(ii) R -dimensional spectral embedding reveals underlying communities: except for vanishing fraction of nodes $i \in [n]$,

$$\|z_i - z_j\| = \begin{cases} o(1) & \text{if } \sigma(i) = \sigma(j), \\ \Omega(1) & \text{if } \sigma(i) \neq \sigma(j) \end{cases}$$

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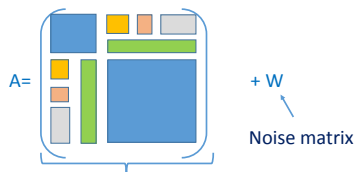
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Corollary

Under these conditions any sensible clustering scheme (eg K -means properly initialized) correctly classifies all but vanishing fraction of nodes.

Proof strategy



- ▶ Write adjacency matrix as $A = \bar{A} + W$ with $\bar{A}_{ij} = \frac{d}{n} B_{\sigma(i), \sigma(j)}$
- ▶ R leading eigen-elements of \bar{A} capture community structure
- ▶ Control perturbation of eigen-elements of a symmetric matrix \bar{A} by addition of symmetric matrix W in terms of **spectral radius** $\rho(W)$ of noise matrix
- ▶ Prove bound on $\rho(W)$ for random noise matrix W

Eigenstructure of \bar{A}

Block structure of $\bar{A} \Rightarrow \bar{A}x$ constant on each block \Rightarrow eigenvectors associated to non-zero eigenvalue are block-constant.

For $t \in \mathbb{R}^K$ define $x := \phi(t) = (t_{\sigma(i)})_{i \in [n]} \in \mathbb{R}^n$.

Then $\bar{A}\phi(t) = d\phi(Mt)$, where $M_{uv} := B_{uv}\alpha_v$.

Lemma

Spectrum of \bar{A} :

R eigen-pairs $(\lambda_u = d\mu_u, \bar{x}_u = \phi(t_u))$ where (μ_u, t_u) : eigen-pairs of M with $\mu_u \neq 0$;

0 : eigenvalue with multiplicity $n - R$

Eigenstructure of \bar{A} (continued)

Lemma

Under distinguishability hypothesis there exists $\epsilon > 0$ function of B, α such that for any choice of normalized leading eigenvectors $\bar{x}_1, \dots, \bar{x}_R$, $\bar{z}_i = \sqrt{n}(\bar{x}_1(i), \dots, \bar{x}_R(i))^T$ verify

$$\sigma(i) \neq \sigma(j) \Rightarrow \|\bar{z}_i - \bar{z}_j\| \geq \epsilon > 0$$

Proof: Let $t_u \in \mathbb{R}^k$ be such that $\sqrt{n}\bar{x}_u = \phi(t_u)$, and $\sqrt{\alpha} = \text{Diag}(\sqrt{\alpha_u})$.

Then: $\{\sqrt{\alpha}t_u\}_{u \in [R]}$: orthonormal family by orthonormality of the \bar{x}_u .

t_u eigenvectors of matrix $M = B\alpha$, hence $\sqrt{\alpha}t_u$: orthonormal family of eigenvectors of matrix $\sqrt{\alpha}B\sqrt{\alpha}$.

Thus $\sqrt{\alpha}B\sqrt{\alpha} = \sum_{u \in [R]} \mu_u(\sqrt{\alpha}t_u)(\sqrt{\alpha}t_u)^T$.

Equivalently: $B = \sum_{u \in [R]} \mu_u t_u t_u^T$.

Hence minimum of $\|\bar{z}_i - \bar{z}_j\|$ over $\sigma(i) \neq \sigma(j)$ strictly positive, for otherwise B has two identical rows, i.e. distinguishability fails.

Proof

- ▶ Matrix \bar{A} of rank R , spectral gaps $|\bar{\lambda}_i - \bar{\lambda}_j| = \Omega(d)$, R -dimensional spectral embedding with $\bar{x}_1, \dots, \bar{x}_R$ separates clusters $V_k = \{i \in [n] : \sigma_i = k\}$
- ▶ Assuming $\rho = \rho(A - \bar{A}) \ll d$, Weyl's inequality: R eigenvalues λ_i close to $\bar{\lambda}_i = \Omega(d)$, others of order $\rho \ll d$
- ▶ Associated eigenvectors x_i such that $\langle x_i, \bar{x}_i \rangle = 1 - O((\rho/d)^2)$

Then $\sum_{i \in [n]} \|z_i - \bar{z}_i\|^2 = n \sum_{u \in [R]} \|x_u - \bar{x}_u\|^2 = n\theta$ with $\theta = O((\rho/d)^2) = o(1)$

Hence (Tchebitchev inequality):

$$|\{i : \|z_i - \bar{z}_i\| \geq \theta^{1/3}\}| \leq n\theta^{1/3} = o(n)$$

Yields desired conclusion: except for vanishing fraction $\theta^{1/3}$ of nodes, spectral representatives z_i $\theta^{1/3}$ -close of corresponding \bar{z}_i , themselves clustered according to community structure

Feige-Ofek: $d \gg \sqrt{\log n} \Rightarrow \rho(A - \bar{A}) \ll d$