Abstract Interpretation
Semantics and applications to verification

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May 5th, 2017
Program of this lecture

Studied so far:

- **semantics**: behaviors of programs
- **properties**: safety, liveness, security...
- **approaches to verification**: typing, use of proof assistants, model checking

Today’s lecture: introduction to abstract interpretation

a **general framework for comparing semantics** introduced by Patrick Cousot and Radhia Cousot (1977)

- **abstraction**: use of a lattice of predicates
- **computing abstract over-approximations**, while preserving soundness
- **computing abstract over-approximations for loops**
Outline

1 Abstraction
   - Notion of abstraction
   - Abstraction and concretization functions
   - Galois connections

2 Abstract interpretation

3 Applications of abstract interpretation

4 A basic static analysis

5 A more realistic static analysis

6 Termination of the Static Analysis

7 Conclusion
Abstraction example 1: signs

**Abstraction: defined by a family of properties to use in proofs**

**Example:**
- objects under study: sets of mathematical integers
- abstract elements: signs

**Lattice of signs**

- $\bot$ denotes only $\emptyset$
- $\pm$ denotes any set of positive integers
- $0$ denotes any subset of $\{0\}$
- $-$ denotes any set of negative integers
- $\top$ denotes any set of integers

**Note:** the order in the abstract lattice corresponds to inclusion...
Abstraction example 1: signs

Definition: abstraction relation

- **concrete elements**: elements of the original lattice \((c \in \mathcal{P}(\mathbb{Z}))\)
- **abstract elements**: predicate \((a: \"\cdot \in \{\pm, 0, \ldots\}\)\)
- **abstraction relation**: \(c \vdash_S a\) when \(a\) describes \(c\)

Examples:

- \(\{1, 2, 3, 5, 7, 11, 13, 17, 19, 23, \ldots\} \vdash_S \pm\)
- \(\{1, 2, 3, 5, 7, 11, 13, 17, 19, 23, \ldots\} \vdash_S \top\)

We use abstract elements **to reason about operations**:

- if \(c_0 \vdash_S \pm\) and \(c_1 \vdash_S \pm\), then \(\{x_0 + x_1 | x_i \in c_i\} \vdash_S \pm\)
- if \(c_0 \vdash_S \pm\) and \(c_1 \vdash_S \pm\), then \(\{x_0 \cdot x_1 | x_i \in c_i\} \vdash_S \pm\)
- if \(c_0 \vdash_S \pm\) and \(c_1 \vdash_S \bot\), then \(\{x_0 \cdot x_1 | x_i \in c_i\} \vdash_S \bot\)
Abstraction

Abstraction example 1: signs

We can also consider the union operation:
- if \( c_0 \vdash_S \pm \) and \( c_1 \vdash_S \pm \), then \( c_0 \cup c_1 \vdash_S \pm \)
- if \( c_0 \vdash_S \pm \) and \( c_1 \vdash_S \bot \), then \( c_0 \cup c_1 \vdash_S \pm \)

But, what can we say about \( c_0 \cup c_1 \), when \( c_0 \vdash_S 0 \) and \( c_1 \vdash_S \pm \)?
- clearly, \( c_0 \cup c_1 \vdash_S \top \)...
- but no other relation holds
- in the abstract, we do not rule out negative values

We can extend the initial lattice:
- \( \geq 0 \) denotes any set of positive or null integers
- \( \leq 0 \) denotes any set of negative or null integers
- \( \neq 0 \) denotes any set of non null integers
- if \( c_0 \vdash_S \pm \) and \( c_1 \vdash_S 0 \), then \( c_0 \cup c_1 \vdash_S \geq 0 \)
Abstraction example 2: constants

Definition: abstraction based on constants

- **Concrete elements:** $\mathcal{P}(\mathbb{Z})$
- **Abstract elements:** $\perp, \top, n$ where $n \in \mathbb{Z}$
  
  \[ D_C^\# = \{\perp, \top\} \cup \{ n \mid n \in \mathbb{Z}\} \]
- **Abstraction relation:** $c \vdash_c n \iff c \subseteq \{n\}$

We obtain a flat lattice:

```
    ┌─┐   ┌─┐   ┌─┐   ┌─┐   ┌─┐
  ┌─┘  ┌─┘  ┌─┘  ┌─┘  ┌─┘  ┌─┘
 ... [−2] [−1] [0] [1] [2] ...
  ∏  ∏  ∏  ∏  ∏  ∏  ∏
```

Abstract reasoning:

- If $c_0 \vdash_c n_0$ and $c_1 \vdash_c n_1$, then $\{k_0 + k_1 \mid k_i \in c_i\} \vdash_c n_0 + n_1$
Abstraction example 3: Parikh vector

Definition: Parikh vector abstraction

- **Concrete elements:** \( P(\mathcal{A}^*) \) (sets of words over alphabet \( \mathcal{A} \))
- **Abstract elements:** \( \{\bot, \top\} \cup (\mathcal{A} \rightarrow \mathbb{N}) \)
- **Abstraction relation:** \( c \vdash \phi : \mathcal{A} \rightarrow \mathbb{N} \) if and only if:

\[
\forall w \in c, \forall a \in \mathcal{A}, \ a \text{ appears } \phi(a) \text{ times in } w
\]

Abstract reasoning:

- Concatenation:

  if \( \phi_0, \phi_1 : \mathcal{A} \rightarrow \mathbb{N} \) and \( c_0, c_1 \) are such that \( c_i \vdash \phi_i \),

\[
\{ w_0 \cdot w_1 \mid w_i \in c_i \} \vdash \phi_0 + \phi_1
\]

Information preserved, information deleted:

- Very precise information about the number of occurrences
- The order of letters is totally abstracted away (lost)
Abstraction example 4: interval abstraction

Definition: abstraction based on intervals

- **concrete elements:** $\mathcal{P}(\mathbb{Z})$
- **abstract elements:** $\perp, \top, (a, b)$ where $a \in \{-\infty\} \cup \mathbb{Z}$, $b \in \mathbb{Z} \cup \{+\infty\}$ and $a \leq b$
- **abstraction relation:**

\[
\emptyset \vdash_I \perp \\
S \vdash_I \top \\
S \vdash_I (a, b) \iff \forall x \in S, \ a \leq x \leq b
\]

Operations: TD
Abstraction example 5: non relational abstraction

**Definition: non relational abstraction**

- **concrete elements**: $\mathcal{P}(X \to Y)$, inclusion ordering
- **abstract elements**: $X \to \mathcal{P}(Y)$, pointwise inclusion ordering
- **abstraction relation**: $c \vdash_{\mathcal{N}_{\mathcal{R}}} a \iff \forall \phi \in c, \forall x \in X, \phi(x) \in a(x)$

**Information preserved, information deleted:**

- **very precise** information about the image of the functions in $c$
- **relations** such as (for given $x_0, x_1 \in X, y_0, y_1 \in Y$) the following are lost:
  
  \[
  \forall \phi \in c, \phi(x_0) = \phi(x_1) \\
  \forall \phi \in c, \forall x, x' \in X, \phi(x) \neq y_0 \lor \phi(x') \neq y_1
  \]
Notion of abstraction relation

**Concrete order:** so far, always inclusion
- the tighter the concrete set, the fewer behaviors
- smaller concrete sets correspond to more precise properties

**Abstraction relation:** \( c \vdash a \) when \( c \) satisfies \( a \)
- if \( c_0 \subseteq c_1 \) and \( c_1 \) satisfies \( a \), in all our examples, \( c_0 \) also satisfies \( a \)

**Abstract order:** in all our examples,
- it matches the abstraction relation as well:
  - if \( a_0 \sqsubseteq a_1 \) and \( c \) satisfies \( a_0 \), then \( c \) also satisfies \( a_1 \)
- great advantage: we can reason about implication in the abstract, without looking back at the concrete properties

We will now formalize this in detail...
Outline

1 Abstraction
   - Notion of abstraction
   - Abstraction and concretization functions
   - Galois connections

2 Abstract interpretation

3 Applications of abstract interpretation

4 A basic static analysis

5 A more realistic static analysis

6 Termination of the Static Analysis

7 Conclusion
Towards adjoint functions

We consider a **concrete lattice** \((C, \subseteq)\) and an **abstract lattice** \((A, \sqsubseteq)\).

So far, we used **abstraction relations**, that are consistent with orderings:

### Abstraction relation compatibility

- \(\forall c_0, c_1 \in C, \forall a \in A, c_0 \subseteq c_1 \land c_1 \vdash a \implies c_0 \vdash a\)
- \(\forall c \in C, \forall a_0, a_1 \in A, c \vdash a_0 \land a_0 \sqsubseteq a_1 \implies c \vdash a_1\)

When we have a \(c\) (resp., \(a\)) and try to map it into a compatible \(a\) (resp. a \(c\)), **the abstraction relation is not a convenient tool**.

Hence, we shall use **adjoint functions** between \(C\) and \(A\).

- from concrete to abstract: **abstraction**
- from abstract to concrete: **concretization**
Concretization function

Our first adjoint function:

**Definition: concretization function**

Concretization function $\gamma : A \rightarrow C$ (if it exists) maps abstract $a$ into the weakest (i.e., most general) concrete $c$ that satisfies $a$ (i.e., $c \vdash a$).

Note: in common cases, there exists a $\gamma$.

- $c \vdash a$ if and only if $c \subseteq \gamma(a)$
Concretization function: a few examples

Signs abstraction:

\[ \gamma_S : \begin{array}{c|c}
\top & \mathbb{Z} \\
\pm & \mathbb{Z}^* \\
0 & \{0\} \\
\equiv & \mathbb{Z}^* \\
\bot & \emptyset
\end{array} \]

Constants abstraction:

\[ \gamma_C : \begin{array}{c|c}
\top & \mathbb{Z} \\
n & \{n\} \\
\bot & \emptyset
\end{array} \]

Non relational abstraction:

\[ \gamma_{NR} : \begin{array}{c|c}
(X \rightarrow \mathcal{P}(Y)) & \mathcal{P}(X \rightarrow Y) \\
\Phi & \{\phi : X \rightarrow Y \mid \forall x \in X, \phi(x) \in \Phi(x)\}
\end{array} \]

Parikh vector abstraction: exercise!
Abstraction function

Our second adjoint function:

Definition: abstraction function

**Abstraction function** $\alpha : C \to A$ (if it exists) maps concrete $c$ into the most precise abstract $a$ that soundly describes $c$ (i.e., $c \vdash a$).

Note: in quite a few cases (including some in this course), there is no $\alpha$.

Summary on adjoint functions:

- $\alpha$ returns the **most precise abstract predicate** that holds true for its argument
  this is called the **best abstraction**
- $\gamma$ returns the **most general concrete meaning** of its argument
  hence, is called the **concretization**
Abstraction: a few examples

**Constants abstraction:**

\[ \alpha_C : (c \subseteq \mathbb{Z}) \mapsto \begin{cases} \bot & \text{if } c = \emptyset \\ n & \text{if } c = \{n\} \\ \top & \text{otherwise} \end{cases} \]

**Non relational abstraction:**

\[ \alpha_{NR} : P(X \rightarrow Y) \mapsto X \rightarrow P(Y) \]

\[ c \mapsto (x \in X) \mapsto \{\phi(x) \mid \phi \in c\} \]

**Signs abstraction** and **Parikh vector abstraction**: exercises
Outline

1. Abstraction
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   - Galois connections

2. Abstract interpretation

3. Applications of abstract interpretation

4. A basic static analysis

5. A more realistic static analysis

6. Termination of the Static Analysis

7. Conclusion
Abstraction  Galois connections

Definition

So far, we have:

- **abstraction** $\alpha : C \rightarrow A$
- **concretization** $\gamma : A \rightarrow C$

How to tie them together?

**They should agree on a same abstraction relation $\vdash$ !**

**Definition: Galois connection**

A **Galois connection** is defined by a **concrete lattice** $(C, \subseteq)$, an **abstract lattice** $(A, \sqsubseteq)$, an **abstraction function** $\alpha : C \rightarrow A$ and a **concretization function** $\gamma : A \rightarrow C$ such that:

$$\forall c \in C, \forall a \in A, \alpha(c) \sqsubseteq a \iff c \subseteq \gamma(a) \quad (\iff c \vdash a)$$

**Notation:**

$$(C, \subseteq) \xleftarrow{\alpha} \xrightarrow{\gamma} (A, \sqsubseteq)$$

Note: in practice, we shall rarely use $\vdash$; we use $\alpha, \gamma$ instead
Example: constants abstraction and Galois connection

**Constants lattice** \( D_C^\# = \{ \bot, \top \} \uplus \{ n \mid n \in \mathbb{Z} \} \)

\[
\begin{align*}
\alpha_C(c) & = \bot \quad \text{if } c = \emptyset & \gamma_C(\top) & \mapsto \mathbb{Z} \\
\alpha_C(c) & = n \quad \text{if } c = \{ n \} & \gamma_C(n) & \mapsto \{ n \} \\
\alpha_C(c) & = \top \quad \text{otherwise} & \gamma_C(\bot) & \mapsto \emptyset
\end{align*}
\]

Thus:

- if \( c = \emptyset \), \( \forall a, c \subseteq \gamma_C(a) \), i.e., \( c \subseteq \gamma_C(a) \iff \alpha_C(c) = \bot \subseteq a \)
- if \( c = \{ n \} \),
  \[\alpha_C(\{ n \}) = n \subseteq c \iff c = n \lor c = \top \iff c = \{ n \} \subseteq \gamma_C(a)\]
- if \( c \) has at least two distinct elements \( n_0, n_1 \), \( \alpha_C(c) = \top \) and \( c \subseteq \gamma_C(a) \Rightarrow a = \top \), i.e., \( c \subseteq \gamma_C(a) \iff \alpha_C(c) = \bot \subseteq a \)

**Constant abstraction: Galois connection**

\( c \subseteq \gamma_C(a) \iff \alpha_C(c) \subseteq a \), therefore,

\[
(\mathcal{P}(\mathbb{Z}), \subseteq) \xrightarrow{\alpha_C} (D_C^\#, \sqsubseteq)
\]
Example: non relational abstraction Galois connection

We have defined:

\[ \alpha_{\mathcal{NR}} : (c \subseteq (X \to Y)) \mapsto (x \in X) \mapsto \{ f(x) \mid f \in c \} \]
\[ \gamma_{\mathcal{NR}} : (\Phi \in (X \to \mathcal{P}(Y))) \mapsto \{ f : X \to Y \mid \forall x \in X, f(x) \in \Phi(x) \} \]

Let \( c \in \mathcal{P}(X \to Y) \) and \( \Phi \in (X \to \mathcal{P}(Y)) \); then:

\[ \alpha_{\mathcal{NR}}(c) \subseteq \Phi \iff \forall x \in X, \alpha_{\mathcal{NR}}(c)(x) \subseteq \Phi(x) \]
\[ \iff \forall x \in X, \{ f(x) \mid f \in c \} \subseteq \Phi(x) \]
\[ \iff \forall f \in c, \forall x \in X, f(x) \in \Phi(x) \]
\[ \iff \forall f \in c, f \in \gamma_{\mathcal{NR}}(\Phi) \]
\[ \iff c \subseteq \gamma_{\mathcal{NR}}(\Phi) \]

Non relational abstraction: Galois connection

\[ c \subseteq \gamma_{\mathcal{NR}}(a) \iff \alpha_{\mathcal{NR}}(c) \subseteq a, \text{ therefore,} \]

\[ (\mathcal{P}(X \to Y), \subseteq) \xleftarrow{\alpha_{\mathcal{NR}}} (X \to \mathcal{P}(Y), \subseteq) \xrightarrow{\gamma_{\mathcal{NR}}} (\mathcal{P}(X \to Y), \subseteq) \]
Galois connection properties

Galois connections have many useful properties.

In the next few slides, we consider a Galois connection $(C, \subseteq) \xrightarrow{\alpha} (A, \subseteq)$ and establish a few interesting properties.

**Extensivity, contractivity**

- $\alpha \circ \gamma$ is **contractive**: $\forall a \in A, \alpha \circ \gamma(a) \subseteq a$
- $\gamma \circ \alpha$ is **extensive**: $\forall c \in C, c \subseteq \gamma \circ \alpha(c)$

**Proof:**

- let $a \in A$; then, $\gamma(a) \subseteq \gamma(a)$, thus $\alpha(\gamma(a)) \subseteq a$
- let $c \in C$; then, $\alpha(c) \subseteq \alpha(c)$, thus $c \subseteq \gamma(\alpha(a))$
Galois connection properties

Monotonicity of adjoints

- α is monotone
- γ is monotone

Proof:

- monotonicity of α: let \( c_0, c_1 \in C \) such that \( c_0 \subseteq c_1 \);
  by extensivity of \( \gamma \circ \alpha \), \( c_1 \subseteq \gamma(\alpha(c_1)) \), so by transitivity, \( c_0 \subseteq \gamma(\alpha(c_1)) \)
  by definition of the Galois connection, \( \alpha(c_0) \subseteq \alpha(c_1) \)

- monotonicity of γ: same principle

Note: many proofs can be derived by duality

Duality principle applied for Galois connections

If \( (C, \subseteq) \xleftrightarrow[\alpha]{\gamma} (A, \sqsubseteq) \), then \( (A, \sqsupseteq) \xleftrightarrow[\gamma]{\alpha} (C, \supseteq) \)
Galois connection properties

**Iteration of adjoints**

- $\alpha \circ \gamma \circ \alpha = \alpha$
- $\gamma \circ \alpha \circ \gamma = \gamma$
- $\alpha \circ \gamma$ (resp., $\gamma \circ \alpha$) is idempotent, hence a lower (resp., upper) closure operator

**Proof:**

- $\alpha \circ \gamma \circ \alpha = \alpha$:  
  let $c \in C$, then $\gamma \circ \alpha(c) \subseteq \gamma \circ \alpha(c)$  
  hence, by the Galois connection property, $\alpha \circ \gamma \circ \alpha(c) \sqsubseteq \alpha(c)$  
  moreover, $\gamma \circ \alpha$ is extensive and $\alpha$ monotone, so $\alpha(c) \sqsubseteq \alpha \circ \gamma \circ \alpha(c)$  
  thus, $\alpha \circ \gamma \circ \alpha(c) = \alpha(c)$
- the second point can be proved similarly (duality); the others follow
Galois connection properties

Properties on iterations of adjoint functions:

- Concrete domain
- Abstract domain

Diagram showing the relationships between concrete and abstract domains with arrows indicating the properties of adjoint functions.
Galois connection properties

\( \alpha \) preserves least upper bounds

\[ \forall c_0, c_1 \in C, \ \alpha(c_0 \cup c_1) = \alpha(c_0) \sqcup \alpha(c_1) \]

By duality:

\[ \forall a_0, a_1 \in A, \ \gamma(c_0 \cap c_1) = \gamma(c_0) \sqcap \gamma(c_1) \]

Proof:
First, we observe that \( \alpha(c_0) \sqcup \alpha(c_1) \subseteq \alpha(c_0 \cup c_1) \), i.e. \( \alpha(c_0 \cup c_1) \) is an upper bound of \( \{ \alpha(c_0), \alpha(c_1) \} \).

We now prove it is the least upper bound. For all \( a \in A \):

\[ \alpha(c_0 \cup c_1) \subseteq a \iff c_0 \cup c_1 \subseteq \gamma(a) \]
\[ \iff c_0 \subseteq \gamma(a) \land c_1 \subseteq \gamma(a) \]
\[ \iff \alpha(c_0) \subseteq a \land \alpha(c_1) \subseteq a \]
\[ \iff \alpha(c_0) \sqcup \alpha(c_1) \subseteq a \]

Note: when \( C, A \) are complete lattices, this extends to families of elements
Galois connection properties

**Uniqueness of adjoints**

- given $\gamma : C \to A$, there exists at most one $\alpha : A \to C$ such that $(C, \subseteq) \xleftarrow{\gamma} (A, \sqsubseteq)$, and, if it exists, $\alpha(c) = \bigcap\{a \in A \mid c \subseteq \gamma(a)\}$
- similarly, given $\alpha : A \to C$, there exists at most one $\gamma : C \to A$ such that $(C, \subseteq) \xrightarrow{\gamma} (A, \sqsubseteq)$, and it is defined dually

**Proof of the first point** (the other follows by duality):
we assume that there exists an $\alpha$ so that we have a Galois connection and prove that, $\alpha(c) = \bigcap\{a \in A \mid c \subseteq \gamma(a)\}$ for a given $c \in C$.

- if $a \in A$ is such that $c \subseteq \gamma(a)$, then $\alpha(a) \subseteq c$ thus, $\alpha(a)$ is a lower bound of $\{a \in A \mid c \subseteq \gamma(a)\}$.
- let $a_0 \in A$ be a lower bound of $\{a \in A \mid c \subseteq \gamma(a)\}$.
  since $\gamma \circ \alpha$ is extensive, $c \subseteq \gamma(\alpha(c))$ and $\alpha(c) \in \{a \in A \mid c \subseteq \gamma(a)\}$.
  hence, $a_0 \subseteq \alpha(c)$

Thus, $\alpha(c)$ is the least upper bound of $\{a \in A \mid c \subseteq \gamma(a)\}$
Construction of adjoint functions

The adjoint uniqueness property is actually a very strong property:

- it allows to construct an abstraction from a concretization
- ... or to understand why no abstraction can be constructed :-)

Turning an adjoint into a Galois connection (1)

Let \((C, \subseteq)\) and \((A, \sqsubseteq)\) be two lattices, such that any subset of \(A\) has a greatest lower bound and let \(\gamma : (A, \sqsubseteq) \to (C, \subseteq)\) be a monotone function.

Then, the function below defines a Galois connection:

\[
\alpha(c) = \bigcap\{a \in A \mid c \subseteq \gamma(a)\}
\]

Example of abstraction with no \(\alpha\): when \(\bigcap\) is not defined on all families, e.g., lattice of convex polyhedra, abstracting sets of points in \(\mathbb{R}^2\).

Exercise: state the dual property and apply the same principle to the concretization
Galois connection characterization

A characterization of Galois connections

Let \((C, \subseteq)\) and \((A, \sqsubseteq)\) be two lattices, and \(\alpha : C \rightarrow A\) and \(\gamma : A \rightarrow C\) be two monotone functions, such that:

- \(\alpha \circ \gamma\) is contractive
- \(\gamma \circ \alpha\) is extensive

Then, we have a Galois connection

\[
(C, \subseteq) \overset{\gamma}{\leftarrow} \overset{\alpha}{\rightarrow} (A, \sqsubseteq)
\]

Proof:

- let \(c \in C\) and \(a \in A\) such that \(\alpha(c) \sqsubseteq a\).
  - then: \(\gamma(\alpha(c)) \sqsubseteq \gamma(a)\) (as \(\gamma\) is monotone)
  - \(c \subseteq \gamma(\alpha(c))\) (as \(\gamma \circ \alpha\) is extensive)
  - thus, \(c \subseteq \gamma(a)\), by transitivity
- the other implication can be proved by duality
Outline

1 Abstraction

2 Abstract interpretation
   - Abstract computation
   - Fixpoint transfer

3 Applications of abstract interpretation

4 A basic static analysis

5 A more realistic static analysis

6 Termination of the Static Analysis

7 Conclusion
Constructing a static analysis

We have set up a notion of **abstraction**:
- it describes **sound** approximations of **concrete properties** with **abstract predicates**
- there are several ways to formalize it (abstraction, concretization...)
- we now wish to **compute sound abstract predicates**

In the following, we assume
- a **Galois connection**

\[
(C, \subseteq) \xleftarrow{\gamma} (A, \sqsubseteq)
\]

- a **concrete semantics \([\cdot]\)**, with a **constructive definition**
  i.e., \([P]\) is defined by constructive equations (\([P] = f(\ldots)\)), least fixpoint formula (\([P] = \text{lfp}_\emptyset f\)...)
Abstract transformer

A fixed concrete element $c_0$ can be abstracted by $\alpha(c_0)$.

We now consider a monotone concrete function $f : C \to C$

- given $c \in C$, $\alpha \circ f(c)$ abstracts the image of $c$ by $f$
- if $c \in C$ is abstracted by $a \in A$, then $f(c)$ is abstracted by $\alpha \circ f \circ \gamma(a)$:
  
  $c \subseteq \gamma(a)$ by assumption
  $f(c) \subseteq f(\gamma(a))$ by monotonicity of $f$
  $\alpha(f(c)) \subseteq \alpha(f(\gamma(a)))$ by monotonicity of $\alpha$

Definition: best and sound abstract transformers

- the best abstract transformer approximating $f$ is $f^\# = \alpha \circ f \circ \gamma$
- a sound abstract transformer approximating $f$ is any operator $f^\# : A \to A$, such that $\alpha \circ f \circ \gamma \sqsubseteq f^\#$ (or equivalently, $f \circ \gamma \sqsubseteq \gamma \circ f^\#$)
Example: lattice of signs

- $f : D^\#_C \to D^\#_C$, $c \mapsto \{-n \mid n \in c\}$
- $f^\# = \alpha \circ f \circ \gamma$

**Lattice of signs:**

```
⊥  ⊤
|   |
|   |
|   |
⊥  ⊤
```

**Abstract negation operator:**

<table>
<thead>
<tr>
<th>$a$</th>
<th>$\ominus^#(a)$</th>
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<tbody>
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<td>⊥</td>
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- here, the best abstract transformer is very easy to compute
- no need to use an approximate one
Abstract computation

Abstract \( n \)-ary operators

We can generalize this to \( n \)-ary operators, such as boolean operators and arithmetic operators.

\[ g : C^n \rightarrow C \]

Definition: sound and exact abstract operators

Let \( g : C^n \rightarrow C \) be an \( n \)-ary operator, monotone in each component. Then:

- the **best abstract operator** approximating \( g \) is defined by:
  \[
g^\# : A^n \rightarrow A \quad (a_0, \ldots, a_{n-1}) \mapsto \alpha \circ g(\gamma(a_0), \ldots, \gamma(a_{n-1})) \]

- a **sound abstract transformer** approximating \( g \) is any operator \( g^\# : A^n \rightarrow A \), such that
  \[
  \forall (a_0, \ldots, a_{n-1}) \in A^n, \alpha \circ g(\gamma(a_0), \ldots, \gamma(a_{n-1})) \sqsubseteq g^\#(a_0, \ldots, a_{n-1})
  \]
  (i.e., equivalently, \( g(\gamma(a_0), \ldots, \gamma(a_{n-1})) \subseteq \gamma \circ g^\#(a_0, \ldots, a_{n-1}) \))
Example: lattice of signs arithmetic operators

Application:

- $\oplus : C^2 \rightarrow C, (c_0, c_1) \mapsto \{n_0 + n_1 \mid n_i \in c_i\}$
- $\otimes : C^2 \rightarrow C, (c_0, c_1) \mapsto \{n_0 \cdot n_1 \mid n_i \in c_i\}$

Best abstract operators:

<table>
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<th>$\oplus^#$</th>
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Example of loss in precision:

- $\{8\} \in \gamma_S(\pm)$ and $\{-2\} \in \gamma_S(-)$
- $\oplus^\#(\pm, -) = T$ is a lot worse than $\alpha_S(\oplus(\{8\}, \{-2\})) = \pm$
Example: lattice of signs set operators

Best abstract operators approximating $\cup$ and $\cap$:

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Example of loss in precision:

- $\gamma(-) \cup \gamma(\pm) = \{ n \in \mathbb{Z} \mid n \neq 0 \} \subset \gamma(T)$
Outline

1. Abstraction

2. Abstract interpretation
   - Abstract computation
   - Fixpoint transfer

3. Applications of abstract interpretation

4. A basic static analysis

5. A more realistic static analysis

6. Termination of the Static Analysis

7. Conclusion
Fixpoint transfer

What about **loops**? semantic functions defined by **fixpoints**?

**Theorem: exact fixpoint transfer**

We assume \((C, \subseteq)\) and \((A, \sqsubseteq)\) are complete lattices. We consider a Galois connection \((C, \subseteq) \xleftarrow{\gamma} (A, \sqsubseteq)\), two functions \(f : C \rightarrow C\) and \(f^\# : A \rightarrow A\) and two elements \(c_0 \in C, a_0 \in A\) such that:

- \(f\) is continuous
- \(f^\#\) is monotone
- \(\alpha \circ f = f^\# \circ \alpha\)
- \(\alpha(c_0) = a_0\)

Then:

- both \(f\) and \(f^\#\) have a least-fixpoint (by Tarski's fixpoint theorem)
- \(\alpha(\text{lfp}_{c_0} f) = \text{lfp}_{a_0} f^\#\)
Fixpoint transfer: proof

- \( \alpha(\text{lfp}_{c_0} f) \) is a fixpoint of \( f^\# \) since:
  
  \[
  f^\#(\alpha(\text{lfp}_{c_0} f)) = \alpha(f(\text{lfp}_{c_0} f)) = \alpha(\text{lfp}_{c_0} f)
  \]
  
  since \( \alpha \circ f = f^\# \circ \alpha \)

- To show that \( \alpha(\text{lfp}_{c_0} f) \) is the least-fixpoint of \( f^\# \),
  we assume that \( X \) is another fixpoint of \( f^\# \) greater than \( a_0 \) and we show that \( \alpha(\text{lfp}_{c_0} f) \subseteq X \), i.e., that \( \text{lfp}_{c_0} f \subseteq \gamma(X) \).

As \( \text{lfp}_{c_0} f = \bigcup_{n \in \mathbb{N}} f^n(c_0) \) (by Kleene’s fixpoint theorem), it amounts to proving that \( \forall n \in \mathbb{N}, f^n(c_0) \subseteq \gamma(X) \).

By induction over \( n \):

- \( f^0(c_0) = c_0 \), thus \( \alpha(f^0(c_0)) = a_0 \subseteq X \); thus, \( f^0(c_0) \subseteq \gamma(X) \).
- let us assume that \( f^n(c_0) \subseteq \gamma(X) \), and let us show that \( f^{n+1}(c_0) \subseteq \gamma(X) \), i.e. that \( \alpha(f^{n+1}(c_0)) \subseteq X \):

  \[
  \alpha(f^{n+1}(c_0)) = \alpha \circ f(f^n(c_0)) = f^\# \circ \alpha(f^n(c_0)) \subseteq f^\#(X) = X
  \]

as \( \alpha(f^n(c_0)) \subseteq X \) and \( f^\# \) is monotone.
Constructive analysis of loops

How to get a constructive fixpoint transfer theorem?

Theorem: fixpoint abstraction

Under the assumptions of the previous theorem, and with the following additional hypothesis:

- lattice $A$ is of finite height

We compute the sequence $(a_n)_{n \in \mathbb{N}}$ defined by $a_{n+1} = a_n \sqcup f^\sharp(a_n)$.

Then, $(a_n)_{n \in \mathbb{N}}$ converges and its limit $a_\infty$ is such that $\alpha(lfp_{c_0} f) = a_\infty$.

Proof: exercise.

Note:

- the assumptions we have made are too restrictive in practice
- more general fixpoint abstraction methods in the next lectures
Outline

1. Abstraction
2. Abstract interpretation
3. Applications of abstract interpretation
4. A basic static analysis
5. A more realistic static analysis
6. Termination of the Static Analysis
7. Conclusion
Comparing existing semantics

1. A **concrete semantics** $[P]$ is given: e.g., big steps operational semantics

2. An **abstract semantics** $[P]^\#$ is given: e.g., denotational semantics

3. **Search for an abstraction relation between them**
   e.g., $[P]^\# = \alpha([P])$, or $[P] \subseteq \gamma([P]^\#)$

**Examples:**
- finite traces semantics as an abstraction of bi-finitary trace semantics
- denotational semantics as an abstraction of trace semantics
- types as an abstraction of denotational semantics
- ...

**Payoff:**
- better understanding of ties across semantics
- chance to generalize existing definitions
Applications of abstract interpretation

Derivation of a static analysis

1. Start from a **concrete semantics** $\llbracket P \rrbracket$
2. **Choose an abstraction** defined by a Galois connection or a concretization function (usually)
3. **Derive an abstract semantics** $\llbracket P \rrbracket^\#$ such that $\llbracket P \rrbracket \subseteq \gamma(\llbracket P \rrbracket^\#)$

**Examples:**
- derivation of an analysis with a numerical lattice (constants, intervals...)
- construction of an analysis for a complex programming language

**Payoff:**
- the derivation of the abstract semantics is quite systematic
- this process offers good opportunities for a modular analysis design

There are many ways to apply abstract interpretation.
Outline

1. Abstraction
2. Abstract interpretation
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7. Conclusion
A very simple language and its semantics

We now apply this to a very simple language, and **derive a static analysis step by step**, from a concrete semantics and an abstraction.

- we assume a **fixed set of** $n$ **integer variables** $x_0, \ldots, x_{n-1}$
- we consider the language defined by the grammar below:

$$
P ::= x_i = n  \quad \text{where } n \in \mathbb{Z}
$$

$$
| x_i = x_j + x_k  \quad \text{basic, three-addresses arithmetics}
| x_i = x_j - x_k  \quad \text{basic, three-addresses arithmetics}
| x_i = x_j \cdot x_k  \quad \text{basic, three-addresses arithmetics}
| P; P  \quad \text{concatenation}
| \text{input}(x_i)  \quad \text{reading of a positive input}
| \text{if}(x_i > 0) P \text{ else } P
| \text{while}(x_i > 0) P
$$

- a state is a vector $\sigma = (\sigma_0, \ldots, \sigma_{n-1}) \in \mathbb{Z}^n$
- a single initial state $\sigma_{\text{init}} = (0, \ldots, 0)$
Concrete semantics

We let \([P] : \mathcal{P}(\mathbb{Z}^n) \rightarrow \mathcal{P}(\mathbb{Z}^n)\) be defined by:

\[
\begin{align*}
[x_i = n](S) &= \{\sigma[i \leftarrow n] | \sigma \in S\} \\
[x_i = x_j + x_k](S) &= \{\sigma[i \leftarrow \sigma_j + \sigma_k] | \sigma \in S\} \\
[x_i = x_j - x_k](S) &= \{\sigma[i \leftarrow \sigma_j - \sigma_k] | \sigma \in S\} \\
[x_i = x_j \cdot x_k](S) &= \{\sigma[i \leftarrow \sigma_j \cdot \sigma_k] | \sigma \in S\} \\
\text{input}(x_i)(S) &= \{\sigma[i \leftarrow n] | \sigma \in S \land n > 0\} \\
P_0; P_1(S) &= [P_1] \circ [P_0](S) \\
\text{if}(x_i > 0) P_0 \text{ else } P_1(S) &= [P_0](\{\sigma \in S | \sigma_i > 0\}) \cup [P_1](\{\sigma \in S | \sigma_i \leq 0\}) \\
\text{while}(x_i > 0) P(S) &= \{\sigma \in \text{lfp}_{S} f | \sigma_i \leq 0\} \text{ where } f : S' \mapsto S' \cup [P](\{\sigma \in S' | \sigma_i > 0\})
\end{align*}
\]

- given a complete program \(P\), the \textbf{reachable states} are defined by \([P](\{\sigma_{\text{init}}\})\)
Abstraction

We compose two abstractions:

- **non relational abstraction**: the values a variable may take is abstracted separately from the other variables
- **sign abstraction**: the set of values observed for each variable is abstracted into the lattice of signs

**Concrete domain**: \((\mathcal{P}(\mathbb{Z}^n), \subseteq)\)

**Abstract domain**: \((D^#, \subseteq), \text{ where } D^# = (D^#_S)^n\) and \(\subseteq\) is the pointwise ordering

**Galois connection** \((\mathcal{P}(\mathbb{Z}), \subseteq) \xRightarrow[\alpha]{\gamma} (D^#, \subseteq), \text{ defined by}\)

\[
\begin{align*}
\alpha &: S \mapsto (\alpha_S(\{\sigma_0 | \sigma \in S\}), \ldots, \alpha_S(\{\sigma_{n-1} | \sigma \in S\})) \\
\gamma &: S^# \mapsto \{\sigma \in \mathbb{Z}^n | \forall i, \sigma_i \in \gamma_S(S^#_i)\}
\end{align*}
\]
Example

Factorial function:

```plaintext
input(x_0);

x_1 = 1;
x_2 = 1;

while(x_0 > 0) {
    x_1 = x_0 \cdot x_1;
    x_0 = x_0 - x_2;
}
```

Abstraction of the semantics:

- abstract pre-condition: \((0, 0, 0)\)
- abstract state before the loop: \((\pm, \pm, \pm)\)
- abstract post-condition (after the loop): \((0, \pm, \pm)\)
Computation of the abstract semantics

We search for an abstract semantics $\llbracket P \rrbracket^\# : D^\# \to D^\#$ such that:

$$\alpha \circ \llbracket P \rrbracket = \llbracket P \rrbracket^\# \circ \alpha$$

We observe that:

$$\alpha(S) = (\alpha_S(\{\sigma_0 \mid \sigma \in S\}), \ldots, \alpha_S(\{\sigma_{n-1} \mid \sigma \in S\}))$$
$$\alpha \circ \llbracket P \rrbracket(S) = (\alpha_S(\{\sigma_0 \mid \sigma \in \llbracket P \rrbracket(S)\}), \ldots, \alpha_S(\{\sigma_{n-1} \mid \sigma \in \llbracket P \rrbracket(S)\}))$$

We start with $x_i = n$:

$$\alpha \circ \llbracket x_i = n \rrbracket(S)$$
$$= (\alpha_S(\{\sigma_0 \mid \sigma \in \{\sigma[i \leftarrow n] \mid \sigma \in S\}\}), \ldots, \alpha_S(\{\sigma_{n-1} \mid \sigma \in \{\sigma[i \leftarrow n] \mid \sigma \in S\}\}))$$
$$= (\alpha_S(\{\sigma_0 \mid \sigma \in S\}), \ldots, \alpha_S(\{\sigma_{n-1} \mid \sigma \in S\}))[i \leftarrow \alpha_S(\{n\})]$$
$$= \alpha(S)[i \leftarrow \alpha_S(\{n\})]$$
$$= \llbracket x_i = n \rrbracket^\#(\alpha(S))$$

where

$$\llbracket x_i = n \rrbracket^\#(S^\#) = S^\#[i \leftarrow \alpha_S(\{n\})]$$
Computation of the abstract semantics

Other assignments are treated in a similar manner:

\[
\begin{align*}
\llbracket x_i = x_j + x_k \rrbracket (S^\#) &= S^\#[i \leftarrow S_j^\# \oplus S_k^\#] \\
\llbracket x_i = x_j - x_k \rrbracket (S^\#) &= S^\#[i \leftarrow S_j^\# \ominus S_k^\#] \\
\llbracket x_i = x_j \cdot x_k \rrbracket (S^\#) &= S^\#[i \leftarrow S_j^\# \otimes S_k^\#] \\
\llbracket \text{input}(x_i) \rrbracket (S^\#) &= S^\#[i \leftarrow \pm]
\end{align*}
\]

Proofs are left as exercises
Computation of the abstract semantics

We now consider the case of tests:

\[
\alpha \circ \lbrack \text{if}(x_i > 0) \ P_0 \ \text{else} \ P_1 \rbrack(S)
\]

\[
= \alpha(\lbrack P_0 \rbrack(\{\sigma \in S \mid \sigma_i > 0\}) \cup \lbrack P_1 \rbrack(\{\sigma \in S \mid \sigma_i \leq 0\}))
\]

\[
= \alpha(\lbrack P_0 \rbrack(\{\sigma \in S \mid \sigma_i > 0\})) \cup \alpha(\lbrack P_1 \rbrack(\{\sigma \in S \mid \sigma_i \leq 0\}))
\]

as \(\alpha\) preserves least upper bounds

\[
= \lbrack P_0 \rbrack^\#(\alpha(\{\sigma \in S \mid \sigma_i > 0\})) \cup \lbrack P_1 \rbrack^\#(\alpha(\{\sigma \in S \mid \sigma_i \leq 0\}))
\]

\[
= \lbrack P_0 \rbrack^\#(\alpha(S) \cap \top[i \leftarrow \pm]) \cup \lbrack P_1 \rbrack^\#(\alpha(S))
\]

where

\[
\lbrack \text{if}(x_i > 0) \ P_0 \ \text{else} \ P_1 \rbrack^\#(S^\#) = \lbrack P_0 \rbrack^\#(S^\# \cap \top[i \leftarrow \pm]) \cup \lbrack P_1 \rbrack^\#(S^\#)
\]

In the case of loops:

\[
\lbrack \text{while}(x_i > 0) \ P \rbrack^\#(S^\#) = \text{lfp}_{S^\#} f^\#
\]

where

\[
f^\#: S^\# \mapsto S^\# \sqcup \lbrack P \rbrack^\#(S^\# \cap \top[i \leftarrow \pm])
\]

Proof: exercise
Abstract semantics

Abstract semantics and soundness

We have derived the following definition of $[P]^{\#}$:

\[
[x_i = n]^{\#}(S^{\#}) = S^{\#}[i \leftarrow \alpha_S(\{n\})]
\]
\[
[x_i = x_j + x_k]^{\#}(S^{\#}) = S^{\#}[i \leftarrow S_j^{\#} \oplus^{\#} S_k^{\#}]
\]
\[
[x_i = x_j - x_k]^{\#}(S^{\#}) = S^{\#}[i \leftarrow S_j^{\#} \ominus^{\#} S_k^{\#}]
\]
\[
[x_i = x_j \cdot x_k]^{\#}(S^{\#}) = S^{\#}[i \leftarrow S_j^{\#} \otimes^{\#} S_k^{\#}]
\]
\[
[\text{input}(x_i)]^{\#}(S^{\#}) = S^{\#}[i \leftarrow \, +]
\]
\[
[\text{if}(x_i > 0) P_0 \text{ else } P_1]^{\#}(S^{\#}) = [P_0]^{\#}(S^{\#} \sqcap [i \leftarrow \, +]) \cup [P_1]^{\#}(S^{\#})
\]
\[
[\text{while}(x_i > 0) P]^{\#}(S^{\#}) = \text{lfp}_{S^{\#}} f^{\#} \text{ where } f^{\#} : S^{\#} \mapsto S^{\#} \sqcup [P]^{\#}(S^{\#} \sqcap [i \leftarrow \, +])
\]

Furthermore, for all program $P$: $\alpha \circ [P] = [P]^{\#} \circ \alpha$

An over-approximation of the final states is computed by $[P]^{\#}(\top)$. 
Example

Factorial function:

```plaintext
input(x0);
x1 = 1;
x2 = 1;
while (x0 > 0) {
    x1 = x0 · x1;
    x0 = x0 - x2;
}
```

Abstract state before the loop: $(\pm, \pm, \pm)$

Iterates on the loop:

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<td>$\top$</td>
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<tr>
<td>x1</td>
<td>$\pm$</td>
<td>$\pm$</td>
<td>$\pm$</td>
</tr>
<tr>
<td>x2</td>
<td>$\pm$</td>
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<td>$\pm$</td>
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Abstract state after the loop: $(\top, \pm, \pm)$
Outline

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Soundness of the Analysis

We have designed the abstract semantics so that it satisfies the following theorem:

**Theorem: analysis soundness**

For any program $P$, we have:

$$\alpha \circ [P] = [P]^{\#} \circ \alpha$$

- By the theorem, $[P]^{\#}$ is sound, i.e., always over-approximates the concrete behaviors of the program.
- Moreover, the computation of $[P]^{\#}$ is fully automatic.

However, we have built $[\cdot]^{\#}$ under pretty restrictive assumptions.
Limitation 1: exact transformers

Assumption 1

For each concrete operation $f : \mathcal{P}(S) \to \mathcal{P}(S)$, we have built a abstract transformer $f^\# : D^\# \to D^\#$ such that: $\alpha \circ f = f^\# \circ \alpha$

Can we always build such a function?

We consider a simple, contrived example, to show that this assumption is too strong:

- **Abstract domain:** basic lattice of signs (first version we presented, with five elements)
- **Concrete function:** $f : S \mapsto \{\phi(x) \mid x \in S\}$  
  where $\phi(x) = 1$ if $x$ is even and $\phi(x) = -1$ if $x$ is odd
- We start from $X = \{2\}$: 
  $\alpha \circ f(X) = \alpha(f(\{2\})) = \alpha(\{1\}) = \pm$  
  However, $\alpha(X) = \pm$, and $f(\gamma(\pm)) = \{-1, 1\}$, thus, the most precise result we may expect for $f^\#(\alpha(X))$ is $\top$...
Limitation 2: Existence of the Best Abstraction

Even worse:

**Assumption 2**
We assumed the existence of the best abstraction function

\[ \alpha : \mathcal{P}(\mathcal{S}) \rightarrow D^# \]

**Abstract domain of convex polyhedra:**
- abstract values: *finite conjunctions of linear inequalities*
- very expressive: relations between variables

Without this assumption, we cannot even use \( \alpha \) anymore
Limitation 3: Finite Height Lattice

Assumption 3
We assumed the abstract lattice has finite height

We consider the lattice of intervals:

Without this assumption, we cannot guarantee the termination of the analysis anymore
Assumptions

We still assume a finite height domain.
We are going to first relax all assumptions related to $\alpha$:

Concretization based framework

We assume:
- an abstract domain $D^\#$ with an abstract order relation $\sqsubseteq$
- a monotone concretization function

$$\gamma : D^\# \rightarrow \mathcal{P}(\mathbb{Z}^n)$$

For a concrete operation $f : \mathcal{P}(\mathbb{Z}^n) \rightarrow \mathcal{P}(\mathbb{Z}^n)$, we set up abstract transfer function $f^\# : D^\# \rightarrow D^\#$ such that:

$$\forall S^\# \subseteq D^\#, \ f \circ \gamma(S^\#) \subseteq \gamma \circ f^\#(S^\#)$$

Exercise: example function $f : S \mapsto \{\phi(x) \mid x \in S\}$
where $\phi(x) = 1$ if $x$ is even and $\phi(x) = -1$ if $x$ is odd
Comparing Frameworks

We are indeed looking at a relaxed framework:

If:

- there exists an abstraction function \( \alpha : P(\mathbb{Z}^n) \rightarrow D^\# \)
- \( \alpha \circ \gamma = \text{id} \) (or, equivalently \( \alpha \) is surjective)
- \( f^\# : D^\# \rightarrow D^\# \) is such that \( \alpha \circ f = f^\# \circ \alpha \)

Then:

\[
\forall S^\# \subseteq D^\#, \ f \circ \gamma(S^\#) \subseteq \gamma \circ f^\#(S^\#)
\]

**Proof:** \( \alpha \circ f \circ \gamma(S^\#) = f^\# \circ \alpha \circ \gamma(S^\#) = f^\#(S^\#) \) thus, \( \alpha \circ f \circ \gamma(S^\#) \subseteq f^\#(S^\#) \), and \( f \circ \gamma(S^\#) \subseteq \gamma \circ f^\#(S^\#) \).

The condition \( \alpha \circ \gamma = \text{id} \) is not too restrictive (no redundant abstract elements).

Why? Is it always satisfied for all lattices seen so far?
A more realistic static analysis

Sound Assignment Operator

Definition

We consider the concrete operation:

$$\text{assign}_{x \leftarrow e} : M \mapsto \{ \sigma[x \leftarrow [e](\sigma)] \mid \sigma \in M \}$$

A sound abstract assignment operator is a function $\text{assign}^\#: D^\# \rightarrow D^#$ such that:

$$\forall S^\# \in D^# \text{, } \text{assign}_{x \leftarrow e} \circ \gamma(S^#) \subseteq \gamma \circ \text{assign}^\#_{x \leftarrow e}(S^#)$$

Then, the analysis of all assign statements relies on $\text{assign}^\#_{x \leftarrow e}$:

- **Case of** $x = v$: $\llbracket x = v \rrbracket^#(S^#) = \text{assign}^\#_{x \leftarrow v}(S^#)$
- **Case of** $x_0 = x_1 \odot x_2$: $\llbracket x_0 = x_1 \odot x_2 \rrbracket^#(S^#) = \text{assign}^\#_{x_0 \leftarrow x_1 \odot x_2}(S^#)$
- **Case of** input$(x)$: $\llbracket \text{input}(x) \rrbracket^#(S^#) = \text{assign}^\#_{x_0 \leftarrow \text{random}}(S^#)$
A more realistic static analysis

Sound Assignment Operator for a Non-Relational Domain

Thus, the real question is how do we define assign\(\#\leftarrow\)?

Case of a non relational domain: \(D^\# = (D^\#_V)^n\), where \(D^\#_V = D^\#_I, D^\#_C, \ldots\)

- **Definition of \([e]^{\#} : D^\# \to D^\#_V\)** by induction over the syntax:
  
  \[
  \begin{align*}
  \lbrack v\rbrack^\#(S^\#) &= v^\# & \text{where } v^\# \text{ is such that } \{v\} \subseteq \gamma(v^\#) \\
  \lbrack x_k\rbrack^\#(S^\#) &= S^\#_k \\
  \lbrack e_0 \odot e_1\rbrack^\#(S^\#) &= \odot^\#(\lbrack e_0\rbrack^\#(S^\#), \lbrack e_1\rbrack^\#(S^\#)) \\
  \lbrack \text{random}\rbrack^\#(S^\#) &= \top \text{ or other approximation of } \mathbb{Z}_+^*
  \end{align*}
  \]

  (very general definition, ways more than that of the expression language)

- **Definition of assign\(\#\leftarrow\):**

  \[
  \text{assign}_{\#\leftarrow e}(S^\#) = S^\#[x \leftarrow \lbrack e\rbrack^\#(S^\#)]
  \]

Case of relational lattices: specific algorithms, next courses
A more realistic static analysis

Sound Join Operator

When using $\alpha$ we relied on commutation with union, but with $\gamma$ this may not work... plus we did not even assume that the abstract domain has a full lattice structure!

Definition

A sound abstract join operator is a binary operator $\sqcup^\#: D^\# \times D^\# \rightarrow D^\#$ such that:

$$\forall S_0^\#, S_1^\# \in D^\#, \quad \gamma(S_0^\#) \cup \gamma(S_1^\#) \subseteq \gamma(S_0^\# \sqcup^\# S_1^\#)$$

- **Non relational domain:** $D^\# = (D^\#_V)^n$, where $D^\#_V = D^\#_I, D^\#_C, \ldots$. We can use the same pointwise operator definition as in the previous lecture (specific case of signs):

$$S_0^\# \sqcup^\# S_1^\# = k \mapsto (S_0^\#(x_k) \sqcup^\# S_1^\#(x_k))$$

- **In a relational abstract domain** (e.g., convex polyedra): convex closure algorithms

Xavier Rival
Abstract Interpretation: Introduction
May 5th, 2017 63 / 78
Sound Condition Test Operator

To analyze condition tests, we also need a sound abstraction of conditions:

**Definition**

We consider the concrete operation:

\[
\text{test}_{x \geq 0} : M \mapsto \{ \sigma \in M \mid \sigma(x) \geq 0 \}
\]

A sound abstract test operator is a function \( \text{test}^\#_{x \geq 0} : D^\# \rightarrow D^\# \) such that:

\[
\text{test}_{x \geq 0}(\gamma(S^\#)) = \{ \sigma \in \gamma(S^\#) \mid \sigma(x) \geq 0 \} \subseteq \gamma(\text{test}^\#_{x \geq 0}(S^\#))
\]

- We observe we get the same definition as in the previous lecture
- Analysis of a condition statement:

\[
[[\text{if}(x \geq 0)\{P\}]]^\#(S^\#) = [[P]]^\#(\text{test}^\#_{x \geq 0}(S^\#)) \sqcup^\# \text{test}^\#_{x < 0}(S^\#)
\]
Approximate Fixpoint Transfer

Transfer theorem

We consider two functions $f : \mathcal{P}(\mathbb{Z}^n) \to \mathcal{P}(\mathbb{Z}^n)$ and $f^\#: D^\# \to D^\#$ and two elements $M_0 \in \mathcal{P}(S)$, $S_0^\# \in D^\#$ such that:

- $f$ is continuous
- $f^\#$ is monotone
- $f \circ \gamma \subseteq \gamma \circ f^\#
- M_0 \subseteq \gamma(S_0^\#)

Then, both $f$ and $f^\#$ have a least-fixpoint and $\text{lfp}_{M_0} f \subseteq \gamma(\text{lfp}_{S_0^\#} f^\#)$

- **Proof**: quite similar to that of the exact case
- **Computation**: we can still compute the abstract lfp by iterative technique, as the abstract lattice is of finite height
- But: this still requires $f^\#$ to be monotone...
New static analysis

We can now summarize the definition of a new static analysis:

\[
\begin{align*}
\llbracket x_i = n \rrbracket (S^\#) &= \text{assign}_{x_i \leftarrow n}(S^\#) \\
\llbracket x_i = x_j + x_k \rrbracket (S^\#) &= \text{assign}_{x_i \leftarrow x_j + x_k}(S^\#) \\
\llbracket x_i = x_j - x_k \rrbracket (S^\#) &= \text{assign}_{x_i \leftarrow x_j - x_k}(S^\#) \\
\llbracket x_i = x_j \cdot x_k \rrbracket (S^\#) &= \text{assign}_{x_i \leftarrow x_j \cdot x_k}(S^\#) \\
\llbracket \text{input}(x_i) \rrbracket (S^\#) &= S^\#[i \leftarrow \pm] \\
\llbracket \text{if}(x_i > 0) \ P_0 \ \text{else} \ P_1 \rrbracket (S^\#) &= \llbracket P_0 \rrbracket (\text{test}_{x_i \geq 0}(S^\#)) \cup \llbracket P_1 \rrbracket (\text{test}_{x_i < 0}(S^\#)) \\
\llbracket \text{while}(x_i > 0) \ P \rrbracket (S^\#) &= \text{test}_{x_i \leq 0}(\text{lfp}_{S^\#} f^\#) \ \text{where} \\
f^\# : S^\# &\mapsto S^\# \cup \llbracket P \rrbracket (\text{test}_{x_i > 0}(S^\#))
\end{align*}
\]
Abstract Semantics Soundness

The analysis function \([\cdot]\)^\# obtained is sound:

\[\text{Theorem} \]

For all program \(P\),
\[\llbracket P\rrbracket(\gamma(S^\#)) \subseteq \gamma(\llbracket P\rrbracket^\#(S^\#))\]

Soundness means no concrete behavior is left out: the analysis can be used to verify programs...

There are now \textbf{more possible origins for imprecisions:}

- the abstraction may not allow to express properties required for the proof
- even if the abstraction expresses the properties required for the proof, specific analysis steps may compute less precise properties
Outline

1. Abstraction
2. Abstract interpretation
3. Applications of abstract interpretation
4. A basic static analysis
5. A more realistic static analysis
6. Termination of the Static Analysis
7. Conclusion
Enforcing Termination with Widening

We now relax the finite height assumption, which would cause the current framework not to guarantee analysis termination.

Computing invariants about infinite executions with widening \( \nabla \)

- **Widening** \( \nabla \) over-approximates \( \cup \): soundness guarantee
- **Widening** \( \nabla \) guarantees the termination of the analyses
- Typical choice of \( \nabla \): remove unstable constraints

**Example**: iteration of the translation \((2, 1)\), with octagons

![Initial](initial.png) \[ X_0 \]

![Iteration 1](iteration_1.png) \[ X_0 \Rightarrow X_1 = X_0 \nabla F(X_0) \]

![Iteration 2: Stable](iteration_2.png) \[ X_1 \Rightarrow F(X_1) \text{ (stable)} \]
To enforce convergence, we use a widening operator:

**Definition**

A widening operator is a binary operator \( \triangledown \) such that:

1. **\( \triangledown \) over-approximates join:**
   \[
   \forall S_0^#, S_1^# \in D^#, \quad \gamma(S_0^#) \cup \gamma(S_1^#) \subseteq \gamma(S_0^# \triangledown S_1^#)
   \]

2. **\( \triangledown \) enforces termination:**
   
   if 
   \[
   S_{n+1}^# = S_n^# \triangledown Q_n^#
   \]

   the sequence \( S_n^# \) converges, whatever the values of \( Q_n^# \)
Termination of the Static Analysis

Widening Operator for the Domain of Intervals

Definition

We let \( \triangledown \) be defined by:

\[
[a, b_0] \triangledown [a, b_1] = \begin{cases} [a, +\infty] & \text{if } b_0 < b_1 \\ [a, b_0] & \text{if } b_0 \geq b_1 \end{cases}
\]

and symmetrically for the left bound of intervals

Intuition:

- an interval is made of 0, 1 or 2 constraints (infinite bounds count for no constraint)
- widening may either keep the same number of constraints (stable) or decrease it
- but then, it can decrease it at most twice after that, all precision is lost, and the analysis necessarily converges
Fixpoint Approximation

Theorem

We consider two functions $f : \mathcal{P}(\mathbb{Z}^n) \to \mathcal{P}(\mathbb{Z}^n)$ and $f^\# : D^\# \to D^\#$ and two elements $M_0 \in \mathcal{P}(\mathbb{S})$, $S_0^\# \in D^\#$ such that:

- $\nabla$ is a widening over $D^\#$
- $f$ is continuous
- $f \circ \gamma \subseteq \gamma \circ f^\#$
- $M_0 \subseteq \gamma(S_0^\#)$
- $S_{n+1}^\# = S_n^\# \nabla f^\#(S_n^\#)$

Then, both $f$ has a least-fixpoint and $(S_n^\#)_{n \in \mathbb{N}}$ converges, with a limit $S_\infty^\#$; moreover:

\[
\begin{align*}
\text{lfp} \quad f & \subseteq \gamma(S_\infty^\#) \\
M_0 & \subseteq \gamma(S_\infty^\#)
\end{align*}
\]

- No more assumption on the monotonicity of $f^\#$!

Xavier Rival

Abstract Interpretation: Introduction

May 5th, 2017
Fixpoint Approximation: Proof

Proof:
- The existence of \texttt{Ifp} \( f \) follows from Kleene’s theorem; moreover:
  \[
  \bigcup_{n \in \mathbb{N}} f^n(M_0)
  \]
- We can prove by induction over \( n \) that:
  \[
  \bigcup_{k=0}^{n} f^k(M_0) \subseteq \gamma(S_n^\#)
  \]
- The definition of the widening entails that the sequence of abstract iterates is stationary, thus defines its limit \( S_\infty^\# \). We let \( N \) be the first rank such that \( S_N^\# = S_\infty^\# \).
- By definition of the least upper bound:
  \[
  \bigcup_{k \in \mathbb{N}} f^k(M_0) \subseteq \gamma(S_\infty^\#)
  \]
Static Analysis of a Loop with a Widening Operator

- Widening in the combined domain: **pointwise application of** $\nabla$
- The analysis simply replaces $\sqcup$ with $\nabla$
- We recover the analysis convergence and soundness theorem:

**Static analysis of a loop while**($x \geq 0$){$P$}

\[
\begin{cases}
S_0^\# & \text{be an approximation of the states before the loop} \\
S_{n+1}^\# = S_n^\# \nabla [\lbrack P \rbrack]^\#(\text{test}_{x \geq 0}^\#(S_n^\#))
\end{cases}
\]

Then this sequence of iterates is monotone and eventually converges after finitely many iterates; we let $S_\infty^\# = \lim S_n^\#$.

Moreover, $S_\infty^\#$ is a sound over-approximation of the states at the loop head and:

\[
[\text{while}(x \geq 0)\{P\}](\gamma(S^\#)) \subseteq \gamma(\text{test}_{x < 0}^\#(S_\infty^\#))
\]
Improving Widening Operators

Widening may lose a lot of precision, if not used carefully

Thus, we generally employ many tricks:

- **do not apply widening right away**, i.e., use $\sqcup$ instead, for the first abstract iterations
- **use thresholds**, to give a chance to intermediate bounds to stabilize
- **alternate join and widening iterations**
- **compute post-widening iterations**: keep iterating the abstract function after the widening iteration computes an abstract post-fixpoint
An example analysis with the interval abstract domain

**Classical techniques:**
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```c
int x = 0;

while(TRUE){
    if(x < 10000) {
        9999 will be a threshold value at loop head
        x = x + 1;
    } else {
        x = -x;
    }
}
```
An example analysis with the interval abstract domain

**Classical techniques:**
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```plaintext
int x = 0;
x \in [0, 0]
while (TRUE) {
    if (x < 10000) {
        9999 will be a threshold value at loop head
        x = x + 1;
    } else {
        x = -x;
    }
}
```
An example analysis with the interval abstract domain

**Classical techniques:**
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```c
int x = 0;
    x ∈ [0, 0]
while(TRUE){
    x ∈ [0, 0]
    if(x < 10000){
        9999 will be a threshold value at loop head
        x = x + 1;
    } else {
        x = -x;
    }
}
```

Entering the loop
Termination of the Static Analysis

An example analysis with the interval abstract domain

**Classical techniques:**
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```c
int x = 0;
x ∈ [0, 0]
while(TRUE){
x ∈ [0, 0]
  if(x < 10000){ 9999 will be a threshold value at loop head
    x ∈ [0, 0]
x = x + 1;
  }
  else {
    x ∈ ∅
x = −x;
  }
}
```

Only true branch possible
An example analysis with the interval abstract domain

**Classical techniques:**
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```cpp
int x = 0;
    x ∈ [0, 0]
while (TRUE) {
    x ∈ [0, 0]
    if (x < 10000) {
        9999 will be a threshold value at loop head
        x = x + 1;
        x ∈ [1, 1]
    } else {
        x ∈ ∅
        x = −x;
        x ∈ ∅
    }
}
```

Incrementation of interval
An example analysis with the interval abstract domain

**Classical techniques:**
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```c
int x = 0;
    x ∈ [0, 0]
while(TRUE) {
    x ∈ [0, 0]
    if (x < 10000) {
        9999 will be a threshold value at loop head
        x ∈ [0, 0]
        x = x + 1;
        x ∈ [1, 1]
    } else {
        x ∈ ∅
        x = −x;
        x ∈ ∅
    }
}
    x ∈ [1, 1]
```

Propagation
An example analysis with the interval abstract domain

Classical techniques:
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```c
int x = 0;
    x ∈ [0, 0]
while(TRUE){
    x ∈ [0, 1]
    if(x < 10000){
        x ∈ [0, 0]
        x = x + 1;
        x ∈ [1, 1]
    } else {
        x ∈ ∅
        x = −x;
        x ∈ ∅
    }
} else {
    x ∈ [1, 1]
}
```

Join at loop head
Termination of the Static Analysis

An example analysis with the interval abstract domain

Classical techniques:
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```plaintext
int x = 0;
  x ∈ [0, 0]
while(TRUE){
  x ∈ [0, 1]
  if(x < 10000){
    9999 will be a threshold value at loop head
    x ∈ [0, 1]
    x = x + 1;
    x ∈ [1, 1]
  } else {
    x ∈ ∅
    x = -x;
    x ∈ ∅
  }
}
x ∈ [1, 1]
```

Still only the true branch is possible
An example analysis with the interval abstract domain

**Classical techniques:**
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```java
int x = 0;
    x ∈ [0, 0]
while (TRUE) {
    x ∈ [0, 1]
    if (x < 10000) {
      9999 will be a threshold value at loop head
      x ∈ [0, 1]
      x = x + 1;
      x ∈ [1, 2]
    } else {
      x ∈ Ø
      x = −x;
      x ∈ Ø
    }
  }
  x ∈ [1, 1]
```

Incrementation of interval
An example analysis with the interval abstract domain

Classical techniques:

- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```plaintext
int x = 0;
    x ∈ [0, 0]
while(TRUE){
    x ∈ [0, 1]
    if(x < 10000){
        9999 will be a threshold value at loop head
        x ∈ [0, 1]
        x = x + 1;
        x ∈ [1, 2]
    }
    else {
        x ∈ ∅
        x = −x;
        x ∈ ∅
    }
    x ∈ [1, 2]
}
```

Propagation
An example analysis with the interval abstract domain

**Classical techniques:**
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```plaintext
int x = 0;
  x ∈ [0, 0]
while(TRUE){
  x ∈ [0, 9999]  instead of [0, +∞[ 
  if(x < 10000){
    9999 will be a threshold value at loop head
    x ∈ [0, 1]
    x = x + 1;
    x ∈ [1, 2]
  }
  else {
    x ∈ ∅
    x = −x;
    x ∈ ∅
  }
} x ∈ [1, 2]
```

*Widening at the loop head, + threshold*
Termination of the Static Analysis

An example analysis with the interval abstract domain

**Classical techniques:**
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```plaintext
int x = 0;
    x ∈ [0, 0]
while(TRUE){
    x ∈ [0, 9999] instead of [0, +∞[  
        if(x < 10000){         9999 will be a threshold value at loop head
            x ∈ [0, 9999]
            x = x + 1;
            x ∈ [1, 2]
        } else {  
            x ∈ ∅
            x = −x;
            x ∈ ∅
        }
}
    x ∈ [1, 2]
}
```

Now both branches are possible...
An example analysis with the interval abstract domain

**Classical techniques:**
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```plaintext
int x = 0;

x ∈ [0, 0]

while(TRUE){
  x ∈ [0, 9999] instead of [0, +∞[
  if(x < 10000){
    9999 will be a threshold value at loop head
    x ∈ [0, 9999]
    x = x + 1;
    x ∈ [1, 10000]
  }
  else {
    x ∈ ∅
    x = −x;
    x ∈ ∅
  }
}

x ∈ [1, 2]
```

Numerical assignments
An example analysis with the interval abstract domain

Classical techniques:
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```c
int x = 0;
for (x = 0; x != 9999; x++)
  if (x < 10000)
    x = x + 1;
else
  x = -x;
```

Join at the end of the loop
An example analysis with the interval abstract domain

Classical techniques:
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```java
int x = 0;
    x ∈ [0, 0]
while(TRUE){
    x ∈ [0, 10000] instead of [−∞, +∞]
    if(x < 10000){
        x ∈ [0, 9999]
        x = x + 1;
        x ∈ [1, 10000]
    } else {
        x ∈ ∅
        x = −x;
        x ∈ ∅
    }
} x ∈ [1, 10000]
```

Join after widening
An example analysis with the interval abstract domain

Classical techniques:
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```
int x = 0;
    x ∈ [0, 0]
while(TRUE){
    x ∈ [0, 10000] instead of [−∞, +∞]
    if(x < 10000){
        9999 will be a threshold value at loop head
        x ∈ [0, 9999]
        x = x + 1;
        x ∈ [1, 10000]
    }
    else {
        x ∈ [10000, 10000] instead of [10000, +∞[  
        x = −x;
        x ∈ ∅
    }
 x ∈ [1, 10000]
}
```

True branch stable, false branch visited for the first time
An example analysis with the interval abstract domain

**Classical techniques:**
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```plaintext
int x = 0;
    x ∈ [0, 0]
while(TRUE){
    x ∈ [0, 10000] instead of [−∞, +∞[  
    if(x < 10 000){  9999 will be a threshold value at loop head  
        x ∈ [0, 9999]
        x = x + 1;
        x ∈ [1, 10000]
    }
    else {
        x ∈ [10000, 10000] instead of [10000, +∞[  
        x = −x;
        x ∈ [−10000, −10000]
    }
}

x ∈ [1, 10000]
```

True branch stable, false branch visited for the first time
An example analysis with the interval abstract domain

Classical techniques:
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```plaintext
int x = 0;
    x ∈ [0, 0]
while(TRUE){
    x ∈ [0, 10000]  instead of [−∞, +∞[ 
    if(x < 10000){ 
        9999 will be a threshold value at loop head 
        x ∈ [0, 9999] 
        x = x + 1; 
        x ∈ [1, 10000] 
    } else { 
        x ∈ [10000, 10000]  instead of [10000, +∞[ 
        x = −x; 
        x ∈ [−10000, −10000] 
    }
}
    x ∈ [−10000, 10000] 
}
```

Join at the end of the loop
An example analysis with the interval abstract domain

Classical techniques:
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```plaintext
int x = 0;
x ∈ [0, 0]
while(TRUE){
x ∈ [−10000, 10000] instead of [−∞, +∞]
if(x < 10000){
    9999 will be a threshold value at loop head
    x ∈ [0, 9999]
x = x + 1;
x ∈ [1, 10000]
} else {
    x ∈ [10000, 10000] instead of [10000, +∞[n
    x = −x;
x ∈ [−10000, −10000]
}
x ∈ [−10000, 10000]
}
```

Join again: no widening after visiting a new branch
An example analysis with the interval abstract domain

Classical techniques:
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```plaintext
int x = 0;
  x ∈ [0, 0]
while(TRUE){
  x ∈ [−10000, 10000] instead of [−∞, +∞]
  if(x < 10000){
    9999 will be a threshold value at loop head
    x ∈ [−10000, 9999]
    x = x + 1;
    x ∈ [1, 10000]
  }
  else {
    x ∈ [10000, 10000] instead of [10000, +∞[ 
    x = −x;
    x ∈ [−10000, −10000]
  }
}
  x ∈ [−10000, 10000]
```

Branches
An example analysis with the interval abstract domain

**Classical techniques:**
- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```plaintext
int x = 0;
  x ∈ [0, 0]
while(TRUE){
  x ∈ [−10000, 10000]  instead of [−∞, +∞[  
  if(x < 10000){  9999 will be a threshold value at loop head
    x ∈ [−10000, 9999]
    x = x + 1;
    x ∈ [−9999, 10000]
  } else {  
    x ∈ [10000, 10000]  instead of [10000, +∞[  
    x = −x;
    x ∈ [−10000, −10000]
  }
}
```

Incrementation of interval in true branch; false branch stable
An example analysis with the interval abstract domain

**Classical techniques:**

- add test values and neighbors as thresholds
- alternate join and widening
- no widening after visiting a new branch

```plaintext
int x = 0;
    x ∈ [0, 0]
while(TRUE){
    x ∈ [−10000, 10000] instead of [−∞, +∞[  
    if(x < 10 000){  
        9999 will be a threshold value at loop head  
        x ∈ [−10000, 9999]  
        x = x + 1;  
        x ∈ [−9999, 10000]  
    } else {  
        x ∈ [10000, 10000] instead of [10000, +∞[  
        x = −x;  
        x ∈ [−10000, −10000]  
    }  
}
    x ∈ [−10000, 10000]
```

Everything is stable; exact ranges inferred
Outline

1. Abstraction
2. Abstract interpretation
3. Applications of abstract interpretation
4. A basic static analysis
5. A more realistic static analysis
6. Termination of the Static Analysis
7. Conclusion
Conclusion

Summary

The last two lectures:

- abstraction and its formalization
- computation of an abstract semantics in a very simplified case
- more realistic static analysis

Next lectures:

- construction of a few non trivial abstractions
- more advanced static analysis techniques