

Automated verification of termination certificates

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Outline

Introduction

Certification Problem Format (CPF)

Definition and proof of the certificate verifier

Conclusions

Work plan for 2012

Motivation

- ▶ **Termination** is an important and difficult problem.
- ▶ In general this problem is undecidable.
- ▶ Many methods and criteria have been developed, and they are being used in various programs (termination provers)
- ▶ These programs are more and more complex, and their result difficult to check by hand.
- ▶ Every year in the **termination competition**, some tools are **disqualified** because of mistakes found in their proofs.
- ▶ For these tools to be used in the certification of critical systems and proof assistants, their results must be **certified**.

⇒ How to certify termination proofs generated by termination provers?

Example of criterion: polynomial interpretations

- ▶ for each function symbol f of arity n , we assume given an integer polynomial P_f with n variables
- ▶ a term t can then be interpreted by an integer polynomial P_t by recursively composing the polynomials interpreting the function symbols occurring in the term t
- ▶ then a program defined by a set R of rules terminates if:
 - ▶ each P_f is monotone in each variable
 - ▶ for every rule $l \rightarrow r \in R$, $P_l > P_r$ on \mathbb{N}

certificate: polynomials P_f

such a certificate is correct if the above conditions are satisfied

Goal

develop a safe, efficient and modular termination certificate verifier

our solution:

- ▶ write a verifier in Coq
- ▶ prove its correctness using the CoLoR library
- ▶ extract it to OCaml

XML

termination certificates are given as XML files

```
<root>
  <child1 attribute1="value1">
    <subchild> .... </subchild>
  </child1>
  <child2>
    <subchild> .... </subchild>
  </child2>
</root>
```

abstractly, an XML file is a tree whose nodes are tagged and may have attributes (leaves are strings)

XML Schema (XSD)

an XSD file describes a class of XML files by defining the possible tags and attributes, and how tagged elements can be composed

XSD type = set of XML elements

It is itself defined as an XML file! with the following tags:

- ▶ `<sequence>XSD_type1 XSD_type2 ...</sequence>`: product type
- ▶ `<choice>XSD_type1 XSD_type2 ...</choice>`: union type
- ▶ `<group name = "<name>">XSD_type</group>`: names a type
- ▶ `<element name = "<tag>">XSD_type</element>`: declares a tag, its attributes and its possible sons

remark: XSD definitions need not be ordered and can be forward or backward referenced

XSD example (part 1/2)

```
<xs:group name="symbol">
  <xs:annotation>
    <xs:documentation>is used as a function symbol in terms, orderings, ...</xs:documentation>
  </xs:annotation>
  <xs:choice>
    <xs:element ref="name"/>
    <xs:element name="sharp">
      <xs:complexType>
        <xs:sequence>
          <xs:group ref="symbol"/>
        </xs:sequence>
      </xs:complexType>
    </xs:element>
    <xs:element name="labeledSymbol">
      <xs:complexType>
        <xs:sequence>
          <xs:group ref="symbol"/>
          <xs:group ref="label"/>
        </xs:sequence>
      </xs:complexType>
    </xs:element>
  </xs:choice>
</xs:group>
```


CPF

the CPF format is regularly modified and extended with new features, it is useful to have a tool that can **automatically generate** in OCaml and Coq:

- ▶ data structures
- ▶ parsers
- ▶ pretty-printers

for that format

XSD example (part 2/2)

OCaml type corresponding to previous XSD definition:

```
type label =  
| Label_numberLabel of nonNegativeInteger list  
| Label_symbolLabel of symbol list  
and symbol =  
| Symbol_name of name  
| Symbol_sharp of symbol  
| Symbol_labeledSymbol of symbol * label
```

problem in Coq: detect mutually inductive types

Dependence relation

type expressions: $T = C | T \Rightarrow T$

type definition for a type constant C : a type for each constructor

C *def-depends* on D , written $C \rightsquigarrow D$, if:

C has a constructor in the type of which D occurs

let \succeq (\simeq) be the transitive (reflexive and symmetric) closure of \rightsquigarrow

- ▶ two types C and D *depend* on each other if $C \simeq D$
- ▶ a type C can be defined *before* D if $C < D$

the transitive closure and then the computation and the ordering of equivalence classes can be done using computations on boolean matrices

General picture

- ▶ define a boolean function:

```
Fixpoint check : cpf -> bool := ...
```

- ▶ prove that it is correct:

```
Lemma check_ok : forall c, check c = true -> WF (red c)
```

where (red c) is the rewrite relation defined in c

for returning useful information in case of failure, instead of bool
we use:

```
Inductive result (A : Type) : Type :=  
| Ok : A -> result A  
| Ko : error -> result A.
```

Translation of CPF types to CoLoR types

problem: CoLoR terms are defined wrt some arity function

```
Record Signature : Type := mkSignature {  
  symbol :> Type;  
  arity : symbol -> nat;  
  beq_symb : symbol -> symbol -> bool;  
  beq_symb_ok : forall x y, beq_symb x y = true <-> x = y  
}.
```

```
Inductive term : Type :=  
  | Var : variable -> term  
  | Fun : forall f : Sig, vector term (arity f) -> term.
```

- ▶ the arity can be computed by examining rules
- ▶ the translation of terms may fail

Remarks

- ▶ **check** requires many auxiliary functions for testing equality on CPF types (symbols, terms, etc.)
- ▶ the induction principles automatically generated by Coq for some types is too weak and needs to be redefined

this is currently done by hand \Rightarrow generate this automatically?

Problem with CoLoR

problem: CoLoR uses **modules and functors**

they need to be instantiated

- ▶ modules **cannot** be defined inside sections
- ▶ certificates are **recursive**

modules need to be first-class object

solution: change CoLoR to use **records** instead.

What did I do?

- ▶ generator of Coq type definitions from XSD
- ▶ generator of OCaml type definitions and parsing functions from XSD
- ▶ replaced modules by records in some CoLoR files
- ▶ translation of CPF types into CoLoR types
- ▶ definition of a certificate verifier for polynomial interpretations
- ▶ correctness proof almost finished for polynomial interpretations

Some figures

cpf.xsd: 1400 lines of XML

ml_of_xsd.ml: 400 lines of OCaml

 cpf.ml: 1200 lines of generated OCaml

coq_of_xsd.ml: 200 lines of OCaml

 cpf.v: 250 lines of generated Coq

rainbow.v: 1200 lines of Coq

 modified CoLoR files: 650 lines of Coq

What did I learn?

- ▶ XML and XML Schema
- ▶ OCaml
- ▶ more on Coq

remark: the use of dependent types in CoLoR makes definitions and proofs more difficult

Work plan for 2012

- ▶ November 2011: finish the new implementation of the OCaml and Coq type definitions generator from XSD
- ▶ December 2011: finish the correctness proof for polynomial interpretations on integers
- ▶ January 2012: extraction to OCaml, linking with CPF parser and testing on the Termination Problem Data Base (TPDB)
- ▶ February 2012 - October 2012: extension to other termination techniques

Thank you for your attention!!!