

# Matchings on infinite graphs

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## Abstract

We prove that the local weak convergence of a sequence of graphs is enough to guarantee the convergence of their normalized matching numbers. The limiting quantity is described by a local recursion defined on the weak limit of the graph sequence. However, this recursion may admit several solutions, implying non-trivial long-range dependencies between the edges of a largest matching. We overcome this lack of correlation decay by introducing a perturbative parameter called the temperature, which we let progressively go to zero. When the local weak limit is a unimodular Galton-Watson tree, the recursion simplifies into a distributional equation, resulting into an explicit formula that considerably extends the well-known one by Karp and Sipser for Erdős-Rényi random graphs.

**Keywords:** local weak convergence, monomer-dimer systems, matching number, cavity method.

## 1 Introduction

A *matching* on a finite graph  $G = (V, E)$  is a subset of pairwise non-adjacent edges  $M \subseteq E$ . The  $|V| - 2|M|$  isolated vertices of  $(V, M)$  are said to be *exposed* by  $M$ . We let  $\mathbb{M}(G)$  denote the set of all possible matchings on  $G$ . The *matching number* of  $G$  is defined as

$$\nu(G) = \max_{M \in \mathbb{M}(G)} |M|, \quad (1)$$

and those  $M$  which achieve this maximum – or equivalently, have the fewest exposed vertices – are called *maximum matchings*.

Our first main result belongs to the theory of convergent graph sequences. Convergence of bounded degree graph sequences was defined by Benjamini and Schramm [7], Aldous and Steele [2], see also Aldous and Lyons [1]. The notion of local weak convergence has then inspired a lot of work [8], [11], [13], [16], [24]. Our first main contribution shows that for any sequence of graphs converging locally, the corresponding sequence of normalized matching numbers converges.

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Our second main contribution concerns sequences of graphs converging locally to Galton-Watson trees. A classical example in this framework is the sequence of Erdős-Rényi graphs with connectivity  $c$  denoted by  $G(n, c/n)$ : the limiting tree is then a Galton-Watson tree with degree distribution a Poisson distribution with parameter  $c$ . In this case, Karp and Sipser [22] showed that

$$\frac{\nu(G(n, c/n))}{n} \xrightarrow{n \rightarrow \infty} 1 - \frac{t_c + e^{-ct_c} + ct_c e^{-ct_c}}{2}, \quad (2)$$

where  $t_c \in (0, 1)$  is the smallest root of  $t = e^{-ct}$  (we will see in the sequel that the convergence is almost sure). The explicit formula (2) rests on the analysis of a heuristic algorithm now called Karp-Sipser algorithm. The latter is based on the following observation : if  $e \in E$  is a pendant edge (i.e. an edge with one end vertex of degree one) in  $G = (V, E)$ , then there is always a maximum matching that contains  $e$ , so all edges that are adjacent to  $e$  may be deleted without affecting  $\nu(G)$ . The first stage of the algorithm consists in iterating this until no more pendant edge is present. This is the *leaf-removal process*.  $G$  is thus simplified into a sub-graph with only isolated vertices, matched pairs, and a so-called *core* with minimum degree at least 2. As long as that core is non-empty, one of its edges is selected uniformly at random, the adjacent edges are deleted, and the whole process starts again. When the algorithm stops, the remaining edges clearly form a matching on  $G$ , but its size may be far below  $\nu(G)$  due to the sub-optimal removals on the core.

On  $G(n, c/n)$ , the dynamics of the deletion process can be approximated in the  $n \rightarrow \infty$  limit by differential equations which can be explicitly solved. In particular, the asymptotic size of both the optimal part constructed in the first stage, and the sub-optimal part constructed on the core can be evaluated up to an  $o(n)$  correcting term (which has been later refined, see [4]). Moreover, the second part happens to be almost perfect, in the sense that only  $o(n)$  vertices are exposed in the core. This guarantees that the overall construction is asymptotically optimal, and the asymptotic formula for  $\nu(G_n)$  follows. More recently, the same technique has been applied to another class of random graphs with a fixed log-concave degree profile [10], resulting in the asymptotical existence of an *almost perfect matching* on these graphs :

$$\frac{\nu(G_n)}{|V_n|} \xrightarrow{n \rightarrow \infty} \frac{1}{2}. \quad (3)$$

In both cases, the proof of optimality – and hence the asymptotic formula for  $\nu(G_n)$  – relies on the fact that the second stage exposes only  $o(n)$  vertices, which is bound to fail as soon as one considers more general graph ensembles where the core does not necessarily admits an almost-perfect matching. We give several such examples in the Appendix. Also, our first main result shows that the normalized matching number converges for such graphs, the computation of the limit requires another set of tools to solve a *recursive distributional equation* (a usual ingredient of the objective method, see [3]). This allows us to derive an explicit formula for the limit that considerably generalizes the aforementioned results.

The rest of our paper is organized as follows: we state our main results in the following Section 2. In Section 3, we extend the Boltzmann-Gibbs distribution over matchings on a finite graph to infinite graphs. This will allow us to derive our Theorem 1 in Section 4. We deal with the specific cases of trees (and random graphs) in Section 5. We finish by an appendix presenting a simpler version of our argument valid for specific graphs having a correlation decay property. We also give simple example of graphs lacking this property.

## 2 Results

Let us start with a brief recall on local weak convergence (see [7, 2] for details). A *rooted graph*  $(G, \circ)$  is a graph  $G = (V, E)$  together with the specification of a particular vertex  $\circ \in V$ , called the *root*. We let  $\mathcal{G}$  denote the set of all locally finite connected rooted graphs considered up to *rooted isomorphism*, i.e.  $(G_1, \circ_1) \equiv (G_2, \circ_2)$  if there exists a bijection  $\gamma: V_1 \rightarrow V_2$  that preserves roots ( $\gamma(\circ_1) = \circ_2$ ) and adjacency ( $uv \in E_1 \iff \gamma(u)\gamma(v) \in E_2$ ). In the space  $\mathcal{G}$ , a sequence  $\{(G_n, \circ_n); n \in \mathbb{N}\}$  *converges locally* to  $(G, \circ)$  if for every radius  $k \in \mathbb{N}$ , there is  $n_k \in \mathbb{N}$  such that

$$n \geq n_k \implies [G_n, \circ_n]_k \equiv [G, \circ]_k.$$

Here,  $[G, \circ]_k$  denotes the finite rooted subgraph induced by the vertices lying at graph-distance at most  $k$  from  $\circ$ . It is not hard to construct a distance which metrizes this notion of convergence and turns  $\mathcal{G}$  into a complete separable metric space. We can thus import the usual machinery of weak convergence of probability measures on Polish spaces (see e.g. [9]). There is a natural procedure for turning a finite deterministic graph  $G = (V, E)$  into a random element of  $\mathcal{G}$ : one simply chooses uniformly at random a vertex  $\circ \in V$  to be the root, and then restrains  $G$  to the connected component of  $\circ$ . The resulting law is denoted by  $\mathcal{U}(G)$ . If  $(G_n)_{n \in \mathbb{N}}$  is a sequence of finite graphs such that  $(\mathcal{U}(G_n))_{n \in \mathbb{N}}$  admits a weak limit  $\rho \in \mathcal{P}(\mathcal{G})$ , we call  $\rho$  the *random weak limit* of the sequence  $(G_n)_{n \in \mathbb{N}}$ . Finally, for any  $d \geq 0$ , we define  $\mathcal{G}_d$  as the space of all connected graphs with maximal degree no more than  $d$ .

Rather than just graphs  $G = (V, E)$ , it will be sometimes convenient to work with *discrete networks*  $G = (V, E, \mathcal{M})$ , in which the additional specification of a *mark map*  $\mathcal{M}: E \rightarrow \mathbb{N}$  allows to attach useful local information to edges, such as their absence/presence in a certain matching. We then simply require the isomorphisms in the above definition to preserve these marks.

The first main implication of our work is that the local weak convergence of a sequence of graphs is enough to guarantee the convergence of their normalized matching numbers to a quantity depending only upon the random weak limit of the graph sequence.

**Theorem 1** *Let  $G_n = (V_n, E_n), n \in \mathbb{N}$ , be a sequence of finite graphs admitting a random weak limit  $\rho$ . Then,*

$$\frac{\nu(G_n)}{|V_n|} \xrightarrow{n \rightarrow \infty} \gamma, \quad (4)$$

where  $\gamma \in [0, \frac{1}{2}]$ , defined in the forthcoming Equation (21), is described by a local recursion defined on the random weak limit  $\rho$ .

Since the work of Heilmann and Lieb [21], it is known that the thermodynamic limit for the monomer-dimer systems exists and basic properties of this limit are derived for lattices. In particular, [21, Lemma 8.7] shows the convergence of the matching number when the underlying graph is a lattice. More recently, Elek and Lippner [17] extend this result by using the framework of the local weak convergence for bounded degree graphs. Here we remove this last assumption. More important, the approach in [17] is highly non-constructive and does not allow to characterize the limit  $\gamma$ . In contrast, we give in the sequel an explicit formula for  $\gamma$  in terms of the random weak limit  $\rho$  which is particularly simple for bounded degree graphs, see Theorem 11 and (15). Our approach starts as in [21] with the introduction of a natural family of probability distributions on the set of matchings parametrized by a single parameter  $z > 0$

called the Boltzmann-Gibbs distribution. The analysis in [21] concentrates on the properties of the partition function also known as the matching polynomial from which a result like (4) could be deduced. Our analysis differs from this approach and concentrates on the analysis of the marginal probabilities of the Boltzmann-Gibbs distribution in a similar spirit as in the (non-rigorous) work of Zdeborová and Mézard [25]. Using a fundamental and simple recurrence relation satisfied by the matching polynomials, we derive a recursion for the marginal probabilities extending the one obtained in [25] for trees and allowing us to compute explicitly the marginal at any nodes as was already noticed by Godsil in [20]. This analysis allows us to define the monomer-dimer model on infinite graphs, see Theorem 6 and then to obtain the limit of the matching number as a function of the marginal probabilities, see Theorem 10. We should stress that the analysis of the marginal probabilities is essential for our second main result which gives an explicit computation of the matching number when the local weak limit concentrates on Galton-Watson trees. Also simple adaptations of the argument in [21] would yield to a result like Theorem 1, the limit would be given in an implicit way which would not be sufficient to get our second main result.

As many other classical graph sequences, Erdős-Rényi graphs and random graphs with a prescribed degree profile admit almost surely a particularly simple random weak limit, namely a *unimodular Galton-Watson (UGW) tree* (see Example 1.1 in [1]). This random rooted tree is parametrized by a probability distribution  $\pi \in \mathcal{P}(\mathbb{N})$  with finite mean, called its *degree distribution*. It is obtained by a Galton-Watson branching process where the root has offspring distribution  $\pi$  and all other genitors have offspring distribution  $\hat{\pi} \in \mathcal{P}(\mathbb{N})$  defined by

$$\forall n \in \mathbb{N}, \hat{\pi}_n = \frac{(n+1)\pi_{n+1}}{\sum_k k\pi_k}.$$

Thanks to the markovian nature of the branching process, the recursion defining  $\gamma$  simplifies into a *recursive distributional equation*, which has been explicitly solved by the authors in a different context [12].

**Theorem 2** *With the notation of Theorem 1, if the random weak limit  $\rho$  is a UGW tree with degree distribution  $\pi$ , we have the explicit formula*

$$\gamma = \frac{1 - \max_{t \in [0,1]} F(t)}{2},$$

where

$$F(t) = t\phi'(1-t) + \phi(1-t) + \phi\left(1 - \frac{\phi'(1-t)}{\phi'(1)}\right) - 1,$$

and  $\phi(t) = \sum_n \pi_n t^n$  is the moment generating function of the degree distribution  $\pi$ .

Differentiating the above expression, we see that any  $t$  achieving the maximum must satisfy

$$\phi'(1)t = \phi'\left(1 - \frac{\phi'(1-t)}{\phi'(1)}\right). \quad (5)$$

For Erdős-Rényi random graphs with connectivity  $c$ , the degree of the limiting UGW tree is Poisson with parameter  $c$  (i.e.  $\phi(t) = \exp(ct - c)$ ), so that (5) becomes  $t = e^{-ce^{-ct}}$ . We thus

recover precisely Karp and Sipser's formula (2). Similarly, for random graphs with a prescribed degree sequence, the log-concave assumption made by Bohmann and Frieze guarantees that the above minimum is achieved at  $t = 0$ , hence (3) follows automatically.

A classical area of combinatorial optimization is formed by bipartite matching [23]. We end this section, with a specialization of our results to bipartite graphs  $G = (V = V^a \cup V^b, E)$ . The natural limit for a sequence of bipartite graphs is the following hierarchal Galton-Watson tree parameterized by two distributions on  $\mathbb{N}$  with finite first moment,  $\pi^a$  and  $\pi^b$  and a parameter  $\lambda \in [0, 1]$ . We denote  $\hat{\pi}^a$  and  $\hat{\pi}^b$  the corresponding distributions given by the transformation (2). We also denote  $\phi^a$  and  $\phi^b$  the generating functions of  $\pi^a$  and  $\pi^b$ . The hierarchal Galton-Watson tree is then defined as follows: with probability  $\lambda$ , the root has offspring distribution  $\pi^a$ , all odd generation genitors have offspring distribution  $\hat{\pi}^b$  and all even generation genitors have offspring distribution  $\hat{\pi}^a$ ; similarly with probability  $1 - \lambda$ , the root has offspring distribution  $\pi^b$ , all odd generation genitors have offspring distribution  $\hat{\pi}^a$  and all even generation genitors have offspring distribution  $\hat{\pi}^b$ . In the first (resp. second) case, we say that the root and all even generations are of type  $a$  (resp.  $b$ ) and all the odd generations are of type  $b$  (resp.  $a$ ). To get a *unimodular hierarchal Galton-Watson (UHGW) tree* with degree distributions  $\pi^a$  and  $\pi^b$ , we need to have:  $\lambda\phi^{a'}(1) = (1 - \lambda)\phi^{b'}(1)$ , so that  $\lambda = \frac{\phi^{b'}(1)}{\phi^{a'}(1) + \phi^{b'}(1)}$ .

**Theorem 3** *With the notation of Theorem 1, assume that the random weak limit  $\rho$  is a UHGW tree with degree distributions  $\pi^a, \pi^b$ . If  $\pi^a$  and  $\pi^b$  have finite first moment, then*

$$\gamma = \frac{\phi^{b'}(1)}{\phi^{a'}(1) + \phi^{b'}(1)} (1 - \max_{t \in [0,1]} F^a(t)), \quad (6)$$

where  $F^a$  is defined by:

$$F^a(t) = \phi^a \left( 1 - \frac{\phi^{b'}(1-t)}{\phi^{b'}(1)} \right) - \frac{\phi^{a'}(1)}{\phi^{b'}(1)} \left( 1 - \phi^b(1-t) - t\phi^{b'}(1-t) \right).$$

Note that if  $\phi^a(x) = \phi^b(x)$ , we find the same limit as in Theorem 2. Note also that our Theorem 3 computes the independence number of random bipartite graphs. Recall that a set of points in a graph  $G$  is said to be independent if no two of them are adjacent. The cardinality of any largest independent set of points in  $G$  is known as the independence number of  $G$  or the stability number of  $G$  and is denoted by  $\alpha(G)$ . By König's theorem, we know that for any bipartite graph  $G$  with vertex set  $V$ ,  $\alpha(G) + \nu(G) = |V|$ . The fact that a limit for  $\frac{\alpha(G_n)}{|V_n|}$  exists, has been proved recently in [5] for Erdős-Rényi and random regular graphs. The actual value for this limit is unknown except for Erdős-Rényi graphs with mean degree  $c < e$ . In this case, the leaf-removal algorithm allows to compute explicitly the limit which agrees with (6) with  $\phi^a(x) = \phi^b(x) = \exp(cx - x)$ .

Motivated by some applications for Cuckoo Hashing [18], [15], recent results have been obtained in the particular case where  $\pi^a(k) = 1$  for some  $k \geq 3$  and  $\pi^b$  is a Poisson distribution with parameter  $\alpha k$ . These degree distributions arise if one consider a sequence of bipartite graphs with  $\lfloor \alpha m \rfloor$  nodes of type  $a$  (called the items),  $m$  nodes of type  $b$  (called the locations) and each node of type  $a$  is connected with  $k$  nodes of type  $b$  chosen uniformly at random (corresponding to the assigned locations the item can be stored in). The most complete result in this domain has been obtained in [19] and follows (see Section 5.3) from our Theorem 3, namely:

**Corollary 4** *Under the assumption of Theorem 3 and with  $\pi^a(k) = 1$  for some  $k \geq 3$  and  $\pi^b$  is a Poisson distribution with parameter  $\alpha k$ . Let  $\xi$  be the unique solution of the equation:*

$$k = \frac{\xi(1 - e^{-\xi})}{1 - e^{-\xi} - \xi e^{-\xi}},$$

and  $\alpha_c = \frac{\xi}{k(1 - e^{-\xi})^{k-1}}$ .

- for  $\alpha \leq \alpha_c$ , all (except  $o_p(n)$ ) vertices of type  $a$  are covered, i.e.  $\frac{\nu(G_n)}{|V_n^a|} \xrightarrow{n \rightarrow \infty} 1$ .
- for  $\alpha > \alpha_c$ , we have:

$$\frac{\nu(G_n)}{|V_n^a|} \xrightarrow{n \rightarrow \infty} 1 - \frac{1}{\alpha} \left( e^{-\xi^*} + \xi^* e^{-\xi^*} + \frac{\xi^*}{k} (1 - e^{-\xi^*}) - 1 \right), \quad (7)$$

where  $\xi^* = k\alpha x^*$  and  $x^*$  is the largest solution of  $x = (1 - e^{-k\alpha x})^{k-1}$ .

In words,  $\alpha_c$  is the load threshold: if  $\alpha \leq \alpha_c$ , there is an assignment of the  $\lfloor \alpha m \rfloor$  items to a table with  $m$  locations that respects the choices of all items, whereas for  $\alpha > \alpha_c$ , such an assignment does not exist and (7) gives the maximal number of items assigned without collision. Note that results in [18], [15] are slightly different in the sense that for the specific sequence of random graphs described above, they show that for  $\alpha \leq \alpha_c$  all vertices of type  $a$  are covered with high probability. There is no hope to obtain such a result for any sequence of graphs converging locally to the UHGW tree defined in our Corollary 4.

### 3 The Monomer-Dimer model on infinite graphs with bounded degree

We start with a finite graph  $G$ . Consider a natural family of probability distributions on the set of matchings  $\mathbb{M}(G)$ , parameterized by a single parameter  $z > 0$  called the *temperature* (note that the standard temperature  $T$  in physics would correspond to  $z = e^{-1/T}$  but this will not be important here): for any  $M \in \mathbb{M}(G)$ ,

$$\mu_G^z(M) = \frac{z^{|V|-2|M|}}{P_G(z)}, \quad (8)$$

where  $P_G$  is the *matching polynomial*,  $P_G(z) = \sum_{M \in \mathbb{M}(G)} z^{|V|-2|M|}$ . In statistical physics, this is called the *monomer-dimer model* at temperature  $z$  on  $G$  (see [21] for a complete treatment). We let  $\mathcal{M}_G^z$  denote a random element of  $\mathbb{M}(G)$  with law  $\mu_G^z$ , and we call it a *Boltzmann random matching at temperature  $z$  on  $G$* . Note that the lowest degree coefficient of  $P_G$  is precisely the number of largest matchings on  $G$ . Therefore,  $\mathcal{M}_G^z$  converges in law to a uniform largest matching as the temperature  $z$  tends to zero. We define the *root-exposure probability (REP)* of the rooted graph  $G$  as

$$\mu_G^z(\circ \text{ is exposed}) = \mathcal{R}_z[G, \circ] = \frac{zP_{G-\circ}(z)}{P_G(z)}, \quad (9)$$

where  $G - \circ$  is the graph obtained from  $G$  by removing its root  $\circ$ . Since the matchings of  $G$  that expose  $\circ$  are exactly the matchings of  $G - \circ$ , we have the identity

$$\mathcal{R}_z[G, \circ] = \frac{zP_{G-\circ}(z)}{P_G(z)}, \quad (10)$$

which already shows that the REP is an analytic function of the fugacity. The remarkable fact that its domain of analyticity contains the right complex half-plane

$$\mathbb{H}_+ = \{z \in \mathbb{C}; \Re(z) > 0\}$$

is a consequence of the powerful Heilmann-Lieb theorem [21] (see [14] for generalizations). The key to the study of the REP is the following elementary but fundamental local recursion :

$$\mathcal{R}_z[G, \circ] = z^2 \left( z^2 + \sum_{v \sim \circ} \mathcal{R}_z[G - \circ, v] \right)^{-1}. \quad (11)$$

Clearly, this recursion determines uniquely the functional  $\mathcal{R}_z$  on the class of finite rooted graphs, and may thus be viewed as an inductive definition of the REP. Remarkably enough, this alternative characterization allows for a continuous extension to infinite graphs with bounded degree, even though the above recursion never ends. We let  $\mathcal{H}$  denote the space of analytic functions on  $\mathbb{H}_+$ , equipped with its usual topology of uniform convergence on compact sets. Our fundamental lemma is as follows :

**Theorem 5 (The fundamental local lemma)**

1. For every fixed  $z \in \mathbb{H}_+$ , the local recursion (11) determines a unique  $\mathcal{R}_z: \mathcal{G}_d \rightarrow z^{-1}\mathbb{H}_+$ .
2. For every fixed  $[G, \circ] \in \mathcal{G}_d$ ,  $z \mapsto \mathcal{R}_z[G, \circ]$  is analytic.
3. The resulting mapping  $[G, \circ] \in \mathcal{G}_d \mapsto \mathcal{R}_{(\cdot)}[G, \circ] \in \mathcal{H}$  is continuous.

This local lemma has strong implications for the monomer-dimer model, which we now list. The first one is the existence of an infinite volume limit for the Gibbs-Boltzmann distribution.

**Theorem 6 (Monomer-dimer model on infinite graphs)** Consider a graph  $G \in \mathcal{G}_d$  and a fugacity  $z > 0$ . As  $M$  ranges over all finite matchings of  $G$ , the cylinder-event marginals

$$\mu_G^z(M \subseteq \mathcal{M}) = z^{-2|M|} \prod_{k=1}^{2|M|} \mathcal{R}_z[G - \{v_1, \dots, v_{k-1}\}, v_k],$$

are consistent and independent of the ordering  $v_1, \dots, v_{2|M|}$  of the vertices spanned by  $M$ . They thus determine a unique probability distribution  $\mu_G^z$  over the matchings of  $G$ . It coincides with the former definition in the case where  $G$  is finite, and extends it continuously in the following sense : for any  $\circ \in V$  and any sequence  $([G_n, \circ_n])_{n \in \mathbb{N}} \in \mathcal{G}_d^{\mathbb{N}}$  converging to  $[G, \circ]$ ,

$$[G_n, \circ_n, \mathcal{M}_n] \xrightarrow[n \rightarrow \infty]{d} [G, \circ, \mathcal{M}],$$

in the local weak sense for random networks, where  $\mathcal{M}_n$  has law  $\mu_{G_n}^z$  and  $\mathcal{M}$  has law  $\mu_G^z$ .

Although it is not our concern here, we obtain as a by-product the strong convergence of the logarithm of the matching polynomial, also called *free entropy* in the monomer-dimer model :

**Corollary 7** *Let  $(G_n)_{n \in \mathbb{N}}$  be a sequence of finite graphs with bounded degree admitting a random weak limit  $[G, \circ]$ . The following convergence holds in the analytic sense on  $\mathbb{H}_+ = \{z \in \mathbb{C}; \Re(z) > 0\}$ .*

$$\frac{1}{|V_n|} \log \frac{P_{G_n}(z)}{P_{G_n}(1)} \xrightarrow{n \rightarrow \infty} \int_1^z \frac{\rho[\mathcal{R}_s[G, \circ]]}{s} ds,$$

where  $\rho[\mathcal{R}_s[G, \circ]]$  denotes the expectation under the measure  $\rho$  of the variable  $\mathcal{R}_s[G, \circ]$ .

A similar result was established in [21] for the lattice case, and in [6] under a restrictive large girth assumption.

### 3.1 Proof of Theorem 5 : the fundamental lemma

The local recursion (11) involves mappings of the form :

$$\phi_{z,d}: (x_1, \dots, x_d) \mapsto z^2 \left( z^2 + \sum_{i=1}^d x_i \right)^{-1},$$

where  $d \in \mathbb{N}$ . In the following lemma, we gather a few elementary properties of this transformation, which are immediate to check but will be of constant use throughout the paper.

**Lemma 8 (Elementary properties)** *For any  $d \in \mathbb{N}$  and  $z \in \mathbb{H}_+$ ,*

1.  $\phi_{z,d}$  maps analytically  $(z\mathbb{H}_+)^d$  into  $z\mathbb{H}_+$
2.  $|\phi_{z,d}|$  is uniformly bounded by  $|z|/\Re(z)$  on  $(z\mathbb{H}_+)^d$ .

From part 1, it follows that the REP of a finite rooted graph belongs to  $\mathcal{H}$ , when viewed as a function of the temperature  $z$ . Part 2 and Montel's theorem guarantee that the family of all those REPs is tight in  $\mathcal{H}$ . This analytic tightness can also be found in [21]. Combined with the following uniqueness property at high temperature, it will quickly lead to the proof of Theorem 5.

The local recursion (11) also involves graph transformations of the form  $[G, \circ] \mapsto [G - \circ, v]$ , where  $v \sim \circ$ . Starting from a given  $[G, \circ] \in \mathcal{G}_d$ , we let  $\text{Succ}^*[G, \circ] \subseteq \mathcal{G}_d$  denote the (denumerable) set of all rooted graphs that can be obtained by successively applying finitely many such transformations. Let

$$\mathbb{D}_{[G, \circ]} := \left\{ z \in \mathbb{C}; \Re(z) > \sqrt{d} \right\} \subseteq \mathbb{H}_+.$$

**Lemma 9 (Uniqueness at high temperature)** *Let  $[G, \circ] \in \mathcal{G}_d$  and  $z \in \mathbb{D}_{[G, \circ]}$ . If*

$$\mathcal{R}_z^1, \mathcal{R}_z^2: \text{Succ}^*[G, \circ] \rightarrow z\mathbb{H}_+$$

*both satisfy the local recursion (11) then  $\mathcal{R}_z^1 = \mathcal{R}_z^2$ .*



*Proof.* Set  $\alpha = 2|z|/\Re(z)$  and  $\beta = \Re(z)^{-2}$ . From (11) and part 2 of Lemma 8 it is clear that the absolute difference  $\Delta = |\mathcal{R}_z^1 - \mathcal{R}_z^2|$  must satisfy

$$\Delta[G, \circ] \leq \alpha \quad \text{and} \quad \Delta[G, \circ] \leq \beta \sum_{v \sim \circ} \Delta[G - \circ, v].$$

In turn, each  $\Delta[G - \circ, v]$  appearing in the second upper-bound may be further expanded into  $\beta \sum_{w \sim v, w \neq \circ} \Delta[G - \circ - v, w]$ . Iterating this procedure  $k$  times, one obtains  $\Delta[G, \circ] \leq \beta^k d^k \alpha$ . Taking the infimum over all  $k$  yields  $\Delta[G, \circ] = 0$ , since the assumption  $z \in \mathbb{D}_{[G, \circ]}$  means precisely  $\beta d < 1$ .  $\square$

*Proof of Theorem 5.* For clarity we divide the proof in three parts.

**Analytic existence.** Fix  $[G, \circ] \in \mathcal{G}_d$ , and consider an arbitrary collection of  $\mathbb{H}_+ \rightarrow z\mathbb{H}_+$  analytic functions  $z \mapsto \mathcal{R}_z^0[H, i]$ , indexed by the elements  $[H, i] \in \text{Succ}^*[G, \circ]$ . For every  $n \geq 1$ , define recursively

$$\mathcal{R}_z^n[H, i] = z^2 \left( z^2 + \sum_{j \sim i} \mathcal{R}_z^{n-1}[H - i, j] \right)^{-1}, \quad (12)$$

for all  $z \in \mathbb{H}_+$  and  $[H, i] \in \text{Succ}^*[G, \circ]$ . By Lemma 8, each sequence  $(z \mapsto \mathcal{R}_z^n[H, i])_{n \in \mathbb{N}}$  is tight in  $\mathcal{H}$ . Consequently, their joint collection as  $[H, i]$  varies in the denumerable set  $\text{Succ}^*[G, \circ]$  is sequentially tight in the product space  $\mathcal{H}^{\text{Succ}^*[G, \circ]}$ . Passing to the limit in (12), we see that any pre-limit  $\mathcal{R}_z: \text{Succ}^*[G, \circ] \rightarrow z\mathbb{H}_+$  must automatically satisfy (11) for each  $z \in \mathbb{H}_+$ . By Lemma 9, this determines uniquely the value of  $\mathcal{R}_z[G, \circ]$  for  $z$  with sufficiently large real part, and hence everywhere in  $\mathbb{H}_+$  by analyticity. To sum up, we have just proved the following : for every  $[G, \circ] \in \mathcal{G}_d$ , the limit

$$\mathcal{R}_z[G, \circ] := \lim_{n \rightarrow \infty} \mathcal{R}_z^n[G, \circ] \quad (13)$$

exists in  $\mathcal{H}$ , satisfies the recursion (11), and does not depend upon the choice of the initial condition  $\mathcal{R}_z^0: \text{Succ}^*[G, \circ] \rightarrow z\mathbb{H}_+$  (provided that the latter is analytic in  $z \in \mathbb{H}_+$ ).

**Pointwise uniqueness.** Let us now show that any  $\mathcal{S}: \text{Succ}^*[G, \circ] \rightarrow z\mathbb{H}_+$  satisfying the recursion (11) at a fixed value  $z = z_0 \in \mathbb{H}_+$  must coincide with the  $z = z_0$  specialization of the analytic solution constructed above. For each  $[H, i] \in \text{Succ}^*[G, \circ]$ , the constant initial function  $\mathcal{R}_{z_0}^0[H, i] := \mathcal{S}[H, i]$  is trivially analytic from  $\mathbb{H}_+$  to  $z\mathbb{H}_+$ , so the iteration (12) must converge to the analytic solution  $\mathcal{R}_z$ . Since  $\mathcal{R}_{z_0}^n = \mathcal{S}$  for all  $n \in \mathbb{N}$ , we obtain  $\mathcal{R}_{z_0} = \mathcal{S}$ , as desired.

**Continuity.** Finally, assume that  $([G_n, \circ])_{n \geq 1} \in \mathcal{G}_d^{\mathbb{N}}$  converges locally to  $[G, \circ]$ , and let us show that

$$\mathcal{R}_z[G_n, \circ] \xrightarrow[n \rightarrow \infty]{\mathcal{H}} \mathcal{R}_z[G, \circ]. \quad (14)$$

It is routine that, up to rooted isomorphisms,  $G, G_1, G_2, \dots$  may be represented on a common vertex set, in such a way that for each fixed  $k \in \mathbb{N}$ ,  $[G_n, \circ]_k = [G, \circ]_k$  for all  $n \geq n_k$ . By construction, any simple path  $v_1 \dots v_k$  starting from the root in  $G$  is now also a simple path starting from the root in each  $G_n, n \geq n_k$ , so the  $\mathcal{H}$ -valued sequence  $(z \mapsto \mathcal{R}_z[G_n - \{v_1, \dots, v_{k-1}\}, v_k])_{n \geq n_k}$  is well defined, and tight (Lemma 8). Again, the denumerable collection of all sequences obtained by letting the simple path  $v_1 \dots v_k$  vary in  $[G, \circ]$  is sequentially tight for the product topology, and any pre-limit must by construction satisfy (11). By pointwise uniqueness, the convergence (14) must hold.  $\square$

### 3.2 Proof of Theorem 6 : convergence of the Boltzmann distribution

Consider an infinite  $[G, \circ] \in \mathcal{G}_d$ , and let  $([G_n, \circ])_{n \geq 1}$  be a sequence of finite rooted connected graphs converging locally to  $[G, \circ]$ . As above, represent  $G, G_1, G_2, \dots$  on a common vertex set, in such a way that for each  $k \in \mathbb{N}$ ,  $[G_n, \circ]_k = [G, \circ]_k$  for all  $n \geq n_k$ . Now fix an arbitrary finite matching  $M$  in  $G$ , and denote by  $v_1, \dots, v_{2|M|}$  the vertices spanned by  $M$ , in any order. By construction,  $M$  is also a matching of  $G_n$  for large enough  $n$ . But the matchings of  $G_n$  that contain  $M$  are exactly the matchings of  $G_n - \{v_1, \dots, v_{2|M|}\}$ , and hence

$$\mu_{G_n}^z(M \subseteq \mathcal{M}) = \frac{P_{G_n - \{v_1, \dots, v_{2|M|}\}}(z)}{P_{G_n}(z)} = z^{-2M} \prod_{k=1}^{2M} \mathcal{R}_z[G_n - \{v_1, \dots, v_{k-1}\}, v_k].$$

But  $[G_n - \{v_1, \dots, v_{k-1}\}, v_k]$  converges locally to  $[G - \{v_1, \dots, v_{k-1}\}, v_k]$ , so by continuity of  $\mathcal{R}_z$ ,

$$\mu_{G_n}^z(M \subseteq \mathcal{M}) \xrightarrow{n \rightarrow \infty} z^{-2M} \prod_{k=1}^{2M} \mathcal{R}_z[G - \{v_1, \dots, v_{k-1}\}, v_k].$$

*Proof of Corollary 7.* Analytic convergence of the free entropy follows from Theorem 6 and Lebesgue dominated convergence Theorem, since for any finite graph  $G = (V, E)$  we have

$$(\log P_G)'(z) = \frac{P'_G(z)}{P_G(z)} = \frac{1}{|V|} \sum_{\circ \in V} \frac{\mathcal{R}_z[G, \circ]}{z} = \frac{\rho[\mathcal{R}_z[G, \circ]]}{z}.$$

The uniform domination  $\left| \frac{\rho[\mathcal{R}_z[G, \circ]]}{z} \right| \leq \frac{1}{\Re(z)}$  is provided by Lemma 8.  $\square$

## 4 The zero-temperature limit

Motivated by the asymptotic study of maximum matchings, we now let the temperature  $z \rightarrow 0$ . We first use the results from previous section to prove a version of Theorem 1 for graphs with bounded degree.

**Theorem 10 (The zero temperature limit in graphs with bounded degree)** *For any  $[G, \circ] \in \mathcal{G}_d$ , the zero temperature limit*

$$\mathcal{R}_*[G, \circ] = \lim_{z \rightarrow 0} \downarrow \mathcal{R}_z[G, \circ]$$

*exists. Moreover,  $\mathcal{R}_*: \mathcal{G}_d \rightarrow [0, 1]$  is the largest solution to the recursion*

$$\mathcal{R}_*[G, \circ] = \left( 1 + \sum_{v \sim \circ} \left( \sum_{w \sim v} \mathcal{R}_*[G - \circ - v, w] \right)^{-1} \right)^{-1}, \quad (15)$$

*with the conventions  $0^{-1} = \infty$ ,  $\infty^{-1} = 0$ . When  $G$  is finite,  $\mathcal{R}_*[G, \circ]$  is the probability that  $\circ$  is exposed in a uniform maximum matching.*

*Proof.*

Fix  $[G, \circ] \in \mathcal{G}_d$ . First, we claim that  $z \mapsto \mathcal{R}_z[G, \circ]$  is non-decreasing on  $\mathbb{R}_+$ . Indeed, this is obvious if  $G$  is reduced to  $\circ$ , since in that case the REP is simply 1. It then inductively extends to any finite graph  $[G, \circ]$ , because iterating twice (11) gives

$$\mathcal{R}_z[G, \circ] = \left( 1 + \sum_{v \sim \circ} \left( z^2 + \sum_{w \sim v} \mathcal{R}_z[G - \circ - v, w] \right)^{-1} \right)^{-1}. \quad (16)$$

For the infinite case,  $[G, \circ]$  is the local limit of the sequence of finite truncations  $([G, \circ]_n)_{n \in \mathbb{N}}$ , so by continuity of the REP,  $\mathcal{R}_z[G, \circ] = \lim_{n \rightarrow \infty} \mathcal{R}_z[G, \circ]_n$  must be non-decreasing in  $z$  as well.

This guarantees the existence of the  $[0, 1]$ -valued limit

$$\mathcal{R}_*[G, \circ] = \lim_{z \rightarrow 0} \downarrow \mathcal{R}_z[G, \circ].$$

Moreover, taking the  $z \rightarrow 0$  limit in (16) guarantees the recursive formula (15).

Finally, consider  $\mathcal{S}_*: \text{Succ}^*[G, \circ] \rightarrow [0, 1]$  satisfying the recursion (15). Let us show by induction over  $n \in \mathbb{N}$  that for every  $[H, i] \in \text{Succ}^*[G, \circ]$  and  $z > 0$ ,

$$\mathcal{S}_*[H, i] \leq \mathcal{R}_z[H, i]_{2n}. \quad (17)$$

The statement is trivial when  $n = 0$  ( $\mathcal{R}_z[H, i]_0 = 1$ ), and is preserved from  $n$  to  $n + 1$  because

$$\begin{aligned} \mathcal{R}_z[H, i]_{2n+2} &= \left( 1 + \sum_{j \sim i} \left( z^2 + \sum_{k \sim j} \mathcal{R}_z[H - i - j, k]_{2n} \right)^{-1} \right)^{-1} \\ &\geq \left( 1 + \sum_{j \sim i} \left( \sum_{k \sim j} \mathcal{S}_*[H - i - j, k] \right)^{-1} \right)^{-1} = \mathcal{S}_*[H - i, j]. \end{aligned}$$

Letting  $n \rightarrow \infty$  and then  $z \rightarrow 0$  in (17) yields  $\mathcal{S}_* \leq \mathcal{R}_*$ , which completes the proof  $\square$

This naturally raises the following question : may the *zero temperature limit* be interchanged with the *infinite volume limit*, as suggested by the diagram below ?

$$\begin{array}{ccc} \mathcal{R}_z[G_n, \circ_n] & \xrightarrow{n \rightarrow \infty} & \mathcal{R}_z[G, \circ] \\ \downarrow z \rightarrow \infty & & \downarrow z \rightarrow \infty \\ \mathcal{R}_*[G_n, \circ_n] & \xrightarrow[n \rightarrow \infty]{} & \mathcal{R}_*[G, \circ] \end{array}$$

Unfortunately, the recursion (15) may admit several distinct solutions, and this translates as follows : in the limit of zero temperature, *correlation decay breaks for the monomer-dimer model*, in the precise sense that the functional  $\mathcal{R}_*: \mathcal{G}_d \rightarrow [0, 1]$  is no longer continuous with respect to local convergence. For example, one can easily construct an infinite rooted tree  $[T, \circ]$  with bounded degree such that

$$\lim_{n \rightarrow \infty} \downarrow \mathcal{R}_*[T, \circ]_{2n} \neq \lim_{n \rightarrow \infty} \uparrow \mathcal{R}_*[T, \circ]_{2n+1}.$$

Despite this lack of correlation decay, the interchange of limits turns out to be valid “on average”, i.e. when looking at a uniformly chosen vertex  $\circ$ .

**Theorem 11 (Convergence of the matching number in graphs with bounded degree)**

Let  $\rho$  be a probability distribution over  $\mathcal{G}_d$ . For any sequence of finite graphs  $(G_n = (V_n, E_n))_{n \in \mathbb{N}}$  satisfying  $|E_n| = O(|V_n|)$  and having  $\rho$  as random weak limit,

$$\frac{\nu(G_n)}{|V_n|} \xrightarrow{n \rightarrow \infty} \frac{1 - \rho[\mathcal{R}_*]}{2}.$$

In order to get our Theorem 1, we need to remove the bounded degree assumption. This is done below. In the case where the limit  $\rho$  is a (two-level) Galton-Watson tree, the recursion (15) simplifies into a *recursive distributional equation* (RDE). The computations for these cases are done in Section 5.

*Proof of Theorem 11.*

Let  $G = (V, E)$  be a finite graph, and set  $\rho = \mathcal{U}(G)$ . First observe the elementary identity  $\rho[\mathcal{R}_*] = 1 - \frac{2\nu(G)}{|V|}$ . The proof of Theorem 11 will easily follow from the following uniform control:

**Lemma 12 (Uniform logarithmic error)** For any  $0 < z < 1$ ,

$$\rho[\mathcal{R}_z] + \frac{|E| \log 2}{|V| \log z} \leq \rho[\mathcal{R}_*] \leq \rho[\mathcal{R}_z]. \quad (18)$$

Indeed, let  $\rho$  be a probability distribution on  $\mathcal{G}_d$ , and let  $(G_n = (V_n, E_n))_{n \in \mathbb{N}}$  be a sequence of finite graphs with  $|V_n| = O(|E_n|)$ , whose random weak limit is  $\rho$ . For each  $n \in \mathbb{N}$ , set  $\rho_n = \mathcal{U}(G_n)$ . With these notations, proving Theorem 11 amounts to establish :

$$\rho_n[\mathcal{R}_*] \xrightarrow{n \rightarrow \infty} \rho[\mathcal{R}_*]. \quad (19)$$

Since  $\rho_n \Rightarrow \rho$ , and since each  $\mathcal{R}_z, z > 0$  is continuous and bounded, we have for every  $z > 0$ ,

$$\rho_n[\mathcal{R}_z] \xrightarrow{n \rightarrow \infty} \rho[\mathcal{R}_z].$$

Thus, setting  $C = \sup_{n \in \mathbb{N}} \frac{|E_n|}{|V_n|}$  and letting  $n \rightarrow \infty$  in (18), we see that for any  $z < 1$ ,

$$\mu[\mathcal{R}_z] + C \frac{\log 2}{\log z} \leq \liminf_{n \rightarrow \infty} \mu_n[\mathcal{R}_*] \leq \limsup_{n \rightarrow \infty} \mu_n[\mathcal{R}_*] \leq \mu[\mathcal{R}_*].$$

Letting finally  $z \rightarrow 0$ , we obtain exactly (19), and it only remains to show Lemma 12.

*Proof of Lemma 12.* Fix  $0 < z < 1$ . Since  $z \mapsto \rho[\mathcal{R}_z]$  is non-decreasing, we have

$$\rho[\mathcal{R}_*] \leq \rho[\mathcal{R}_z] \leq \frac{-1}{\log z} \int_z^1 s^{-1} \rho[\mathcal{R}_s] ds.$$

Use  $\rho[\mathcal{R}_s] = \frac{sP'_G(s)}{|V|P_G(s)}$  to rewrite this as

$$\rho[\mathcal{R}_*] \leq \rho[\mathcal{R}_z] \leq \frac{1}{|V| \log z} \log \frac{P_G(z)}{P_G(1)}.$$

Now,  $P_G(1)$  is the total number of matchings and is thus clearly at most  $2^{|E|}$ , while  $P_G(z)$  is at least  $z^{|V| - 2\nu(G)}$ . These two bounds yield exactly (18).  $\square$

## The case of unbounded degree

In this section, we establish Theorem 1 in full generality, removing the restriction of bounded degree from Theorem 11. To this end, we introduce the  $d$ -truncation  $G^d$  ( $d \in \mathbb{N}$ ) of a graph  $G = (V, E)$ , obtained from  $G$  by *isolating* all vertices with degree more than  $d$ , i.e. removing any edge incident to them. This transformation is clearly continuous with respect to local convergence. Moreover, its effect on the matching number can be easily controlled :

$$\nu(G^d) \leq \nu(G) \leq \nu(G^d) + \#\{v \in V; \deg_G(v) > d\}. \quad (20)$$

Now, consider a sequence of finite graphs  $(G_n)_{n \in \mathbb{N}}$  admitting a random weak limit  $(G, \circ)$ . First, fixing  $d \in \mathbb{N}$ , we may apply Theorem 11 to the sequence  $(G_n^d)_{n \in \mathbb{N}}$  to obtain :

$$\frac{\nu(G_n^d)}{|V_n|} \xrightarrow{n \rightarrow \infty} \frac{1 - \rho_d[\mathcal{R}_*]}{2},$$

where  $\rho_d$  is the  $d$ -truncation of  $\rho$ . Second, we may rewrite (20) as

$$\left| \frac{\nu(G_n^d)}{|V_n|} - \frac{\nu(G_n)}{|V_n|} \right| \leq \mathcal{U}(G_n) [\deg(\circ) > d] = \frac{\#\{v \in V_n; \deg_{G_n}(v) > d\}}{|V_n|}.$$

Letting  $n \rightarrow \infty$ , we obtain

$$\limsup_{n \rightarrow \infty} \left| \frac{1 - \rho_d[\mathcal{R}_*]}{2} - \frac{\nu(G_n)}{|V_n|} \right| \leq \rho(\deg(\circ) > d),$$

This last line is, by an elementary application of Cauchy criterion, enough to guarantee the convergence promised by Theorem 1, i.e.

$$\frac{\nu(G_n)}{|V_n|} \xrightarrow{n \rightarrow \infty} \gamma, \quad \text{where} \quad \gamma := \lim_{d \rightarrow \infty} \frac{1 - \rho_d[\mathcal{R}_*]}{2}. \quad (21)$$

Note that because of the possible absence of correlation decay, the largest solution  $\mathcal{R}_*[G, \circ]$  is not a continuous function of  $(G, \circ) \in \mathcal{G}$ . In particular, we do not know whether it is always the case that

$$\gamma = \frac{1 - \rho[\mathcal{R}_*]}{2}, \quad (22)$$

as established in Theorem 11 for graphs with bounded degree. However, (22) holds in the particular cases where we have an explicit formula for  $\rho[\mathcal{R}_*]$  which depends continuously upon the degree distribution as will be the case in Section 5.

## 5 Computations on (hierarchal) Galton-Watson trees

### 5.1 The case of Galton-Watson trees

We now investigate the special case where the limiting random graph is a UGW tree  $T$ . Specifically, we fix a distribution  $\pi \in \mathcal{P}(\mathbb{N})$  with finite support (we will relax this assumption in the sequel) and we consider a UGW tree  $T$  with degree distribution  $\pi$  as defined in Section 2. The

random matchings  $\mathcal{M}_T^z, z \geq 0$  are perfectly well-defined, and all the previously established results for graphs with bounded degree hold almost surely. However, this is not the end of the story yet : the self-similar recursive structure of  $T$  gives to the fixed-point characterizations (11) and (15) a very special form that is worth making explicit.

Before we start, let us insist on the fact that  $\mathcal{R}_z[T]$  ( $z > 0$ ) is random : it is the quenched probability that the root is exposed at temperature  $z$ , given the random tree  $T$ . In light of Theorem 1, it becomes important to ask for its distribution. Let  $\mathcal{P}([0, 1])$  denote the space of Borel probability measures on  $[0, 1]$ . Given  $z > 0$ ,  $\nu \in \mathcal{P}(\mathbb{N})$  and  $\mu \in \mathcal{P}([0, 1])$ , we denote by  $\Theta_{\nu, z}(\mu)$  the law of the  $[0, 1]$ -valued r.v.

$$Y = \frac{z^2}{z^2 + \sum_{i=1}^{\mathcal{N}} X_i},$$

where  $\mathcal{N} \sim \nu$  and  $X_1, X_2, \dots \sim \mu$ , all of them being independent. This defines an operator  $\Theta_{\nu, z}$  on  $\mathcal{P}([0, 1])$ . The corresponding fixed point equation  $\mu = \Theta_{\nu, z}(\mu)$  belongs to the general class of *recursive distributional equations*, or RDE. Equivalently, it can be rewritten as

$$X \stackrel{d}{=} \frac{z^2}{z^2 + \sum_{i=1}^{\mathcal{N}} X_i},$$

where  $X_1, X_2, \dots$  are i.i.d. copies of the unknown random variable  $X$ . Note that the same RDE appears in the analysis of the spectrum and rank of adjacency matrices of random graphs [11], [12]. With this notations in hands, the infinite system of equations (11) defining  $\mathcal{R}_z[T]$  clearly leads to the following distributional characterization:

**Lemma 13** *For any  $z > 0$ ,  $\mathcal{R}_z[T]$  has distribution  $\Theta_{\pi, z}(\mu_z)$ , where  $\mu_z$  is the unique solution to the RDE.*

$$\mu_z = \Theta_{\hat{\pi}, z}(\mu_z).$$

The same program can be carried out in the zero temperature limit. Specifically, given  $\nu, \nu' \in \mathcal{P}(\mathbb{N})$  and  $\mu \in \mathcal{P}([0, 1])$ , we define  $\Theta_{\nu, \nu'}(\mu)$  as the law of the  $[0, 1]$ -valued r.v.

$$Y = \frac{1}{1 + \sum_{i=1}^{\mathcal{N}} \left( \sum_{j=1}^{\mathcal{N}'_i} X_{ij} \right)^{-1}}, \quad (23)$$

where  $\mathcal{N} \sim \nu$ ,  $\mathcal{N}'_i \sim \nu'$ , and  $X_{ij} \sim \mu$ , all of them being independent. This defines an operator  $\Theta_{\nu, \nu'}$  on  $\mathcal{P}([0, 1])$  whose fixed points will play a crucial role in our study. Then, Theorem 10 implies:

**Lemma 14** *The random variable  $\mathcal{R}_*[T]$  has law  $\Theta_{\pi, \hat{\pi}}(\mu_*)$ , where  $\mu_*$  is the largest solution to the RDE  $\mu_* = \Theta_{\hat{\pi}, \hat{\pi}}(\mu_*)$ .*

Recall that the mean of  $\Theta_{\pi, \hat{\pi}}(\mu_*)$  gives precisely the asymptotic size of a maximum matching for any sequence of finite random graphs whose random weak limit is  $T$  (Theorem 11). We will solve this RDE in the next section in the more general set-up of UHGW trees. Combined with Theorem 1 and a simple continuity argument to remove the bounded degree assumption, this will prove Theorem 2.

## 5.2 The case of hierarchal Galton-Watson trees

As in previous section, we first assume that both  $\pi^a$  and  $\pi^b$  have a finite support. We can define a RDE but with some care about the types  $a$  and  $b$ . The corresponding results read as follows:

**Lemma 15** *For any  $z > 0$ , conditionally on the root being of type  $b$  (resp.  $a$ ),  $\mathcal{R}_z[T]$  has distribution  $\Theta_{\pi^b, z}(\mu_z^a)$  (resp.  $\Theta_{\pi^a, z}(\mu_z^b)$ ), where  $\mu_z^a$  is the unique solution to the RDE:*

$$\mu_z^a = \Theta_{\hat{\pi}^a, z} \circ \Theta_{\hat{\pi}^b, z}(\mu_z^a),$$

and  $\mu_z^b = \Theta_{\hat{\pi}^b, z}(\mu_z^a)$ .

For  $z = 0$ : conditionally on the root being of type  $a$  (resp.  $b$ ), the random variable  $\mathcal{R}_*[T]$  has law  $\Theta_{\pi^a, \hat{\pi}^b}(\mu_*^a)$  (resp.  $\Theta_{\pi^b, \hat{\pi}^a}(\mu_*^b)$ ), where  $\mu_*^a$  is the largest solution to the RDE

$$\mu_*^a = \Theta_{\hat{\pi}^a, \hat{\pi}^b}(\mu_*^a), \quad (24)$$

and  $\mu_*^b$  is the largest solution to the RDE

$$\mu_*^b = \Theta_{\hat{\pi}^b, \hat{\pi}^a}(\mu_*^b),$$

We now analyze the RDE (24). We define:

$$F^a(x) = \phi^a(1 - \hat{\phi}^b(1 - x)) - \frac{\phi^{a'}(1)}{\phi^{b'}(1)}(1 - \phi^b(1 - x) - x\phi^{b'}(1 - x)). \quad (25)$$

Observe that

$$F^{a'}(x) = \frac{\phi^{a'}(1)}{\phi^{b'}(1)}\phi^{b''}(1 - x)(\hat{\phi}^a(1 - \hat{\phi}^b(1 - x)) - x).$$

Hence any  $x$  where  $F^a$  admits a local maximum must satisfy  $x = \hat{\phi}^a(1 - \hat{\phi}^b(1 - x))$ . We define the historical records of  $F^a$  as in previous section.

**Theorem 16** *If  $p_1 < \dots < p_r$  are the locations of the historical records of  $F^a$ , then the RDE (24) admits exactly  $r$  solutions ; moreover, these solutions can be stochastically ordered, say  $\mu_1 < \dots < \mu_r$ , and for any  $i \in \{1, \dots, r\}$ ,*

- $\mu_i(\{0\}^c) = p_i$  ;
- $\Theta_{\pi^a, \hat{\pi}^b}(\mu_i)$  has mean  $F^a(p_i)$ .

The proof of Theorem 16 relies on two lemmas.

**Lemma 17** *The operators  $\Theta_{\pi^a, \hat{\pi}^b}$  and  $\Theta_{\hat{\pi}^a, \hat{\pi}^b}$  are continuous (with respect to weak convergence) and strictly increasing (with respect to stochastic ordering) on  $\mathcal{P}([0, 1])$ .*

*Proof of Lemma 17.* It follows directly from the fact that, for any  $n \geq 0$  and any  $n_1, \dots, n_n \geq 0$ , the mapping

$$x \mapsto \frac{1}{1 + \sum_{i=1}^n \left( \sum_{j=1}^{n_i} x_{ij} \right)^{-1}}$$

is continuous and increasing from  $[0, 1]^{n_1 + \dots + n_n}$  to  $[0, 1]$ . □

**Lemma 18** For any  $\mu \in \mathcal{P}([0, 1])$ , letting  $p = \mu(\{0\}^c)$ , we have

1.  $\Theta_{\hat{\pi}^a, \hat{\pi}^b}(\mu)(\{0\}^c) = \hat{\phi}^a(1 - \hat{\phi}^b(1 - p))$
2. if  $\Theta_{\hat{\pi}^a, \hat{\pi}^b}(\mu) \leq \mu$ , then the mean of  $\Theta_{\hat{\pi}^a, \hat{\pi}^b}(\mu)$  is at least  $F^a(p)$ .
3. if  $\Theta_{\hat{\pi}^a, \hat{\pi}^b}(\mu) \geq \mu$ , then the mean of  $\Theta_{\hat{\pi}^a, \hat{\pi}^b}(\mu)$  is at most  $F^a(p)$ ;

In particular, if  $\mu$  is a fixed point of  $\Theta_{\hat{\pi}^a, \hat{\pi}^b}$ , then  $p = \hat{\phi}^a(1 - \hat{\phi}^b(1 - p))$  and  $\Theta_{\hat{\pi}^a, \hat{\pi}^b}(\mu)$  has mean  $F^a(p)$ .

*Proof of Lemma 18.* In equation (23) it is clear that  $Y > 0$  if and only if for any  $i \in \{1, \dots, \mathcal{N}\}$ , there exists  $j \in \{1, \dots, \mathcal{N}'\}$  such that  $X_{ij} > 0$ . With the notation introduced above, this rewrites:

$$\Theta_{\hat{\pi}^a, \hat{\pi}^b}(\mu)(\{0\}^c) = \hat{\phi}^a \left( 1 - \hat{\phi}^b(1 - \mu(\{0\}^c)) \right),$$

hence the first result follows.

Now let  $X \sim \mu$ ,  $Y \sim \Theta_{\hat{\pi}^a, \hat{\pi}^b}(\mu)$ ,  $\mathcal{N}^a \sim \pi^a$ ,  $\hat{\mathcal{N}}^a \sim \hat{\pi}^a$ , and let  $S, S_1, \dots$  have the distribution of the sum of a  $\hat{\pi}^b$ -distributed number of i.i.d. copies of  $X$ , all these variables being independent. Then,  $\Theta_{\hat{\pi}^a, \hat{\pi}^b}(\mu)$  has mean

$$\begin{aligned} \mathbb{E} \left[ \frac{1}{1 + \sum_{i=1}^{\mathcal{N}^a} S_i^{-1}} \right] &= \mathbb{E} \left[ \left( 1 - \frac{\sum_{i=1}^{\mathcal{N}^a} S_i^{-1}}{1 + \sum_{i=1}^{\mathcal{N}^a} S_i^{-1}} \right) \mathbf{1}_{\{\forall i=1 \dots \mathcal{N}^a, S_i > 0\}} \right] \\ &= \phi^a(1 - \hat{\phi}^b(1 - p)) - \phi^{a'}(1) \mathbb{E} \left[ \frac{S^{-1}}{S^{-1} + 1 + \sum_{i=1}^{\hat{\mathcal{N}}^a} S_i^{-1}} \mathbf{1}_{\{S > 0, \forall i=1 \dots \hat{\mathcal{N}}^a, S_i > 0\}} \right] \\ &= \phi^a(1 - \hat{\phi}^b(1 - p)) - \phi^{a'}(1) \mathbb{E} \left[ \frac{Y}{Y + S} \mathbf{1}_{\{S > 0\}} \right], \end{aligned}$$

where the second and last lines follow from (2) and  $Y \sim \Theta_{\hat{\pi}^a, \hat{\pi}^b}(\mu)$ , respectively. Now, for any  $s > 0$ ,  $x \mapsto x/(x + s)$  is increasing and hence, depending on whether  $\Theta_{\hat{\pi}^a, \hat{\pi}^b}(\mu) \geq \mu$  or  $\Theta_{\hat{\pi}^a, \hat{\pi}^b}(\mu) \leq \mu$ ,  $\Theta_{\hat{\pi}^a, \hat{\pi}^b}(\mu)$  has mean at most/least:

$$\begin{aligned} &\phi^a(1 - \hat{\phi}^b(1 - p)) - \phi^{a'}(1) \mathbb{E} \left[ \frac{X}{X + S} \mathbf{1}_{\{S > 0\}} \right] \\ &= \phi^a(1 - \hat{\phi}^b(1 - p)) - p \phi^{a'}(1) \mathbb{E} \left[ \frac{1}{1 + \mathcal{N}^*} \mathbf{1}_{\{\mathcal{N}^* \geq 1\}} \right] \text{ with } \mathcal{N}^* = \sum_{i=1}^{\hat{\mathcal{N}}^b} \mathbf{1}_{\{X_i > 0\}}, \end{aligned}$$

where  $\hat{\mathcal{N}}^b \sim \hat{\pi}^b$ . But using the definition (2) and the combinatorial identity  $(n + 1) \binom{n}{d} = (d + 1) \binom{n+1}{d+1}$ , one easily derive :

$$\begin{aligned} &\phi^a(1 - \hat{\phi}^b(1 - p)) - p \phi^{a'}(1) \mathbb{E} \left[ \frac{1}{1 + \mathcal{N}^*} \mathbf{1}_{\{\mathcal{N}^* \geq 1\}} \right] \\ &= \phi^a(1 - \hat{\phi}^b(1 - p)) - p \phi^{a'}(1) \sum_{n \geq 1} \hat{\pi}_n^b \sum_{d=1}^n \binom{n}{d} \frac{p^d (1 - p)^{n-d}}{d + 1} = F^a(p). \end{aligned}$$

□



*Proof of Theorem 16.* Let  $p \in [0, 1]$  such that  $\widehat{\phi}^a(1 - \widehat{\phi}^b(1 - p)) = p$ , and define  $\mu_0 = \text{Bernoulli}(p)$ . From lemma 18 we know that  $\Theta_{\widehat{\pi}^a, \widehat{\pi}^b}(\mu_0)(\{0\}^c) = p$ , and since  $\text{Bernoulli}(p)$  is the largest element of  $\mathcal{P}([0, 1])$  putting mass  $p$  on  $\{0\}^c$ , we have  $\Theta_{\widehat{\pi}^a, \widehat{\pi}^b}(\mu_0) \leq \mu_0$ . Immediately, Lemma 17 guarantees that the limit

$$\mu_\infty = \lim_{k \rightarrow \infty} \searrow \Theta_{\widehat{\pi}^a, \widehat{\pi}^b}^k(\mu_0)$$

exists in  $\mathcal{P}([0, 1])$  and is a fixed point of  $\Theta_{\widehat{\pi}^a, \widehat{\pi}^b}$ . Moreover, by Fatou's lemma, the number  $p_\infty = \mu_\infty(\{0\}^c)$  must satisfy  $p_\infty \leq p$ . But then the mean of  $\Theta_{\widehat{\pi}^a, \widehat{\pi}^b}(\mu_\infty)$  must be both

- equal to  $F^a(p_\infty)$  by Lemma 18 with  $\mu = \mu_\infty$  ;
- at least  $F^a(p)$  since this holds for all  $\Theta_{\widehat{\pi}^a, \widehat{\pi}^b} \circ \Theta_{\widehat{\pi}^a, \widehat{\pi}^b}^k(\mu_0)$ ,  $k \geq 1$  (Lemma 18 with  $\mu = \Theta_{\widehat{\pi}^a, \widehat{\pi}^b}^k(\mu_0)$ ).

We have just shown both  $F^a(p) \leq F^a(p_\infty)$  and  $p_\infty \leq p$ . From this, we now deduce the one-to-one correspondence between historical records of  $F^a$  and fixed points of  $\Theta_{\widehat{\pi}^a, \widehat{\pi}^b}$ . We treat each inclusion separately:

1. If  $F^a$  admits an historical record at  $p$ , then clearly  $p_\infty = p$ , so  $\mu_\infty$  is a fixed point satisfying  $\mu_\infty(\{0\}^c) = p$ .
2. Conversely, considering a fixed point  $\mu$  with  $\mu(\{0\}^c) = p$ , we want to deduce that  $F^a$  admits an historical record at  $p$ . We first claim that  $\mu$  is the above defined limit  $\mu_\infty$ . Indeed,  $\mu \leq \text{Bernoulli}(p)$  implies  $\mu \leq \mu_\infty$  ( $\Theta_{\widehat{\pi}^a, \widehat{\pi}^b}$  is increasing), and in particular  $p \leq p_\infty$ . Therefore,  $p = p_\infty$  and  $F^a(p) = F^a(p_\infty)$ . In other words, the two ordered distributions  $\Theta_{\widehat{\pi}^a, \widehat{\pi}^b}(\mu) \leq \Theta_{\widehat{\pi}^a, \widehat{\pi}^b}(\mu_\infty)$  share the same mean, hence are equal. This ensures  $\mu = \mu_\infty$ . Now, if  $q < p$  is any historical record location, we know from part 1 that

$$\nu_\infty = \lim_{k \rightarrow \infty} \searrow \Theta_{\widehat{\pi}^a, \widehat{\pi}^b}^k(\text{Bernoulli}(q))$$

is a fixed point of  $\Theta_{\widehat{\pi}^a, \widehat{\pi}^b}$  satisfying  $\nu_\infty(\{0\}^c) = q$ . But  $q < p$ , so  $\text{Bernoulli}(q) < \text{Bernoulli}(p)$ , hence  $\nu_\infty \leq \mu_\infty$ . Moreover, this limit inequality is strict because  $\nu_\infty(\{0\}^c) = q < p = \mu_\infty(\{0\}^c)$ . Consequently,  $\Theta_{\widehat{\pi}^a, \widehat{\pi}^b}(\nu_\infty) < \Theta_{\widehat{\pi}^a, \widehat{\pi}^b}(\mu_\infty)$  and taking expectations,  $F^a(q) < F^a(p)$ . Thus,  $F^a$  admits an historical record at  $p$ .

□

We may now finish the proof of Theorem 3.

*Proof of Theorem 3 : case of bounded degrees.* We assume that  $\pi_a$  and  $\pi_b$  have bounded support. We set  $\lambda = \frac{\phi^{b'}(1)}{\phi^{a'}(1) + \phi^{b'}(1)}$ ,  $\lambda$  is the probability that the root is of type  $a$ . Theorems 11 and 16 and Lemma 15 give:

$$\frac{\nu(G_n)}{|V_n|} \xrightarrow{n \rightarrow \infty} \frac{\lambda(1 - \max_{x \in [0, 1]} F^a(x)) + (1 - \lambda)(1 - \max_{x \in [0, 1]} F^b(x))}{2}, \quad (26)$$

where  $F^a$  is defined in (25) and  $F^b$  is defined similarly by

$$F^b(x) = \phi^b(1 - \widehat{\phi}^a(1 - x)) - \frac{\phi^{b'}(1)}{\phi^{a'}(1)}(1 - \phi^a(1 - x) - x\phi^{a'}(1 - x)). \quad (27)$$

For any  $x$  which is an historical record of  $F^a$ , we define  $y = \widehat{\phi}^b(1 - x)$  so that  $\widehat{\phi}^a(1 - y) = x$ . Then we have:

$$\begin{aligned}\lambda(1 - F^a(x)) &= \lambda \left( 1 - \phi^a(1 - y) + \phi^{a'}(1) \left( \frac{1}{\phi^{b'}(1)} - \frac{\phi^b(1 - \widehat{\phi}^a(1 - y))}{\phi^{b'}(1)} - y\widehat{\phi}^a(1 - y) \right) \right) \\ &= (1 - \lambda)(1 - F^b(y)).\end{aligned}$$

By symmetry, this directly implies that  $\lambda(1 - \max_{x \in [0,1]} F^a(x)) = (1 - \lambda)(1 - \max_{x \in [0,1]} F^b(x))$  so that (26) is equivalent to (6). This proves Theorems 2 and 3 for distributions with bounded support.  $\square$

*Proof of Theorem 3 : general case.* To keep notation simple, we only prove Theorem 2. The following proof clearly extends to the case of UHGW trees. Let  $G_1, G_2, \dots$  be finite random graphs whose local weak limit is a Galton-Watson tree  $T$ , and assume that the degree distribution  $\pi$  of  $T$  (with generating function  $\phi$ ) has a finite mean :  $\phi'(1) = \sum_n n\pi_n < \infty$ . For any rooted graph  $G$  and any fixed integer  $d \geq 1$ , recall that  $G^d$  is the graph obtained from  $G$  by deleting all edges adjacent to a vertex  $v$  whenever  $\deg(v) > d$ . Hence  $T^d$  is a Galton-Watson tree whose degree distribution  $\pi^d$  is defined by

$$\forall i \geq 0, \pi_i^d = \pi_i \mathbf{1}_{i \leq d} + \mathbf{1}_{i=0} \sum_{k \geq d+1} \pi_k.$$

By Theorem 1, Equation (21) and our weaker version of Theorem 2 for distributions with bounded support,

$$\frac{\nu(G_n)}{|V_n|} \xrightarrow{n \rightarrow \infty} \lim_{d \rightarrow \infty} \min_{x \in [0,1]} g^d(x), \quad (28)$$

with  $\phi_d(x) = \sum_{k=0}^d \pi_k x^k$  and

$$g^d(x) = 1 - \frac{1}{2}(1 - x)\phi'_d(x) - \frac{1}{2}\phi_d(x) - \frac{1}{2}\phi_d \left( 1 - \frac{\phi'_d(x)}{\phi'_d(1)} \right).$$

Also, as  $d \rightarrow \infty$ , we have  $\phi_d \rightarrow \phi$  and  $\phi'_d \rightarrow \phi'$  uniformly on  $[0, 1]$ , so

$$\min_{x \in [0,1]} g^d(x) \xrightarrow{n \rightarrow \infty} \min_{x \in [0,1]} g(x), \quad (29)$$

with  $g(x) = 1 - \frac{1}{2}(1 - x)\phi'(x) - \frac{1}{2}\phi(x) - \frac{1}{2}\phi \left( 1 - \frac{\phi'(x)}{\phi'(1)} \right)$ . Finally, combining (28) and (29), we easily obtain the desired

$$\frac{\nu(G_n)}{|V_n|} \xrightarrow{n \rightarrow \infty} \min_{x \in [0,1]} g(x).$$

$\square$

### 5.3 Proof of Corollary 4

We have

$$\begin{aligned}F^a(x) &= \left(1 - e^{-k\alpha x}\right)^k - \frac{1}{\alpha} \left(1 - e^{-k\alpha x} - k\alpha x e^{-k\alpha x}\right) \\ F^{a'}(x) &= k^2 \alpha e^{-k\alpha x} \left( \left(1 - e^{-k\alpha x}\right)^{k-1} - x \right).\end{aligned}$$

One can check that

$$\begin{aligned}\max_{x \in [0,1]} F^a(x) &= \max(0, F^a(x^*)) \\ &= \mathbf{1}(\alpha > \alpha_c) F^a(x^*).\end{aligned}$$

This shows the first point and gives the formula in the second point. To conclude the proof, just notice that for such graphs for  $\alpha \leq \alpha_c$ ,  $F^a$  has a unique historical record 0 while for  $\alpha > \alpha_c$ ,  $F^b$  has a unique historical record at  $\exp(-k\alpha x^*)$ .

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## Appendix : (Absence of) correlation decay at zero temperature

In this section, we explain the notion of correlation decay for UGW trees with degree distribution  $\pi$ . This framework is simple and allows us to show why standard techniques cannot be applied even in simple situations. An exploration of the correlation decay in great generality is left for future work. In this section, the sequence  $(G_n = (V_n, E_n))_{n \in \mathbb{N}}$  is a sequence of finite graphs whose local weak limit under uniform rooting is a UGW tree  $T$  with degree distribution  $\pi$  with finite first moment. The generating function of  $\pi$  is  $\phi(x) = \sum \pi_n x^n$  and  $\hat{\phi}(x) = \phi'(x)/\phi'(1)$  is the generating function of  $\hat{\pi}_n$  defined by (2). For any vertex  $v$ ,  $T_v$  denotes the subtree consisting of  $v$  and its offspring in  $T$ , rooted at  $v$ .

For any finite graph  $G$ , the leaf removal algorithm proceeds as follows: start with the empty matching and then as long as there is a pendant edge  $e = (u, v)$  with  $u$  of degree one, add this edge to the matching and remove the edge  $e$  and all its adjacent edges from the graph. The algorithm stops when there is no more pendant edge. The graph  $G$  is thus simplified into a sub-graph with only isolated vertices, matched pairs and a so-called core denoted by  $C(G)$  with minimum degree at least 2. Let  $LR(G)$  be the number of isolated vertices produced by the leaf-removal algorithm on  $G$ . As explained in the introduction, there exists a maximum matching containing the matched pairs produced by the leaf-removal algorithm. Hence in this maximum matching, the  $LR(G)$  isolated vertices will be exposed and we get the bounds:

$$LR(G) \leq 1 - 2\nu(G) \leq LR(G) + |C(G)|.$$

It is clear that these bounds will be tight if we can prove that  $|C(G_n)|/|V_n| \rightarrow 0$ . We now relate this condition to a simple RDE. The analysis of the leaf-removal algorithm has been done in [22, Section 4] (see also [12, Proposition 15] for a more closely related framework). The idea is to analyze the leaf-removal step by step where in one step, all the pendant edges of the current graph are removed. We denote by  $G_k$  the graph obtained after  $k$  steps. We now put labels on the vertices. First, all isolated vertices of  $G$  are of type  $L$ . After  $k \geq 0$  steps, for all the pendant edges  $e = (u, v)$  of  $G_k$  with  $u$  of degree one in  $G_k$  and  $v$  of degree at least 2, we say that  $u$  is of type  $L$  (a leaf of  $G_k$ ) and  $v$  is of type  $N$  ( $v$  will be covered, i.e. not exposed). All the pendant edges  $e = (u, v)$  of  $G_k$  with both  $u$  and  $v$  of degree one in  $G_k$  are of type  $P$  (they are paired). Let  $L_k(G)$  (resp.  $N_k(G), P_k(G)$ ) denote the sets of vertices of type  $L$  (resp.  $N, P$ ) after  $k$  steps. Then the number of isolated vertices produced by the leaf-removal algorithm after  $k$  steps is given by  $LR_k(G) = |L_k(G)| - |N_k(G)|$  and we have the bounds:

$$|L_k(G)| - |N_k(G)| \leq 1 - 2\nu(G) \leq |L_k(G)| - |N_k(G)| + |V| - (|L_k(G)| + |N_k(G)| + |P_k(G)|). \quad (30)$$

Now the computations of the limits  $\lim_n \frac{|L_k(G_n)|}{|V_n|}$ ,  $\lim_n \frac{|N_k(G_n)|}{|V_n|}$  and  $\lim_n \frac{|P_k(G_n)|}{|V_n|}$  can be done thanks to a simple analysis on the limiting tree. Consider a (possibly infinite) UGW tree  $T$ . For any  $v$  children of the root  $\circ$ , let  $p_k$  (resp.  $q_k$ ) be the probability that  $v$  is of type  $L$  (resp.  $N$ ) after  $k$ -steps of the leaf-removal algorithm. By construction  $v$  is of type  $N$  after  $k$  steps if and only if one of its children is of type  $L$  after  $k$  steps, hence we have  $q_k = 1 - \hat{\phi}(1 - p_k)$ . Similarly,  $v$  is of type  $L$  after  $k$  steps if and only if all its children are of type  $N$  after  $k-1$  steps, hence  $p_k = \hat{\phi}(q_{k-1})$ . Hence for all  $k \geq 1$ , we have  $p_k = \hat{\phi}(1 - \hat{\phi}(1 - p_{k-1}))$  and  $p_0 = 0$ . Since  $x \mapsto \hat{\phi}(1 - \hat{\phi}(1 - x))$  is non-decreasing,  $p_k$  converges to  $p$ , the smallest solution to the fixed point equation  $x = \hat{\phi}(1 - \hat{\phi}(1 - x))$  and  $q_k$  converges to  $q = 1 - \hat{\phi}(1 - p)$ . A careful analysis (done in

[12, Proposition 15]) shows that

$$\begin{aligned}\mathbb{P}(\circ \in L_k) &= \phi(q_{k-1}) + (1 - q_{k-1} - p_k)\phi'(q_{k-1}), \\ \mathbb{P}(\circ \in N_k) &= 1 - \phi(1 - p_k) - p_k\phi'(q_{k-1}), \\ \mathbb{P}(\circ \in P_k) &= p_k\phi'(q_{k-1}).\end{aligned}$$

Now a simple coupling argument shows that for any finite  $k \geq 0$ ,

$$\lim_n \frac{|L_k(G_n)|}{|V_n|} \xrightarrow{n \rightarrow \infty} \mathbb{P}(\circ \in L_k), \quad \lim_n \frac{|N_k(G_n)|}{|V_n|} \xrightarrow{n \rightarrow \infty} \mathbb{P}(\circ \in N_k), \quad \lim_n \frac{|P_k(G_n)|}{|V_n|} \xrightarrow{n \rightarrow \infty} \mathbb{P}(\circ \in P_k)$$

Using (30), we obtain

$$\begin{aligned}\liminf_{n \rightarrow \infty} 1 - 2 \frac{\nu(G_n)}{|V_n|} &\geq \phi(q) + (1 - q)\phi'(q) - 1 + \phi(1 - p), \\ \limsup_{n \rightarrow \infty} 1 - 2 \frac{\nu(G_n)}{|V_n|} &\leq \phi(q) + (1 - q)\phi'(q) - 1 + \phi(1 - p) + (1 - \phi(1 - p) + \phi(q) + (1 - p - q)\phi'(q)).\end{aligned}\tag{31}$$

First consider the simple case where the fixed point equation  $x = \widehat{\phi}(1 - \widehat{\phi}(1 - x))$  has a unique solution, namely  $p = p^*$  where  $p^*$  is the unique solution to the fixed point equation  $x = \widehat{\phi}(1 - x)$ . In this case, we have  $q = 1 - p$  so that

$$\lim_{n \rightarrow \infty} 1 - 2 \frac{\nu(G_n)}{n} = 2\phi(1 - p^*) + p^*\phi'(1 - p^*) - 1 = F(p^*).\tag{32}$$

Indeed if  $p = p^*$ , then  $p^*$  is the unique maximum of the function  $F$  and (32) is in accordance with Theorem 2. In the particular case where  $\pi$  is a Poisson distribution with mean  $c$ , we have  $p = p^*$  iff  $c \leq e$ . In words, the leaf-removal algorithm leaves a core of size  $o(n)$  and produces a maximum matching on the complementary part of the core.

We now consider the case, where  $p < p^*$ . We still have a lower bound (31) on the number of exposed vertices and the analysis of the leaf-removal algorithm is not sufficient to conclude. However an analysis 'at zero temperature' is still possible in some cases. Recall that the zero temperature RDE was derived in (23) and solved in Theorem 16. Note first how this RDE is related to previous RDE. Let  $Z_v = \mathbf{1}(\exists \text{maximum matching in } T_v \text{ exposing } v)$ . From (15), we see that we obtain the following recursion:

$$\forall v \in T, \quad Z_v = \prod_{u \succ v} \left( 1 - \prod_{w \succ u} (1 - Z_w) \right).$$

This recursion corresponds to the two steps rule described above so that with previous notation:  $Z_v = \mathbf{1}(v \text{ is of type } L \text{ in } T_v)$ . The RDE associated to this recursion has been solved in previous section and the solutions are Bernoulli random variables with parameter  $p$  solution of the fixed point equation  $\widehat{\phi}(1 - \widehat{\phi}(1 - p)) = p$ . However the RDE (23) contains more information and allows to exclude some of these solutions.

By Theorem 16, we see that if the global maximum of  $F$  is its only historical record, then the RDE (23) has a unique solution. Observe that  $F'(x) = \phi''(1 - x) (\widehat{\phi}(1 - \widehat{\phi}(1 - x)) - x)$  and therefore, any  $x$  where  $F$  admits a local maximum must satisfy the fixed point equation

$x = \hat{\phi}(1 - \hat{\phi}(1 - x))$ . Note also that  $F'(0) = \pi_1 \hat{\phi}'(1) \geq 0$ , hence the first historical record is  $p_1 = p$  defined previously. To summarize, if  $F(p) = \max_{x \in [0,1]} F(x)$ , then the RDE (23) has a unique solution. In this case, we say that there is correlation decay (at zero temperature). Then, by a standard coupling argument, it follows that

$$\lim_n 1 - 2 \frac{\nu(G_n)}{n} = \max_{x \in [0,1]} F(x) = F(p) \geq F(p^*). \quad (33)$$

Again, this is in agreement with our Theorem 2 and to the best of our knowledge, all results in the literature on the matching number for such random graphs have been proved for graphs satisfying this correlation decay also they do not use the approach described here. For such cases, a direct analysis at 'zero temperature', i.e. with  $z = 0$  in our notations, is sufficient to compute the matching number.

A sufficient condition for weak correlation decay to occur is that  $\phi''$  is log-concave. This is in particular true in the Erdős-Rényi case, where  $\phi(x) = \exp(cx - 1)$  for any value of  $c > 0$ . The corresponding function  $F$  is given in Figure 1 for various values of  $c$ .

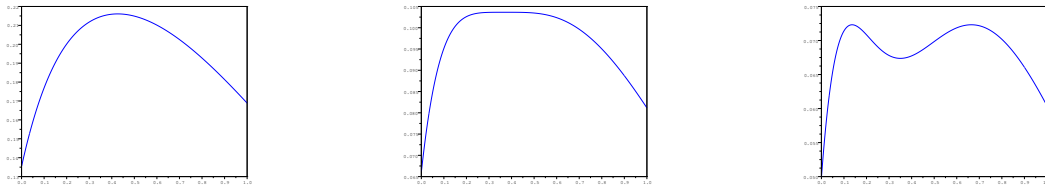


Figure 1: From left to right : plot of  $F$  for  $c = 2$ ,  $c = e$  and  $c = 3$ .

We denote by  $p' = \arg \max_{x \in [0,1]} F(x)$ . If correlation decay does not hold, then we have  $p' > p$  and the RDE (23) has several solutions. Our main result allows us to select the proper one.

We end this appendix with examples of graphs exhibiting absence of correlation decay. Note

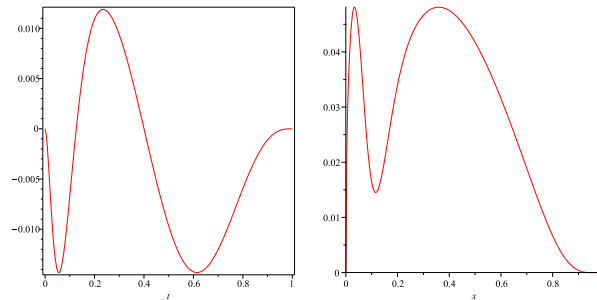


Figure 2: Plots of  $F$  for various degree distributions  $\phi(t) : \frac{3}{4}t^3 + \frac{1}{4}t^{15}$  (left) and  $\frac{50}{101}t^3 + \frac{50}{101}t^{20} + \frac{1}{101}t^{700}$  (right)

that for the degree distribution on the left of Figure 2, we have  $\lim_n \frac{\nu(G_n)}{n} = \frac{1-F(p^*)}{2}$  whereas for the oterh degree distribution, we have  $\lim_n \frac{\nu(G_n)}{n} < \frac{1-F(p^*)}{2}$ . In both cases,  $p^*$  corresponds

to the middle extremum (being a maximum on the left and a local minimum on the right). Both cases correspond to graphs with minimal degree 3 and with no almost perfect matching:  $\lim_n \frac{\nu(G_n)}{n} < 1/2$ .