

Network Attacks

Chapter 2

Network & Security

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SUMMARY OF CHAPTER 2

- Denial of Service
- Spoofing
- Hijacking
- Conclusion and References

DENIAL OF SERVICE

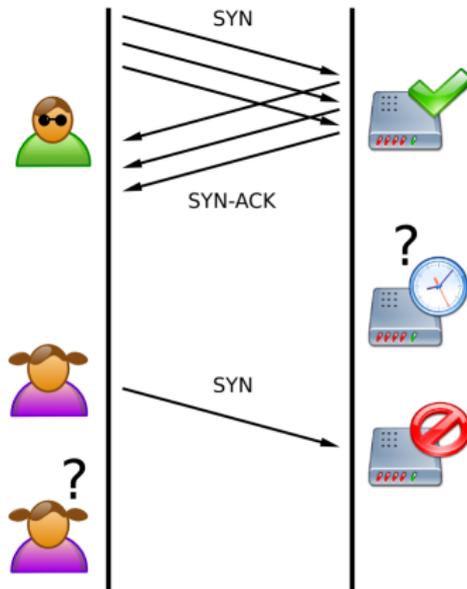
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Ping of Death (for Historical Purposes)

- Ping size should be 64 bytes (84 with IP header).
- Send IP packets that exceed the maximum legal length (65535 bytes).
- One of the earliest denial of service attack.
- Unix, Linux, Mac, Windows, printers, and routers were vulnerable (< 1997).

- Upon reception of the **SYN packet**, the server allocates necessary memory for the connection and enters it in a queue of **half open connections**.
- This situation having not been foreseen, the server can **no more accept new connections** once the queue overflows.
- The attacker can forge the source address of his SYN packets to remain **anonymous**.
- Current versions of operating systems are protected against such attacks.

SYN Flooding



SYN Flooding: Protections

- Increase the **size of the queue**.
- Reduce **timeout** during which server is waiting for an ACK.
- Drop the **oldest SYN** in the queue.
- **Filtering** eg on IP addresses.
- **SYN-Cache**: cache the SYN and send a SYN/ACK. If the ACK arrives, a complete connection is created.

SYN Flooding: Protections

- **SYN-Cookies:** Once the connection queue is almost filled up, the server uses SYN cookies.
- Upon reception of a SYN:
 - The server sends a SYN/ACK containing a **SYN cookie**.
 - The server **erases the SYN entry**.
- Upon reception of a ACK:
 - The server checks whether it contains a **valid cookie**. If so this highly likely means that the client has already sent a SYN and is so a honest client.

- SYN cookies are specific **Initial Sequence Numbers**.
 - t is a 5-bit counter incremented every 64 seconds modulo 32.
 - m is the Maximum Segment Size encoded on 3 bits.
 - s is the 24-bit result of a cryptographic function computed on t , the server IP address and port number, the client IP address and port number.

$$ISN = [t \bmod 32] || m || s$$

SYN Flooding: SYN-Cookie Check

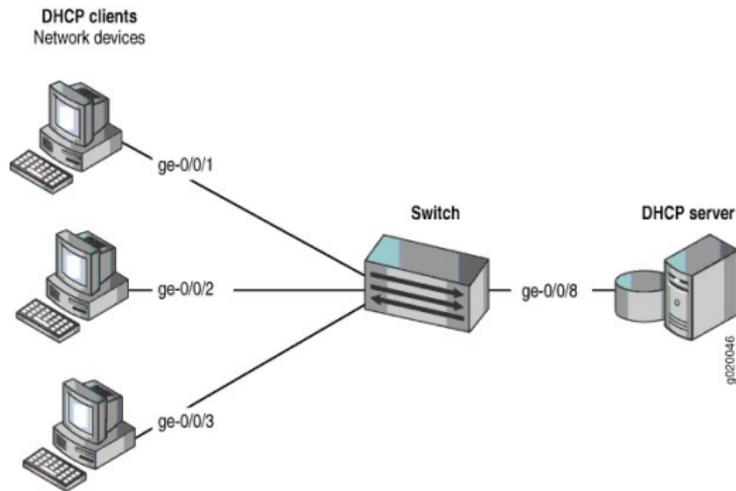
- Upon reception of an ACK, the server carries out the following operations.
 - Check that the received value t is valid with respect to the current time. Otherwise, this means the connection is expired.
 - Recompute s to check its validity.
 - Decodes the value m , which allows the server to reconstruct the SYN queue entry.

Kamikaze Packets (Xmas Tree Packets)

- TCP Packets with flags **URG**, **PSH**, and **FIN**.
- When many **Kamikaze packets** are sent, an unexpected behavior of **routers** may occur.
- Certain routers may **reboot**.

DHCP Starvation

- The attacker floods a **DHCP server** with DHCP requests from spoofed (counterfeit) MAC addresses.
- The server's **pool of IP addresses** is exhausted.

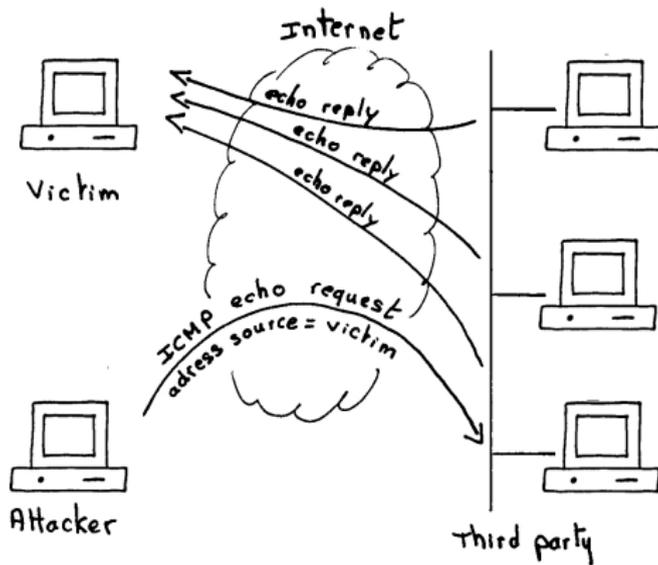


www.juniper.net

Smurf Attack

- Drown the target with the help of **traffic amplifiers**.
- Typical case: **ICMP echo-request** (ping).
- The hacker sends a **ping packet** with the target address as source address.
- The “pinged” machine sends its **response** to the target.
- If the hacker sends the packet to a **broadcast address**, all machines of the network will reply to the target.

Smurf Attack



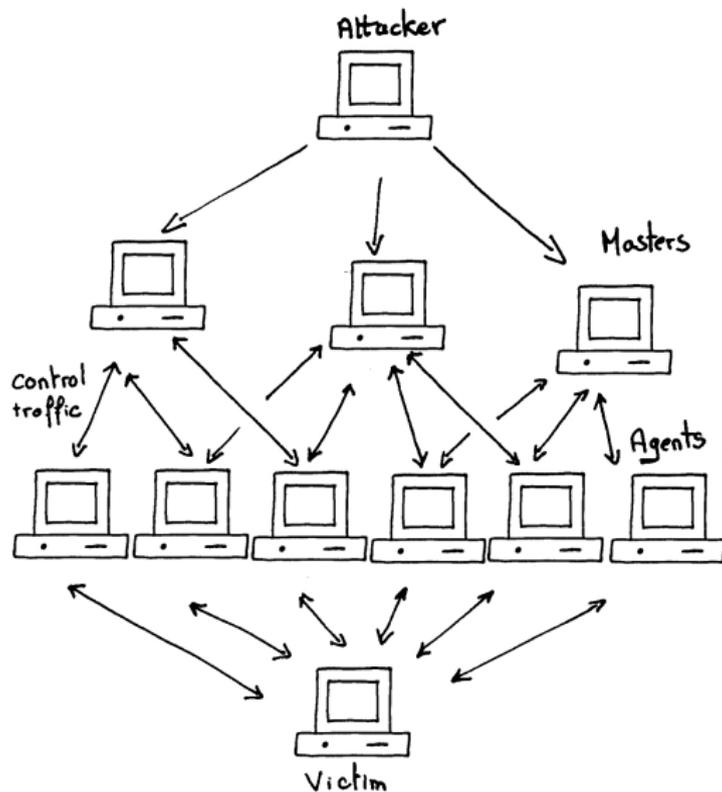
Smurf Attack: Protection

- Configure individual hosts and routers **not to respond to ping requests** to broadcast addresses.
- Configure routers **not to forward packets** directed to broadcast addresses.
- Magnifying the traffic can be done with applications where the **replies are much bigger than the requests**.

DDoS: Distributed Denial of Service

- To increase the efficiency of Denial of Service, hackers hack into **several machines and install agents** on them.
- Several **master** machines control the agents.
- The **hacker** sends commands to the **masters** which in turn execute the attack through the **agents**.

DDoS: Architecture



- **Botnet**: A network of hacked machines controlled by a hacker.
- The power (bandwidth) of the attack is **multiplied by the agents/bots**.
- Typically the bots connect to an Internet Relay Chat and wait for commands from their master.
- It is more **difficult to trace** the hackers (2 intermediate layers).
- Since attack comes from several sources, it is much more **difficult to filter it**.

- Hackers rent botnets to spammers for as low as \$350 per week for 5000 bots.
- In the press: 3 men arrested in **the Netherlands** in 2005; they managed a 1.5 million computer botnet.

DDoS: Historical Example

- **February 7, 2000:** The Internet portal of Yahoo was inaccessible for several hours.
- **February 8, 2000:** Amazon, Buy.com, eBay, and CNN were also victims of a DDoS attack, which significantly reduced their activities.
- **February 9, 2000:** E*Trade and ZDNet were both victim of a DDoS attack.

- **Trinoo.**
- **The Tribe Flood Network.**
- **Stacheldraht.**
- **Tribe Flood Network 2000 (tfn2k).**
 - Agents (bots) do not answer to the masters.
 - Masters send 20 command packets.
 - Masters use ICMP, TCP, UDP.
 - Communication encrypted.
- **Loic**, eg used by **Anonymous** in Operation Chanology (against the scientology Church) in 2010. Particularity: People voluntarily install Loic on their computer to join the botnet.

SPOOFING

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- In certain cases, the **IP source address** is used to **authorize** a connection.
 - Routers and firewalls can filter packets according to their source.
 - Some programs (rlogin, rsh) can authorize certain sources to connect without authentication.
- It is easy to forge a packet's source address and to **abuse the trust** of that source.
- The reply to a forged message is sent to the forged address.
- Easy to use with protocols based on **UDP**.
- The applications to be hacked (typically rlogin, rsh, ...) use **TCP**.

DNS Spoofing (UDP)

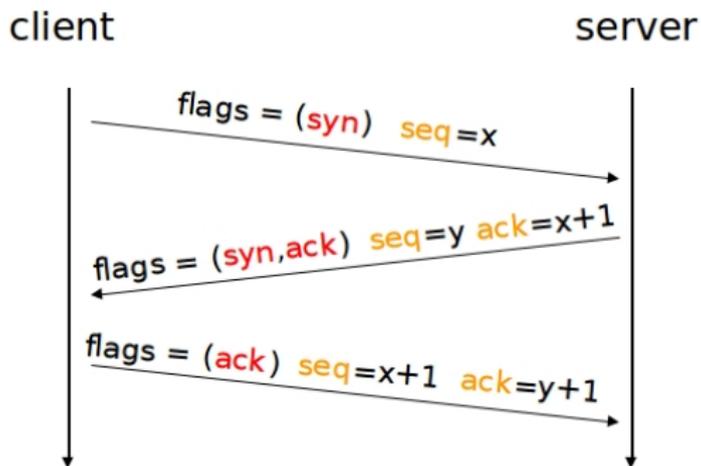
- A user sends a **DNS request** to a local DNS server.
- An attacker sends a **DNS response** faster than the DNS server.
- DNS is mostly based on **UDP**.

DNS Cache Poisoning (UDP)

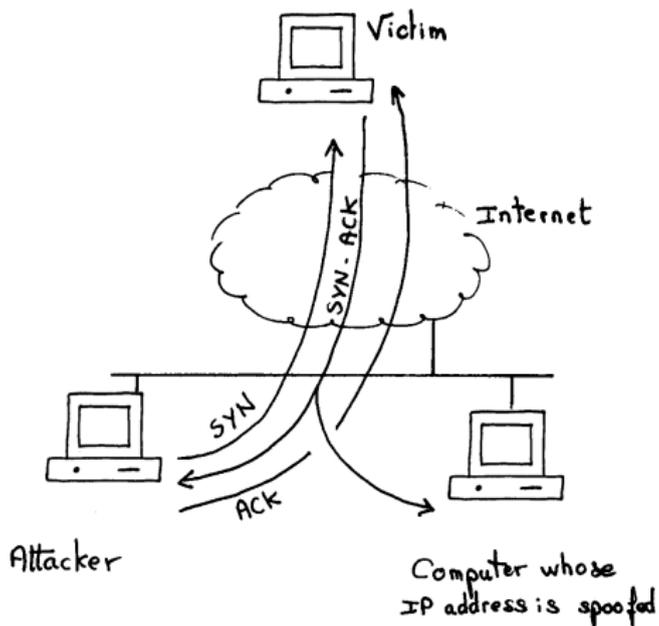
- The attacker sends a DNS request to a **local DNS server**.
- The local DNS server queries a **master DNS server**.
- The attacker **spoofs the master DNS server**, providing the local DNS server with a **fake DNS response**.
- However: the local DNS server's query includes an **identifier**.
- The attacker must **guess the identifier**.
- The attacker **floods the local DNS server** with DNS responses.
- The attacker may send **many DNS requests**.

- TCP is a sliding window protocol, it uses **sequence numbers** to keep track of sent and received data.
- To avoid using the same sequence numbers, a random **initial sequence number** (ISN) is chosen for each new connection.

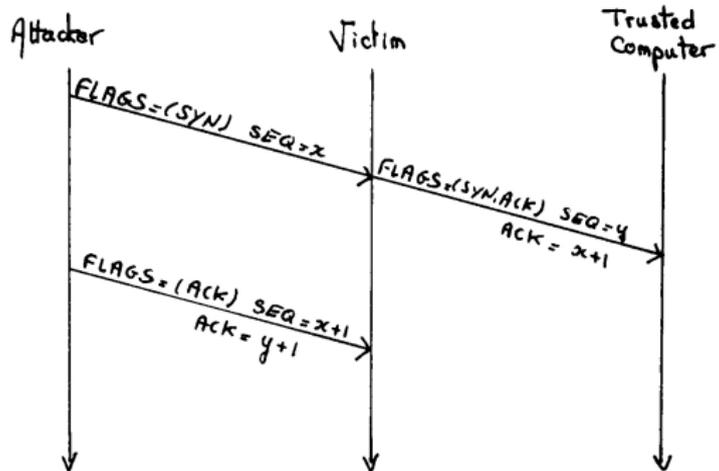
TCP/IP Spoofing: TCP Handshake



TCP/IP Spoofing Within a LAN



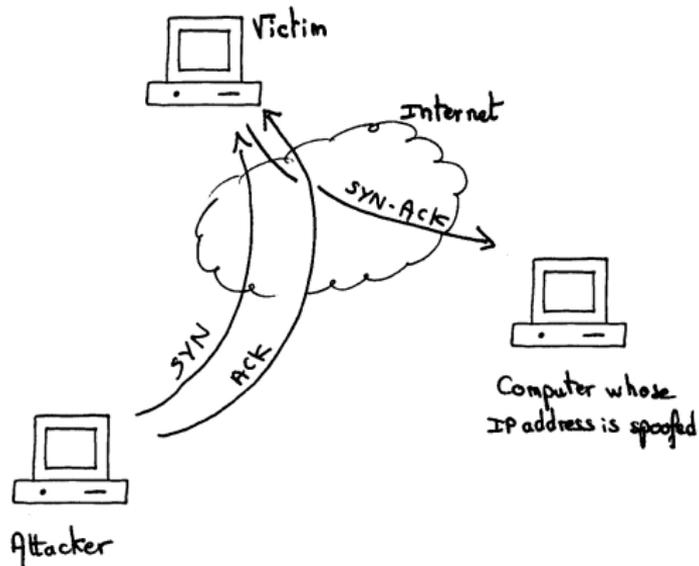
TCP/IP Spoofing Within a LAN



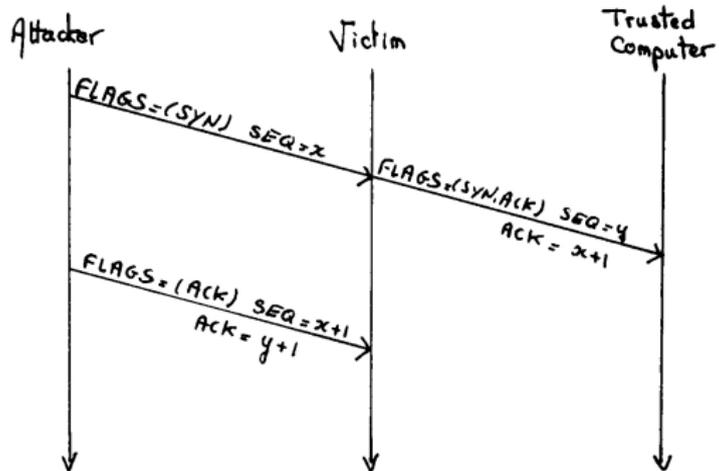
TCP/IP Spoofing Within a LAN

- The victim **resets** the handshake protocol.
- The hacker must **prevent the victim from responding**.

TCP/IP Spoofing From Outside



TCP/IP Spoofing From Outside



TCP/IP Spoofing From Outside: ISN Prediction

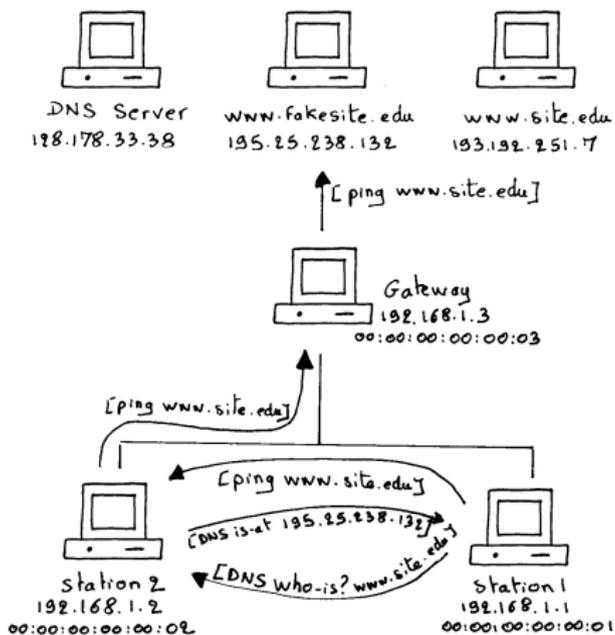
- The original standard (**RFC 793**) requires that the ISN be incremented once every four microseconds.
- In some simple TCP implementations the **next ISN can be predicted**.
- Hacker's procedure (ISN prediction):
 - He opens a few **authentic connections** (for example SMTP) to obtain the current ISN and increment samples.
 - He launches his **forged connection** using the last ISN plus an increment obtained from those samples.
 - He can launch **multiple forged connections** with different increments hoping that at least one is correct.

TCP/IP Spoofing From Outside: ISN Prediction

14:18:25.90	kevin.1000	>	bob.514:	S	1382726990	
14:18:26.09	bob.514	>	kevin.1000:	S	2021824000	ack 1382726991
14:18:26.17	kevin.1000	>	bob.514:	R	1382726991	128'000
14:18:26.50	kevin.999	>	bob.514:	S	1382726991	
14:18:26.69	bob.514	>	kevin.999:	S	2021952000	ack 1382726992
14:18:26.77	kevin.999	>	bob.514:	R	1382726992	128'000
14:18:27.01	kevin.998	>	bob.514:	S	1382726992	
14:18:27.17	bob.514	>	kevin.998:	S	2022080000	ack 1382726993
14:18:27.25	kevin.998	>	bob.514:	R	1382726993	128'000
14:18:27.54	kevin.997	>	bob.514:	S	1382726993	
14:18:27.71	bob.514	>	kevin.997:	S	2022208000	ack 1382726994
14:18:27.79	kevin.997	>	bob.514:	R	1382726994	128'000
14:18:28.05	kevin.996	>	bob.514:	S	1382726994	
14:18:28.22	bob.514	>	kevin.996:	S	2022336000	ack 1382726995
14:18:28.30	kevin.996	>	bob.514:	R	1382726995	128'000

- ARP: Address Resolution Protocol.
 - Protocol that helps finding a **layer 2** (Ethernet) address from a **layer 3** (IP) address.
- Very simple and **insecure**:
 - client: **who knows** the ethernet address of 10.1.2.3?
 - anybody: 10.1.2.3 **has** ethernet address 010203040506.
- It is easy to forge responses (even **non-solicited**) to redirect traffic.

ARP Poisoning



- **Dynamic ARP Inspection** (analyze consistencies of ARP packets).
- **DHCP Snooping** (detect fake DHCP servers).

HIJACKING

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- Instead of stealing a password, the hacker can wait until a user authenticates himself and then steal his session.
- This technique can be applied to **several layers**, eg, modem, TCP, HTTP.

Session Hijacking: Modem Session

- The modem gives access to a serial line (for ex. remote access).
- A user may drop the line without quitting the online session.
- The terminal's session **remains active** for a while.
- The next user (or hacker) who connects to the modem **finds the preceding user's session**.

- If a hacker can spy on a TCP connection, he can **insert a TCP packet** with correct sequence numbers.
- Inserting an additional packet in a TCP connection creates a **packet avalanche**:
 - The source, who has never sent the packet, **does not agree with the acknowledged sequence number** and emits an acknowledgement.
 - The destination, who has seen the packet, insists on the sequence number and also sends an acknowledgement.

Session Hijacking: HTTP “Session”

- HTTP protocol is **not session-oriented**.
- It is made of **independent requests/responses**.
- E-commerce web-sites use **artificial means** to recognize requests belonging to a session: cookies or personalized URLs.
- If the hacker can spy on these data, he can create **requests that would be part of the same session**.

Session Hijacking

The screenshot shows the Amazon product page for the book "Computer System Security: Basic Concepts and Solved Exercises (Computer and Communication Sciences)". The page includes the Amazon logo, search bar, and navigation links. The product title is "Computer System Security: Basic Concepts and Solved Exercises (Computer and Communication Sciences)", published on July 13, 2007, by Gildas Avoine, Philippe Oechslin, and Pascal Junod. The current price is \$61.15, with a list price of \$78.95. The page also features a "FREE TWO-DAY SHIPPING FOR COLLEGE STUDENTS" banner and a "Buy New" button.

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CONCLUSION AND REFERENCES

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- Most of the presented attacks are **known for a long while**.
- But they are still **up to date**.
- **ARP Poisoning** is an efficient attack.
- **Countermeasures** exist but they are (too) rarely deployed.

- SYN-cookies:

<http://cr.yp.to/syncookies.html>

- Christmas Tree Attacks:

<https://www.youtube.com/watch?v=bVrxL2AL4yQ>

- DNS spoofing:

<https://www.checkpoint.com/defense/advisories/public/dnsvideo/index.html>

- TCP hijacking:

<http://www.cs.berkeley.edu/~daw/security/shimo-post.txt>

- ARP poisoning:

<http://www.royabubakar.com/blog/2013/11/04/arp-poisoning-attack-and-mitigation-for-cisco-catalyst/>