

# Machine learning and convex optimization with submodular functions

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Workshop on combinatorial optimization - Cargese, 2013

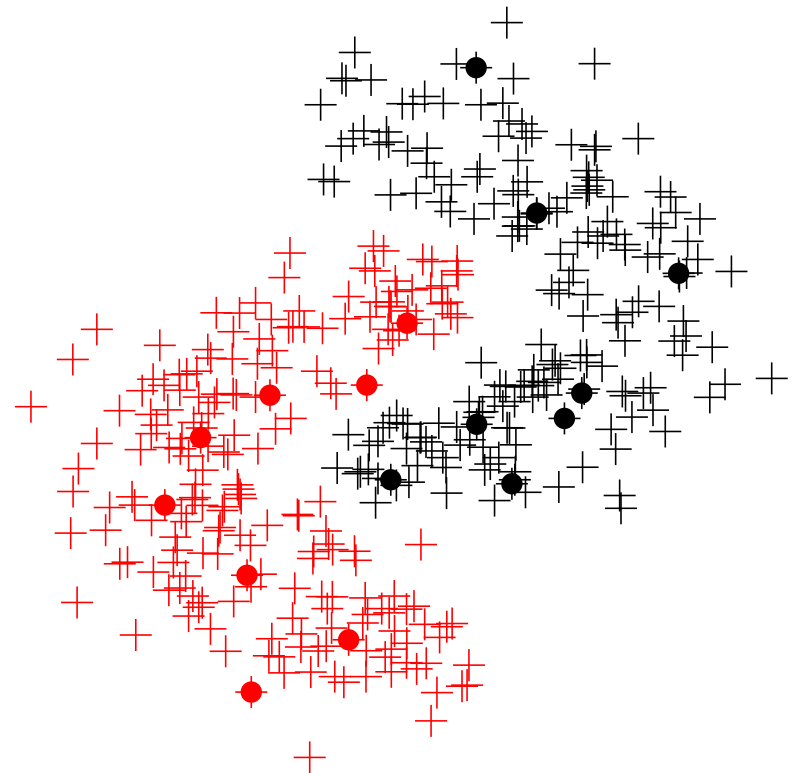
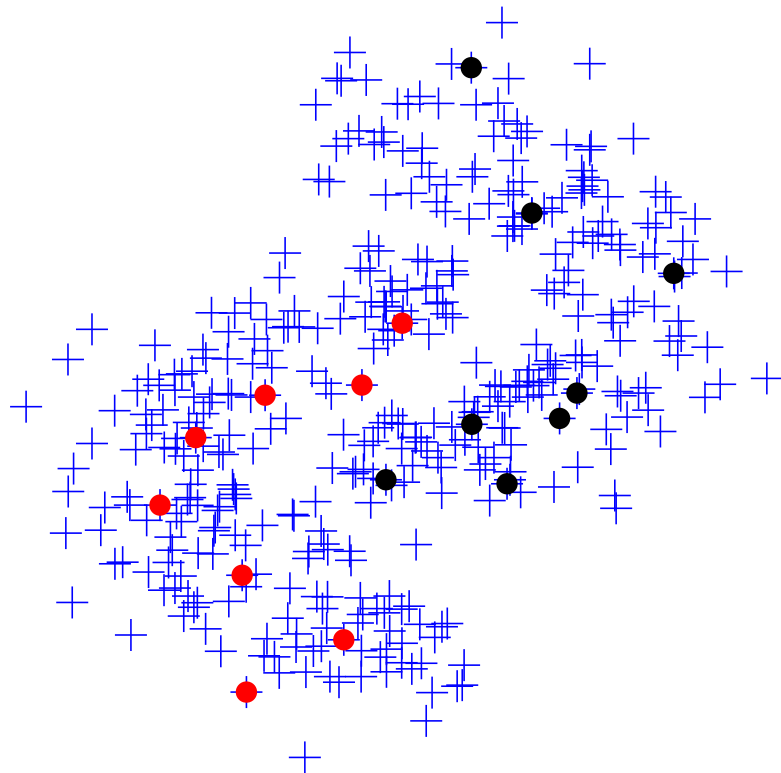
# Submodular functions - References

- **References based on combinatorial optimization**
  - *Submodular Functions and Optimization* (Fujishige, 2005)
  - *Discrete convex analysis* (Murota, 2003)
- **Tutorial paper based on convex optimization** (Bach, 2011b)
  - [www.di.ens.fr/~fbach/submodular\\_fot.pdf](http://www.di.ens.fr/~fbach/submodular_fot.pdf)
- **Slides for this lecture**
  - [www.di.ens.fr/~fbach/fbach\\_cargese\\_2013.pdf](http://www.di.ens.fr/~fbach/fbach_cargese_2013.pdf)

# Submodularity (almost) everywhere

## Clustering

- Semi-supervised clustering

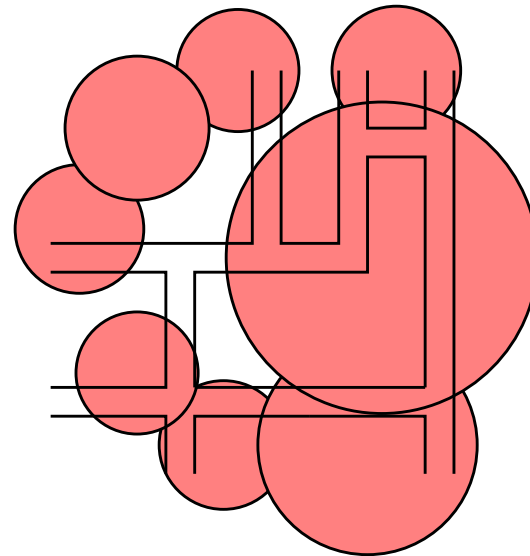
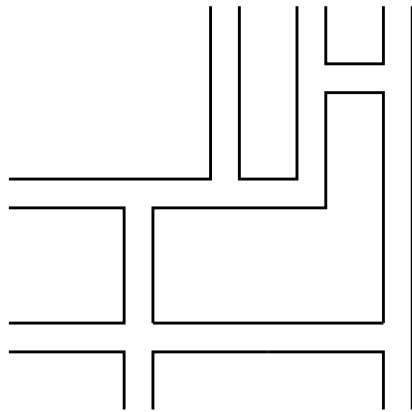


- Submodular function minimization

# Submodularity (almost) everywhere

## Sensor placement

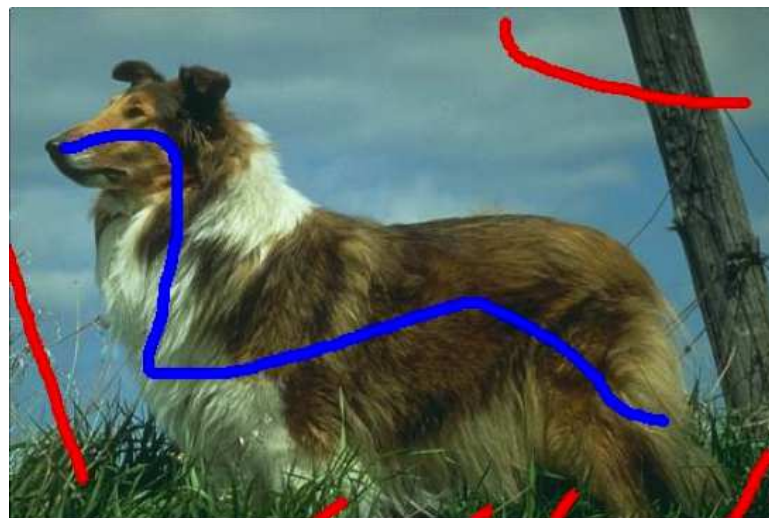
- Each sensor covers a certain area (Krause and Guestrin, 2005)
  - Goal: maximize coverage



- Submodular function maximization
- Extension to experimental design (Seeger, 2009)

# Submodularity (almost) everywhere

## Graph cuts and image segmentation

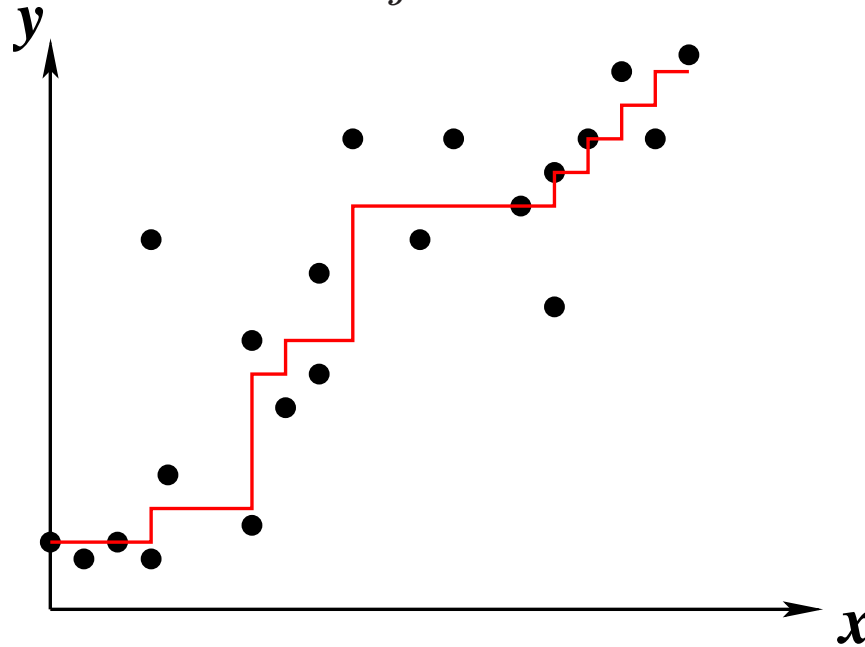


- Submodular function minimization

# Submodularity (almost) everywhere

## Isotonic regression

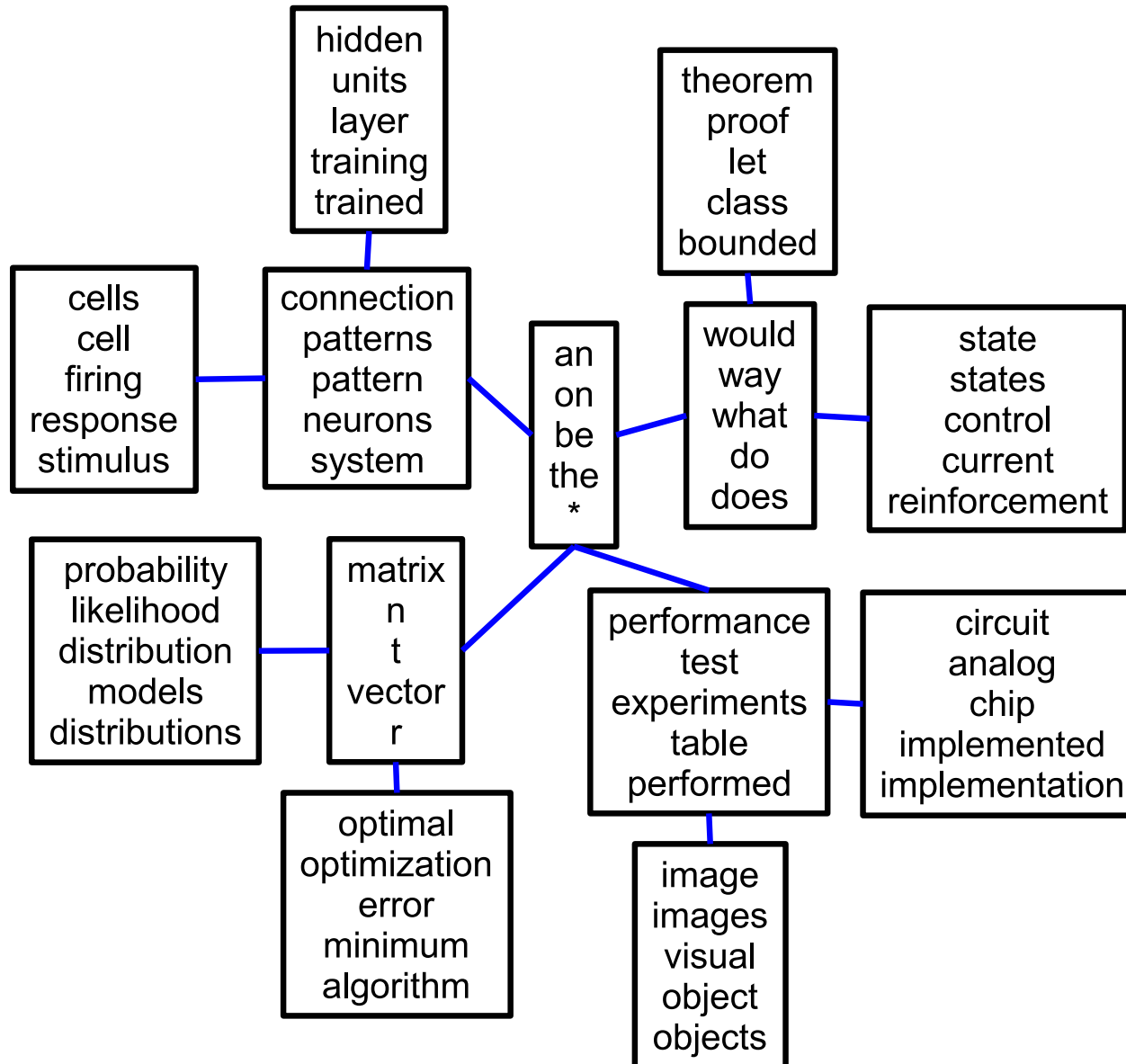
- Given real numbers  $x_i, i = 1, \dots, p$ 
  - Find  $y \in \mathbb{R}^p$  that minimizes  $\frac{1}{2} \sum_{j=1}^p (x_j - y_j)^2$  such that  $\forall i, y_i \leq y_{i+1}$



- Submodular convex optimization problem

# Submodularity (almost) everywhere

## Structured sparsity - I

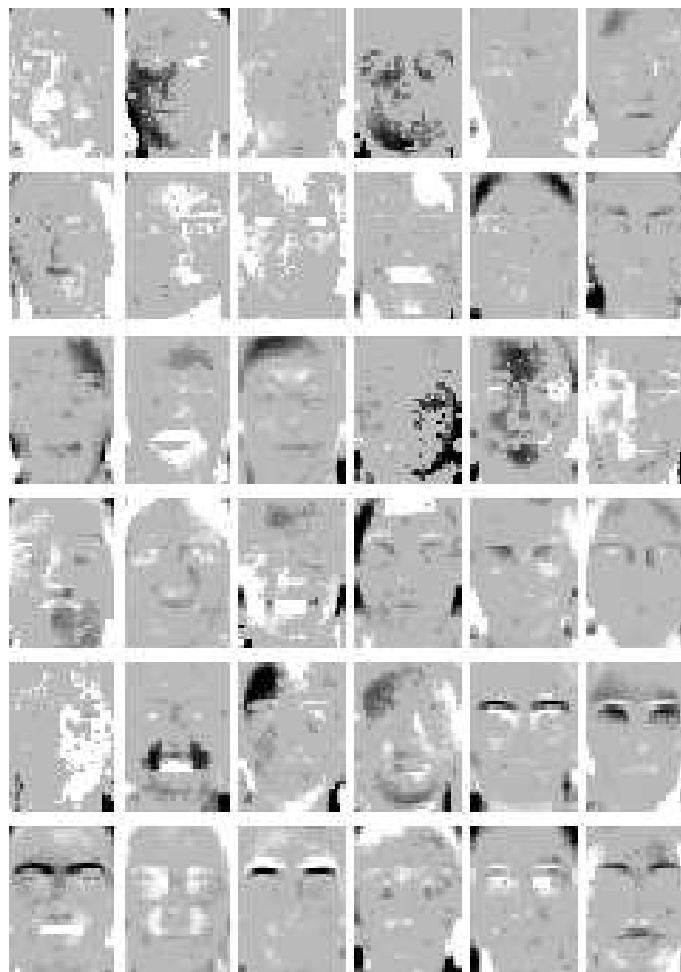


# Submodularity (almost) everywhere

## Structured sparsity - II



raw data



sparse PCA

- No structure: many zeros do not lead to better interpretability

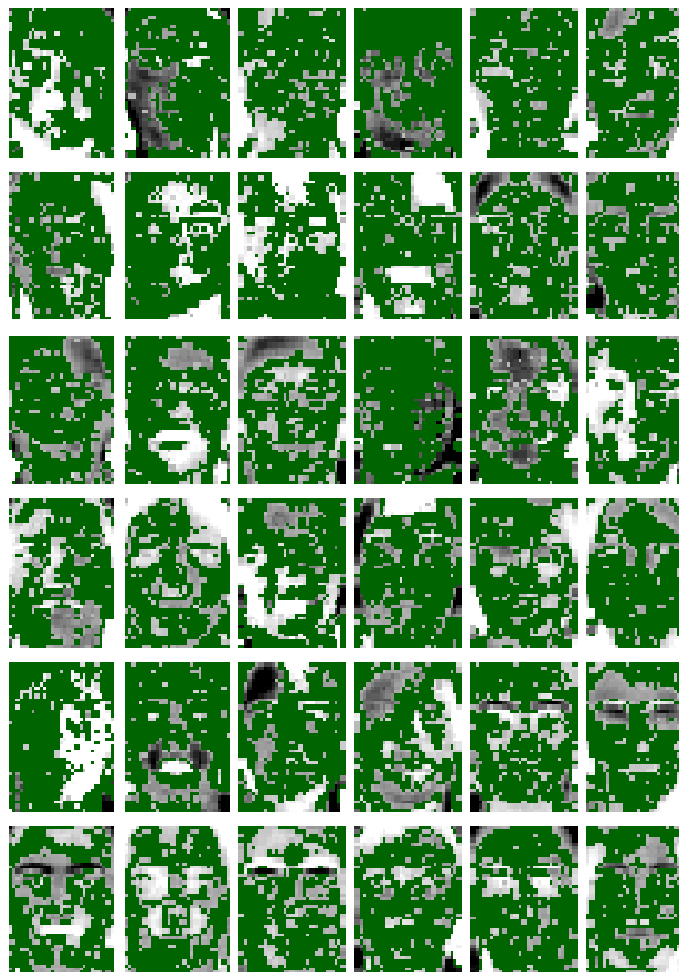


# Submodularity (almost) everywhere

## Structured sparsity - II



raw data



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# Submodularity (almost) everywhere

## Structured sparsity - II



raw data



Structured sparse PCA

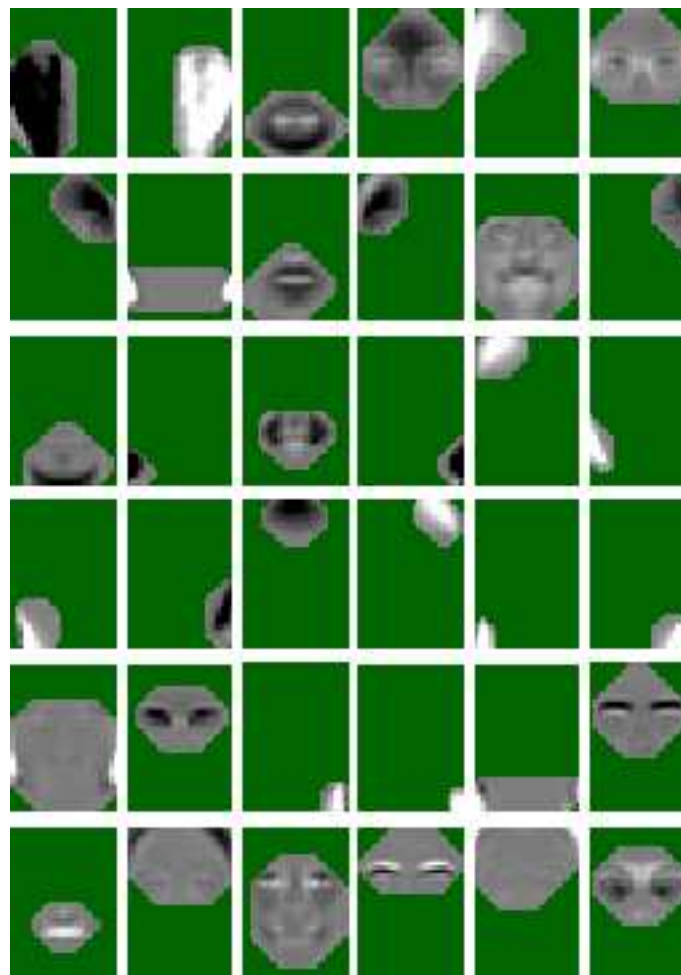
- Submodular convex optimization problem

# Submodularity (almost) everywhere

## Structured sparsity - II



raw data



Structured sparse PCA

- Submodular convex optimization problem



# Submodularity (almost) everywhere

## Image denoising

- Total variation denoising (Chambolle, 2005)

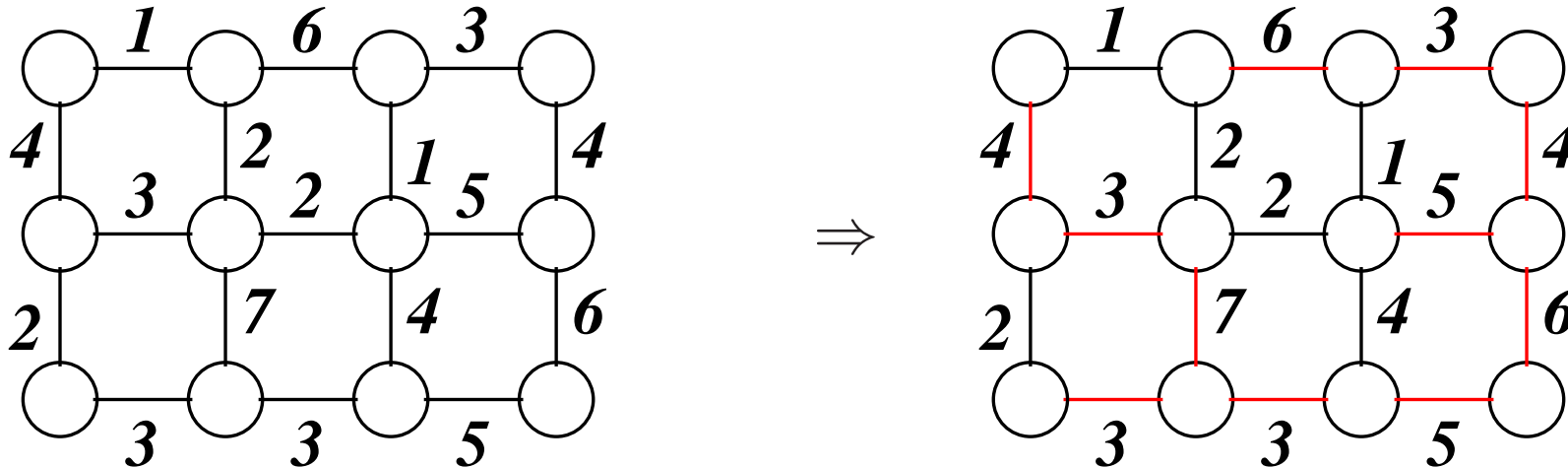


- Submodular convex optimization problem

# Submodularity (almost) everywhere

## Maximum weight spanning trees

- Given an undirected graph  $G = (V, E)$  and weights  $w : E \mapsto \mathbb{R}_+$ 
  - find the maximum weight spanning tree



- Greedy algorithm for submodular polyhedron - matroid

# Submodularity (almost) everywhere

## Combinatorial optimization problems

- Set  $V = \{1, \dots, p\}$
- Power set  $2^V =$  set of all subsets, of cardinality  $2^p$
- Minimization/maximization of a set function  $F : 2^V \rightarrow \mathbb{R}$ .

$$\min_{A \subseteq V} F(A) = \min_{A \in 2^V} F(A)$$

# Submodularity (almost) everywhere

## Combinatorial optimization problems

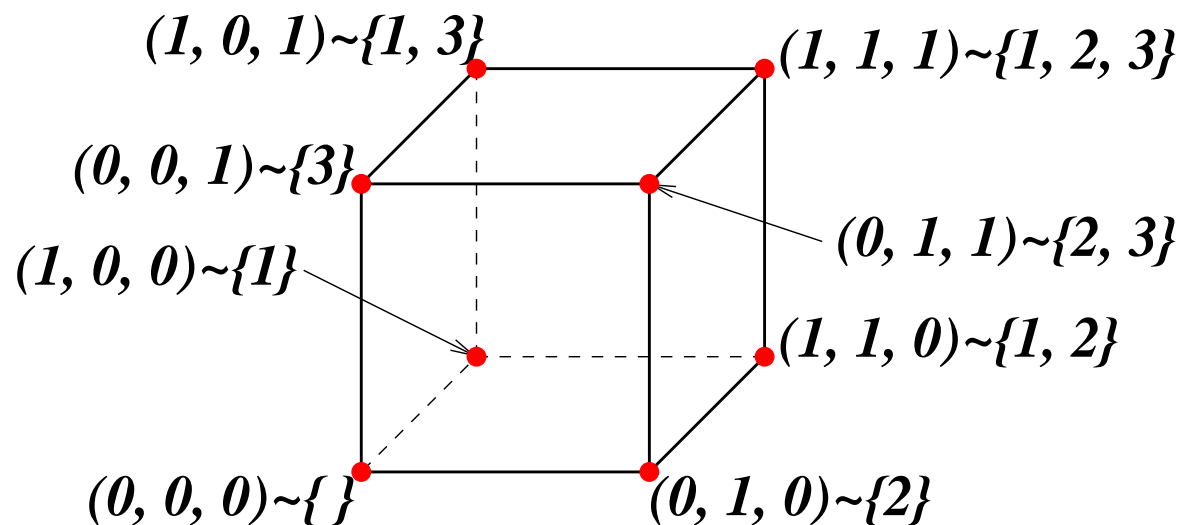
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$$\min_{A \subset V} F(A) = \min_{A \in 2^V} F(A)$$

- Reformulation as (pseudo) Boolean function

$$\min_{w \in \{0,1\}^p} f(w)$$

with  $\forall A \subset V, f(1_A) = F(A)$



# Submodularity (almost) everywhere

## Convex optimization with combinatorial structure

- **Supervised learning / signal processing**

- Minimize regularized empirical risk from data  $(x_i, y_i)$ ,  $i = 1, \dots, n$ :

$$\min_{f \in \mathcal{F}} \frac{1}{n} \sum_{i=1}^n \ell(y_i, f(x_i)) + \lambda \Omega(f)$$

- $\mathcal{F}$  is often a vector space, formulation often convex

- **Introducing discrete structures within a vector space framework**

- Trees, graphs, etc.

- Many different approaches (e.g., stochastic processes)

- **Submodularity allows the incorporation of discrete structures**



# Outline

## 1. Submodular functions

- Review and examples of submodular functions
- Links with convexity through Lovász extension

## 2. Submodular minimization

- Non-smooth convex optimization
- Parallel algorithm for special case

## 3. Structured sparsity-inducing norms

- Relaxation of the penalization of supports by submodular functions
- Extensions (symmetric,  $\ell_q$ -relaxation)

# Submodular functions

## Definitions

- **Definition:**  $F : 2^V \rightarrow \mathbb{R}$  is **submodular** if and only if

$$\forall A, B \subset V, \quad F(A) + F(B) \geq F(A \cap B) + F(A \cup B)$$

- NB: equality for *modular* functions
- Always assume  $F(\emptyset) = 0$

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- NB: equality for *modular* functions
- Always assume  $F(\emptyset) = 0$

- **Equivalent definition:**

$$\forall k \in V, \quad A \mapsto F(A \cup \{k\}) - F(A) \text{ is non-increasing}$$

$$\Leftrightarrow \forall A \subset B, \forall k \notin A, \quad F(A \cup \{k\}) - F(A) \geq F(B \cup \{k\}) - F(B)$$

- “**Concave property**”: Diminishing return property

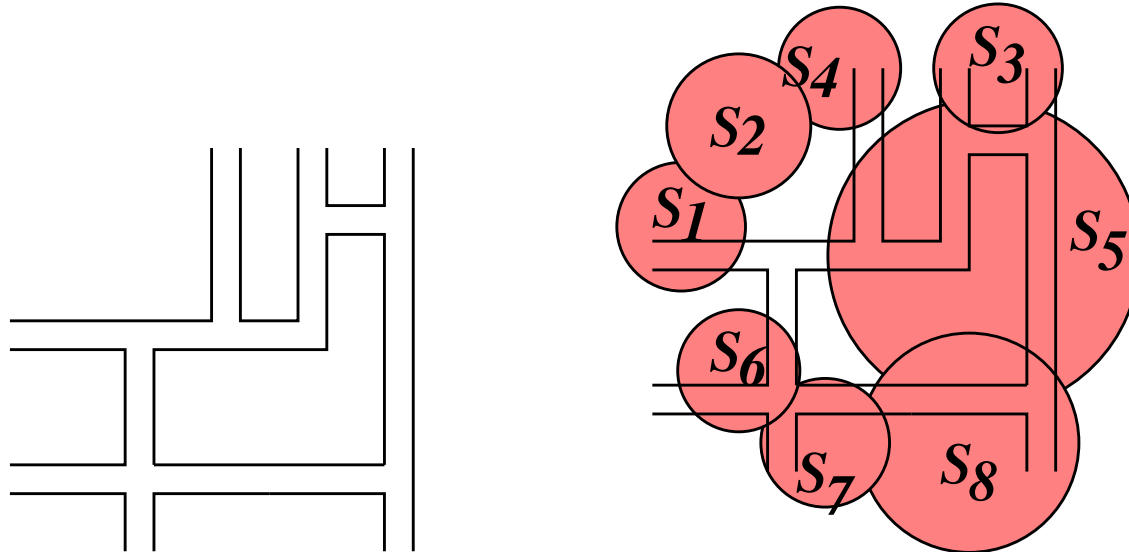
# Examples of submodular functions

## Cardinality-based functions

- Notation for modular function:  $s(A) = \sum_{k \in A} s_k$  for  $s \in \mathbb{R}^p$ 
  - If  $s = 1_V$ , then  $s(A) = |A|$  (cardinality)
- **Proposition:** If  $s \in \mathbb{R}_+^p$  and  $g : \mathbb{R}_+ \rightarrow \mathbb{R}$  is a concave function, then  $F : A \mapsto g(s(A))$  is submodular
- **Proposition 2:** If  $F : A \mapsto g(s(A))$  is submodular for all  $s \in \mathbb{R}_+^p$ , then  $g$  is concave
- Classical example:
  - $F(A) = 1$  if  $|A| > 0$  and 0 otherwise
  - May be rewritten as  $F(A) = \max_{k \in V} (1_A)_k$

# Examples of submodular functions

## Covers

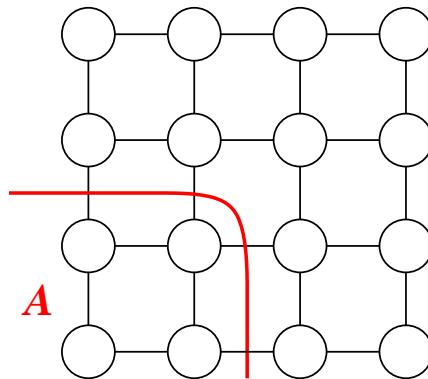


- Let  $W$  be any “base” set, and for each  $k \in V$ , a set  $S_k \subset W$
- Set cover defined as  $F(A) = \left| \bigcup_{k \in A} S_k \right|$
- *Proof of submodularity  $\Rightarrow$  homework*

# Examples of submodular functions

## Cuts

- Given a (un)directed graph, with vertex set  $V$  and edge set  $E$ 
  - $F(A)$  is the total number of edges going from  $A$  to  $V \setminus A$ .



- Generalization with  $d : V \times V \rightarrow \mathbb{R}_+$

$$F(A) = \sum_{k \in A, j \in V \setminus A} d(k, j)$$

- *Proof of submodularity  $\Rightarrow$  homework*

# Examples of submodular functions

## Entropies

- Given  $p$  random variables  $X_1, \dots, X_p$  with finite number of values
  - Define  $F(A)$  as the joint entropy of the variables  $(X_k)_{k \in A}$
  - $F$  is **submodular**
- *Proof of submodularity* using data processing inequality (Cover and Thomas, 1991): if  $A \subset B$  and  $k \notin B$ ,

$$F(A \cup \{k\}) - F(A) = H(X_A, X_k) - H(X_A) = H(X_k | X_A) \geq H(X_k | X_B)$$

- Symmetrized version  $G(A) = F(A) + F(V \setminus A) - F(V)$  is **mutual information** between  $X_A$  and  $X_{V \setminus A}$
- Extension to continuous random variables, e.g., Gaussian:  
 $F(A) = \log \det \Sigma_{AA}$ , for some positive definite matrix  $\Sigma \in \mathbb{R}^{p \times p}$

# Examples of submodular functions

## Flows

- Net-flows from multi-sink multi-source networks (Megiddo, 1974)
- See details in Fujishige (2005); Bach (2011b)
- **Efficient formulation for set covers**



# Examples of submodular functions

## Matroids

- The pair  $(V, \mathcal{I})$  is a matroid with  $\mathcal{I}$  its family of independent sets, iff:
  - (a)  $\emptyset \in \mathcal{I}$
  - (b)  $I_1 \subset I_2 \in \mathcal{I} \Rightarrow I_1 \in \mathcal{I}$
  - (c) for all  $I_1, I_2 \in \mathcal{I}$ ,  $|I_1| < |I_2| \Rightarrow \exists k \in I_2 \setminus I_1, I_1 \cup \{k\} \in \mathcal{I}$
- **Rank function** of the matroid, defined as  $F(A) = \max_{I \subset A, A \in \mathcal{I}} |I|$  is submodular (*direct proof*)
- **Graphic matroid**
  - $V$  **edge set** of a certain graph  $G = (U, V)$
  - $\mathcal{I}$  = set of subsets of edges which do not contain any cycle
  - $F(A) = |U|$  minus the number of connected components of the subgraph induced by  $A$

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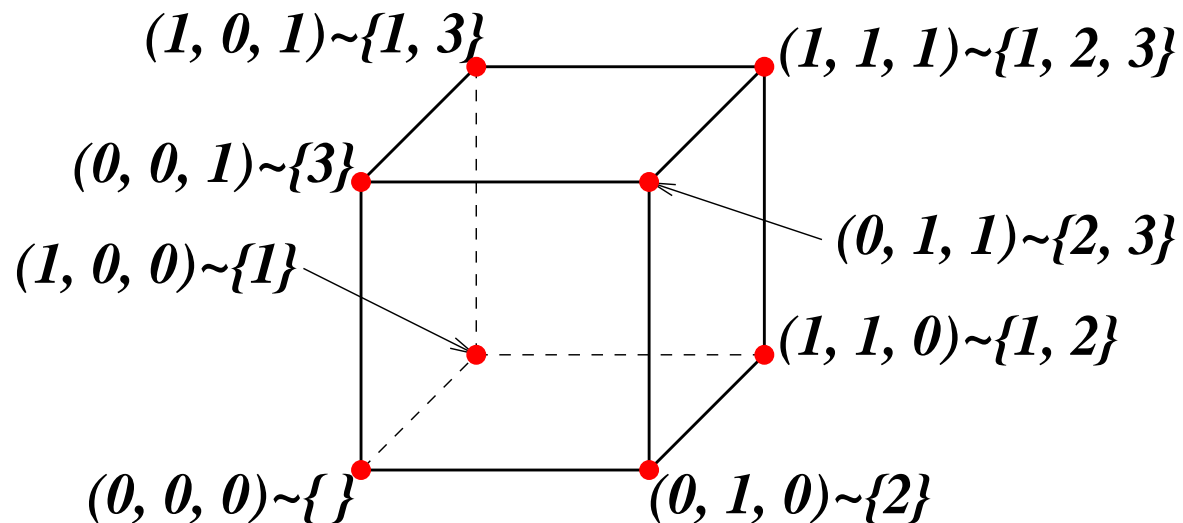
## 3. Structured sparsity-inducing norms

- Relaxation of the penalization of supports by submodular functions
- Extensions (symmetric,  $\ell_q$ -relaxation)

# Choquet integral (Choquet, 1954) - Lovász extension

- Subsets may be identified with elements of  $\{0, 1\}^p$
- Given **any** set-function  $F$  and  $w$  such that  $w_{j_1} \geq \dots \geq w_{j_p}$ , define:

$$\begin{aligned}
 f(w) &= \sum_{k=1}^p w_{j_k} [F(\{j_1, \dots, j_k\}) - F(\{j_1, \dots, j_{k-1}\})] \\
 &= \sum_{k=1}^{p-1} (w_{j_k} - w_{j_{k+1}}) F(\{j_1, \dots, j_k\}) + w_{j_p} F(\{j_1, \dots, j_p\})
 \end{aligned}$$



# Choquet integral (Choquet, 1954) - Lovász extension

## Properties

$$\begin{aligned} f(w) &= \sum_{k=1}^p w_{j_k} [F(\{j_1, \dots, j_k\}) - F(\{j_1, \dots, j_{k-1}\})] \\ &= \sum_{k=1}^{p-1} (w_{j_k} - w_{j_{k+1}}) F(\{j_1, \dots, j_k\}) + w_{j_p} F(\{j_1, \dots, j_p\}) \end{aligned}$$

- For any set-function  $F$  (even not submodular)
  - $f$  is piecewise-linear and positively homogeneous
  - If  $w = 1_A$ ,  $f(w) = F(A) \Rightarrow$  extension from  $\{0, 1\}^p$  to  $\mathbb{R}^p$

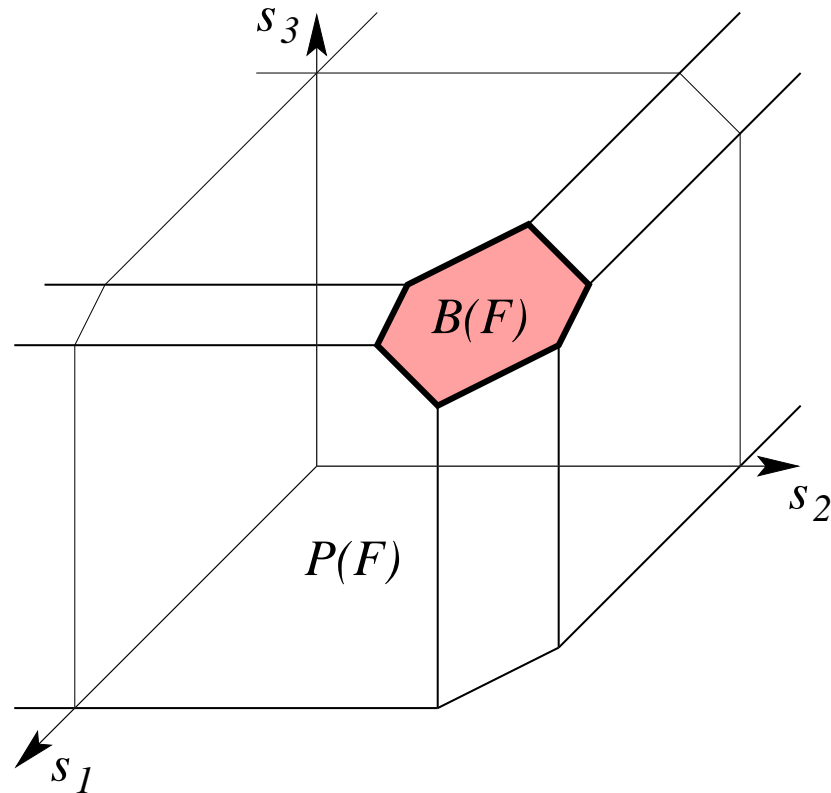
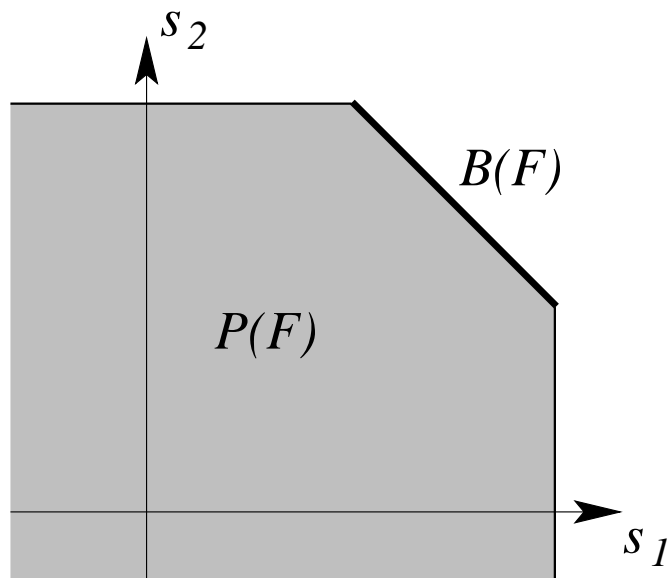
# Submodular functions

## Links with convexity (Edmonds, 1970; Lovász, 1982)

- **Theorem** (Lovász, 1982):  $F$  is submodular if and only if  $f$  is convex
- Proof requires additional notions from Edmonds (1970):
  - **Submodular and base polyhedra**

# Submodular and base polyhedra - Definitions

- Submodular polyhedron:  $P(F) = \{s \in \mathbb{R}^p, \forall A \subset V, s(A) \leq F(A)\}$
- Base polyhedron:  $B(F) = P(F) \cap \{s(V) = F(V)\}$



- Property:  $P(F)$  has non-empty interior

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- Many facets (up to  $2^p$ ), many extreme points (up to  $p!$ )

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- Base polyhedron:  $B(F) = P(F) \cap \{s(V) = F(V)\}$
- Many facets (up to  $2^p$ ), many extreme points (up to  $p!$ )
- **Fundamental property** (Edmonds, 1970): If  $F$  is submodular, maximizing linear functions may be done by a “greedy algorithm”
  - Let  $w \in \mathbb{R}_+^p$  such that  $w_{j_1} \geq \dots \geq w_{j_p}$
  - Let  $s_{j_k} = F(\{j_1, \dots, j_k\}) - F(\{j_1, \dots, j_{k-1}\})$  for  $k \in \{1, \dots, p\}$
  - Then  $f(w) = \max_{s \in P(F)} w^\top s = \max_{s \in B(F)} w^\top s$
  - Both problems attained at  $s$  defined above
- Simple proof by convex duality



# Submodular functions

## Links with convexity

- **Theorem** (Lovász, 1982): If  $F$  is submodular, then

$$\min_{A \subset V} F(A) = \min_{w \in \{0,1\}^p} f(w) = \min_{w \in [0,1]^p} f(w)$$

- Consequence: Submodular function minimization may be done in polynomial time (through ellipsoid algorithm)
- **Representation of  $f(w)$  as a support function** (Edmonds, 1970):

$$f(w) = \max_{s \in B(F)} s^\top w$$

- Maximizer  $s$  may be found efficiently through the greedy algorithm

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# Submodular function minimization

## Dual problem

- Let  $F : 2^V \rightarrow \mathbb{R}$  be a submodular function (such that  $F(\emptyset) = 0$ )
- **Convex duality** (Edmonds, 1970):

$$\begin{aligned} \min_{A \subseteq V} F(A) &= \min_{w \in [0,1]^p} f(w) \\ &= \min_{w \in [0,1]^p} \max_{s \in B(F)} w^\top s \\ &= \max_{s \in B(F)} \min_{w \in [0,1]^p} w^\top s = \max_{s \in B(F)} s_-(V) \end{aligned}$$

# Exact submodular function minimization

## Combinatorial algorithms

- Algorithms based on  $\min_{A \subset V} F(A) = \max_{s \in B(F)} s_-(V)$
- Output the subset  $A$  and a base  $s \in B(F)$  as a **certificate of optimality**
- Best algorithms have **polynomial complexity** (Schrijver, 2000; Iwata et al., 2001; Orlin, 2009) (typically  $O(p^6)$  or more)
- Update a sequence of convex combination of vertices of  $B(F)$  obtained from the greedy algorithm using a specific order:
  - **Based only on function evaluations**
- Recent algorithms using efficient reformulations in terms of generalized graph cuts (Jegelka et al., 2011)

# Approximate submodular function minimization

- For most machine learning applications, no need to obtain exact minimum
  - For convex optimization, see, e.g., Bottou and Bousquet (2008)

$$\min_{A \subset V} F(A) = \min_{w \in \{0,1\}^p} f(w) = \min_{w \in [0,1]^p} f(w)$$

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- Important properties of  $f$  for convex optimization
  - Polyhedral function
  - Representation as maximum of linear functions

$$f(w) = \max_{s \in B(F)} w^\top s$$

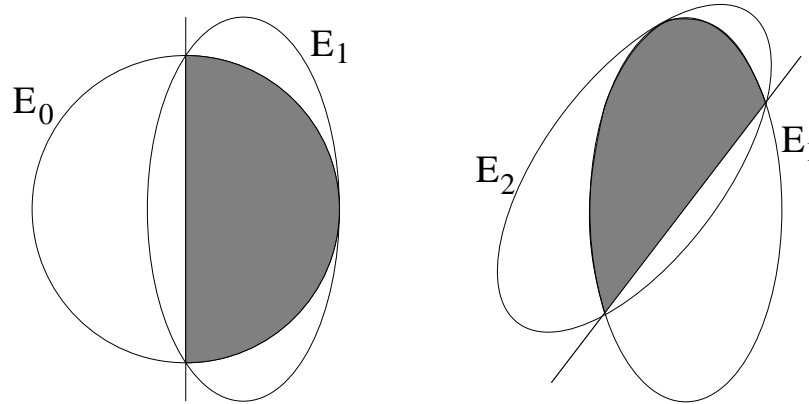
- Stability vs. speed vs. generality vs. ease of implementation

# Projected subgradient descent (Shor et al., 1985)

- Subgradient of  $f(w) = \max_{s \in B(F)} s^\top w$  through the greedy algorithm
- Using **projected subgradient descent** to minimize  $f$  on  $[0, 1]^p$ 
  - Iteration:  $w_t = \Pi_{[0,1]^p}(w_{t-1} - \frac{C}{\sqrt{t}}s_t)$  where  $s_t \in \partial f(w_{t-1})$
  - Convergence rate:  $f(w_t) - \min_{w \in [0,1]^p} f(w) \leq \frac{\sqrt{p}}{\sqrt{t}}$  with primal/dual guarantees (Nesterov, 2003)
- Fast iterations but slow convergence
  - need  $O(p/\varepsilon^2)$  iterations to reach precision  $\varepsilon$
  - need  $O(p^2/\varepsilon^2)$  function evaluations to reach precision  $\varepsilon$

# Ellipsoid method (Nemirovski and Yudin, 1983)

- Build a sequence of minimum volume ellipsoids that enclose the set of solutions



- Cost of a single iteration:  $p$  function evaluations and  $O(p^3)$  operations
- Number of iterations:  $2p^2 \left( \max_{A \subset V} F(A) - \min_{A \subset V} F(A) \right) \log \frac{1}{\varepsilon}$ .  
–  $O(p^5)$  operations and  $O(p^3)$  function evaluations
- Slow in practice (the bound is “tight”)



# Analytic center cutting planes (Goffin and Vial, 1993)

- **Center of gravity method**

- improves the convergence rate of ellipsoid method
- cannot be computed easily

- **Analytic center** of a polytope defined by  $a_i^\top w \leq b_i, i \in I$

$$\min_{w \in \mathbb{R}^p} - \sum_{i \in I} \log(b_i - a_i^\top w)$$

- **Analytic center cutting planes (ACCPM)**

- Each iteration has complexity  $O(p^2|I| + |I|^3)$  using Newton's method
- No linear convergence rate
- Good performance in practice

# Simplex method for submodular minimization

- Mentioned by Girlich and Pinaruk (1997); McCormick (2005)
- **Formulation as linear program:**  $s \in B(F) \Leftrightarrow s = S^\top \eta$ ,  $S \in \mathbb{R}^{d \times p}$

$$\begin{aligned} \max_{s \in B(F)} s_-(V) &= \max_{\eta \geq 0, \eta^\top \mathbf{1}_d = 1} \sum_{i=1}^p \min\{(S^\top \eta)_i, 0\} \\ &= \max_{\eta \geq 0, \alpha \geq 0, \beta \geq 0} -\beta^\top \mathbf{1}_p \text{ such that } S^\top \eta - \alpha + \beta = 0, \eta^\top \mathbf{1}_d = 1. \end{aligned}$$

- **Column generation for simplex methods:** only access the rows of  $S$  by maximizing linear functions
  - no complexity bound, may get global optimum if enough iterations

# Separable optimization on base polyhedron

- **Optimization of convex functions** of the form  $\Psi(w) + f(w)$  with  $f$  Lovász extension of  $F$ , and  $\Psi(w) = \sum_{k \in V} \psi_k(w_k)$
- **Structured sparsity**
  - Total variation denoising - isotonic regression
  - Regularized risk minimization penalized by the Lovász extension

# Total variation denoising (Chambolle, 2005)

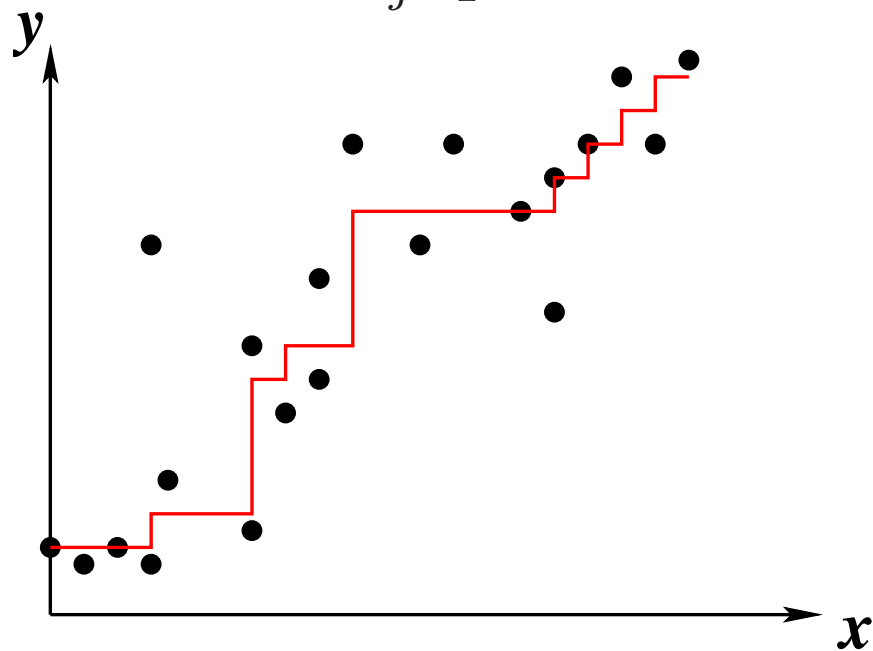
- $F(A) = \sum_{k \in A, j \in V \setminus A} d(k, j) \Rightarrow f(w) = \sum_{k, j \in V} d(k, j)(w_k - w_j)_+$
- $d$  symmetric  $\Rightarrow f =$  total variation



# Isotonic regression

- Given real numbers  $x_i, i = 1, \dots, p$

– Find  $y \in \mathbb{R}^p$  that minimizes  $\frac{1}{2} \sum_{j=1}^p (x_j - y_j)^2$  such that  $\forall i, y_i \leq y_{i+1}$



- For a directed chain,  $f(y) = 0$  if and only if  $\forall i, y_i \leq y_{i+1}$
- Minimize  $\frac{1}{2} \sum_{j=1}^p (x_j - y_j)^2 + \lambda f(y)$  for  $\lambda$  large

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- **Structured sparsity**
  - Total variation denoising - isotonic regression
  - Regularized risk minimization penalized by the Lovász extension
- **Proximal methods** (see second part)
  - Minimize  $\Psi(w) + f(w)$  for smooth  $\Psi$  as soon as the following “proximal” problem may be obtained efficiently

$$\min_{w \in \mathbb{R}^p} \frac{1}{2} \|w - z\|_2^2 + f(w) = \min_{w \in \mathbb{R}^p} \sum_{k=1}^p \frac{1}{2} (w_k - z_k)^2 + f(w)$$

- **Submodular function minimization**

# Separable optimization on base polyhedron

## Convex duality

- Let  $\psi_k : \mathbb{R} \rightarrow \mathbb{R}$ ,  $k \in \{1, \dots, p\}$  be  $p$  functions. Assume
  - Each  $\psi_k$  is strictly convex
  - $\sup_{\alpha \in \mathbb{R}} \psi'_j(\alpha) = +\infty$  and  $\inf_{\alpha \in \mathbb{R}} \psi'_j(\alpha) = -\infty$
  - Denote  $\psi_1^*, \dots, \psi_p^*$  their Fenchel-conjugates (then with full domain)



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  - Denote  $\psi_1^*, \dots, \psi_p^*$  their Fenchel-conjugates (then with full domain)

$$\begin{aligned} \min_{w \in \mathbb{R}^p} f(w) + \sum_{j=1}^p \psi_j(w_j) &= \min_{w \in \mathbb{R}^p} \max_{s \in B(F)} w^\top s + \sum_{j=1}^p \psi_j(w_j) \\ &= \max_{s \in B(F)} \min_{w \in \mathbb{R}^p} w^\top s + \sum_{j=1}^p \psi_j(w_j) \\ &= \max_{s \in B(F)} - \sum_{j=1}^p \psi_j^*(-s_j) \end{aligned}$$

# Separable optimization on base polyhedron

## Equivalence with submodular function minimization

- For  $\alpha \in \mathbb{R}$ , let  $A^\alpha \subset V$  be a minimizer of  $A \mapsto F(A) + \sum_{j \in A} \psi'_j(\alpha)$
- Let  $w^*$  be the unique minimizer of  $w \mapsto f(w) + \sum_{j=1}^p \psi_j(w_j)$
- **Proposition** (Chambolle and Darbon, 2009):
  - Given  $A^\alpha$  for all  $\alpha \in \mathbb{R}$ , then  $\forall j, w_j^* = \sup(\{\alpha \in \mathbb{R}, j \in A^\alpha\})$
  - Given  $w^*$ , then  $A \mapsto F(A) + \sum_{j \in A} \psi'_j(\alpha)$  has minimal minimizer  $\{w^* > \alpha\}$  and maximal minimizer  $\{w^* \geq \alpha\}$
- Separable optimization equivalent to a sequence of submodular function minimizations
  - NB: extension of known results from parametric max-flow

# Equivalence with submodular function minimization

## Proof sketch (Bach, 2011b)

- Duality gap for  $\min_{w \in \mathbb{R}^p} f(w) + \sum_{j=1}^p \psi_j(w_j) = \max_{s \in B(F)} - \sum_{j=1}^p \psi_j^*(-s_j)$

$$\begin{aligned} & f(w) + \sum_{j=1}^p \psi_j(w_j) - \sum_{j=1}^p \psi_j^*(-s_j) \\ &= f(w) - w^\top s + \sum_{j=1}^p \left\{ \psi_j(w_j) + \psi_j^*(-s_j) + w_j s_j \right\} \\ &= \int_{-\infty}^{+\infty} \left\{ (F + \psi'(\alpha))(\{w \geq \alpha\}) - (s + \psi'(\alpha))_-(V) \right\} d\alpha \end{aligned}$$

- Duality gap for convex problems = sums of duality gaps for combinatorial problems

# Separable optimization on base polyhedron

## Quadratic case

- Let  $F$  be a submodular function and  $w \in \mathbb{R}^p$  the unique minimizer of  $w \mapsto f(w) + \frac{1}{2}\|w\|_2^2$ . Then:
  - (a)  $s = -w$  is the point in  $B(F)$  with minimum  $\ell_2$ -norm
  - (b) For all  $\lambda \in \mathbb{R}$ , the maximal minimizer of  $A \mapsto F(A) + \lambda|A|$  is  $\{w \geq -\lambda\}$  and the minimal minimizer of  $F$  is  $\{w > -\lambda\}$
- **Consequences**
  - Threshold at 0 the minimum norm point in  $B(F)$  to minimize  $F$  (Fujishige and Isotani, 2011)
  - Minimizing submodular functions with cardinality constraints (Nagano et al., 2011)

# From convex to combinatorial optimization

- Solving  $\min_{w \in \mathbb{R}^p} \sum_{k \in V} \psi_k(w_k) + f(w)$  to solve  $\min_{ACV} F(A)$ 
  - Thresholding solutions  $w$  at zero if  $\forall k \in V, \psi'_k(0) = 0$
  - For quadratic functions  $\psi_k(w_k) = \frac{1}{2}w_k^2$ , equivalent to projecting 0 on  $B(F)$  (Fujishige, 2005)

# From convex to combinatorial optimization and vice-versa...

- Solving  $\min_{w \in \mathbb{R}^p} \sum_{k \in V} \psi_k(w_k) + f(w)$  to solve  $\min_{A \subset V} F(A)$ 
  - Thresholding solutions  $w$  at zero if  $\forall k \in V, \psi'_k(0) = 0$
  - For quadratic functions  $\psi_k(w_k) = \frac{1}{2}w_k^2$ , equivalent to projecting 0 on  $B(F)$  (Fujishige, 2005)
- Solving  $\min_{A \subset V} F(A) - t(A)$  to solve  $\min_{w \in \mathbb{R}^p} \sum_{k \in V} \psi_k(w_k) + f(w)$ 
  - General decomposition strategy (Groenevelt, 1991)
  - Efficient only when submodular minimization is efficient

**Solving**  $\min_{A \subset V} F(A) - t(A)$  **to solve**  $\min_{w \in \mathbb{R}^p} \sum_{k \in V} \psi_k(w_k) + f(w)$

- General **recursive divide-and-conquer** algorithm (Groenevelt, 1991)
- NB: Dual version of Fujishige (2005)
  1. Compute minimizer  $t \in \mathbb{R}^p$  of  $\sum_{j \in V} \psi_j^*(-t_j)$  s.t.  $t(V) = F(V)$
  2. Compute minimizer  $A$  of  $F(A) - t(A)$
  3. If  $A = V$ , then  $t$  is optimal. Exit.
  4. Compute a minimizer  $s_A$  of  $\sum_{j \in A} \psi_j^*(-s_j)$  over  $s \in B(F_A)$  where  $F_A : 2^A \rightarrow \mathbb{R}$  is the restriction of  $F$  to  $A$ , i.e.,  $F_A(B) = F(A)$
  5. Compute a minimizer  $s_{V \setminus A}$  of  $\sum_{j \in V \setminus A} \psi_j^*(-s_j)$  over  $s \in B(F^A)$  where  $F^A(B) = F(A \cup B) - F(A)$ , for  $B \subset V \setminus A$
  6. Concatenate  $s_A$  and  $s_{V \setminus A}$ . Exit.

**Solving**  $\min_{w \in \mathbb{R}^p} \sum_{k \in V} \psi_k(w_k) + f(w)$  **to solve**  $\min_{ACV} F(A)$

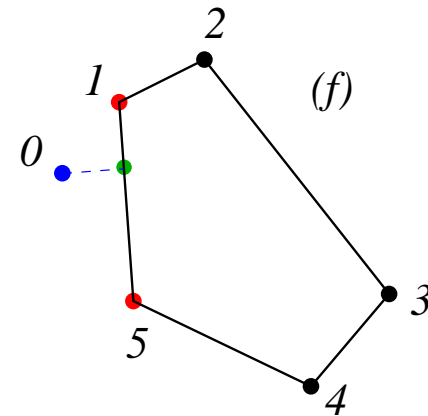
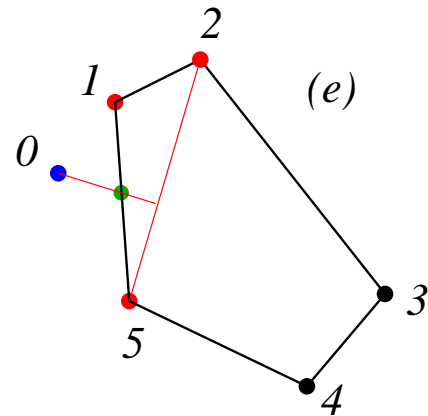
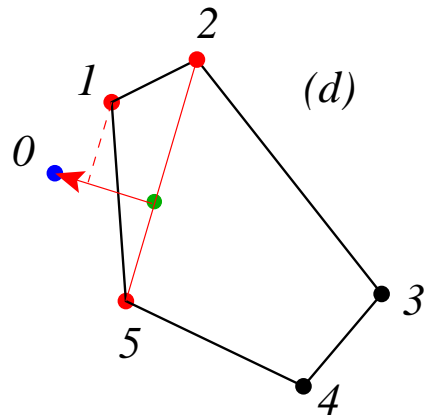
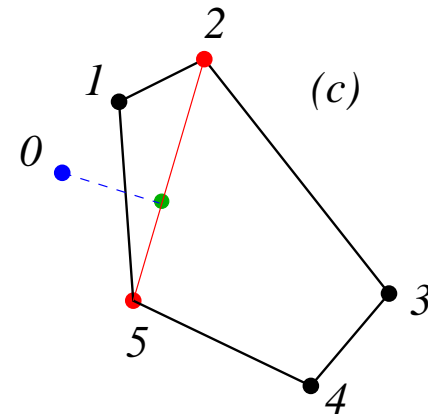
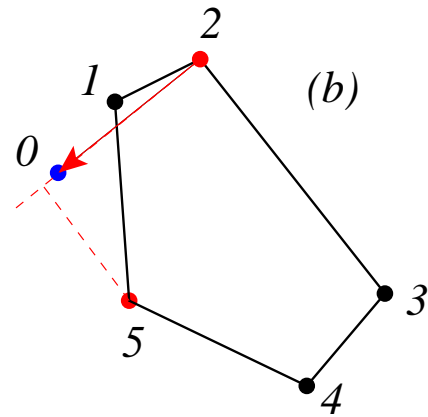
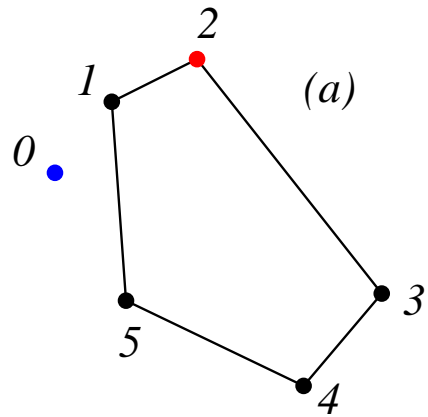
- Dual problem:  $\max_{s \in B(F)} - \sum_{j=1}^p \psi_j^*(-s_j)$
- Constrained optimization when linear functions can be maximized
  - **Frank-Wolfe algorithms**
- Two main types for convex functions



# Approximate quadratic optimization on $B(F)$

- **Goal:**  $\min_{w \in \mathbb{R}^p} \frac{1}{2} \|w\|_2^2 + f(w) = \max_{s \in B(F)} -\frac{1}{2} \|s\|_2^2$
- Can only maximize linear functions on  $B(F)$
- **Two types of “Frank-wolfe” algorithms**
- **1. Active set algorithm ( $\Leftrightarrow$  min-norm-point)**
  - Sequence of maximizations of linear functions over  $B(F)$   
+ overheads (affine projections)
  - Finite convergence, but no complexity bounds

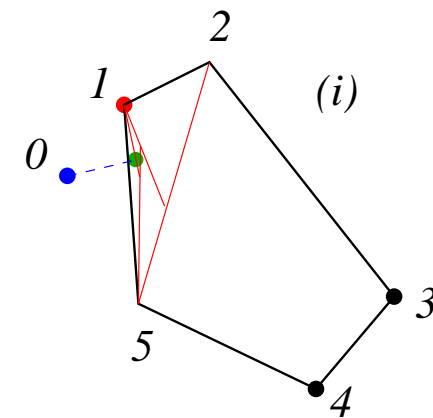
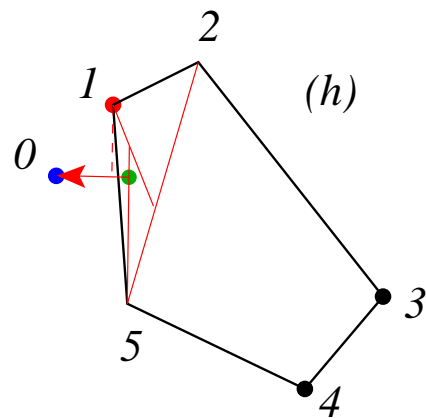
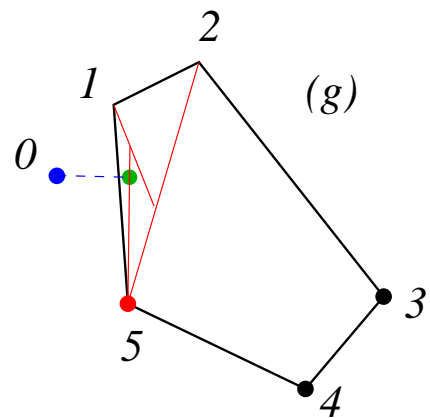
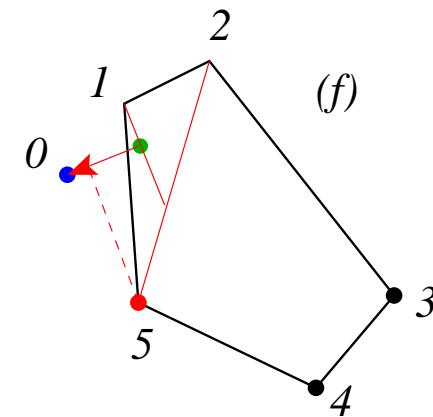
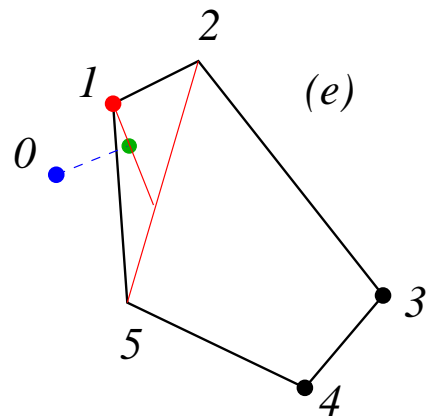
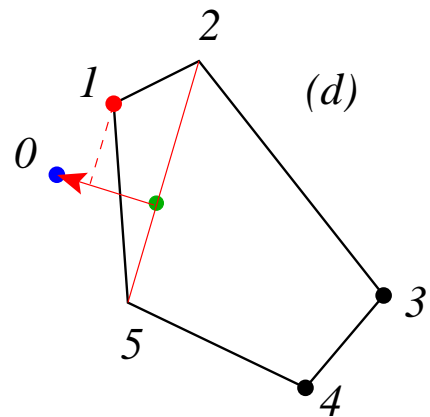
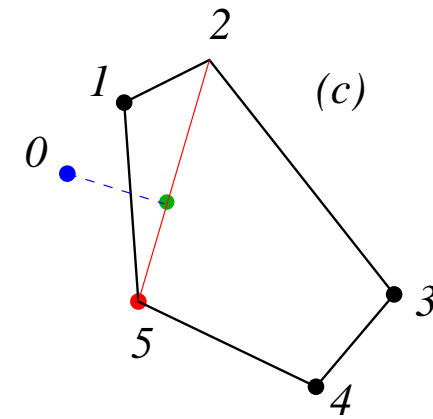
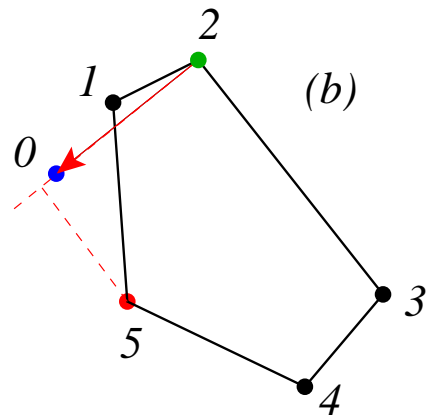
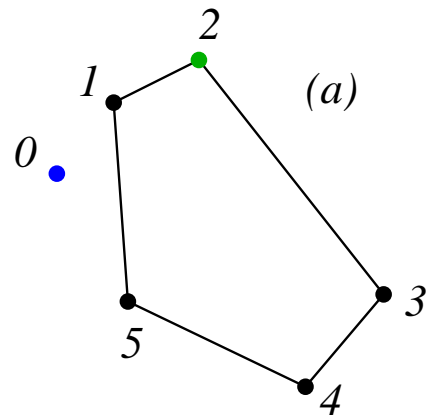
# Minimum-norm-point algorithm (Wolfe, 1976)



# Approximate quadratic optimization on $B(F)$

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- **Two types of “Frank-wolfe” algorithms**
- **1. Active set algorithm ( $\Leftrightarrow$  min-norm-point)**
  - Sequence of maximizations of linear functions over  $B(F)$   
+ overheads (affine projections)
  - Finite convergence, but no complexity bounds
- **2. Conditional gradient**
  - Sequence of maximizations of linear functions over  $B(F)$
  - Approximate optimality bound

# Conditional gradient with line search



# Approximate quadratic optimization on $B(F)$

- **Proposition:**  $t$  steps of **conditional gradient** (with line search) outputs  $s_t \in B(F)$  and  $w_t = -s_t$ , such that

$$f(w_t) + \frac{1}{2}\|w_t\|_2^2 - \text{OPT} \leq f(w_t) + \frac{1}{2}\|w_t\|_2^2 + \frac{1}{2}\|s_t\|_2^2 \leq \frac{2D^2}{t}$$

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- **Improved primal candidate through isotonic regression**
  - $f(w)$  is linear on any set of  $w$  with fixed ordering
  - May be optimized using isotonic regression (“pool-adjacent-violator”) in  $O(n)$  (see, e.g., Best and Chakravarti, 1990)
  - Given  $w_t = -s_t$ , keep the ordering and reoptimize

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  - Given  $w_t = -s_t$ , keep the ordering and reoptimize
- **Better bound for submodular function minimization?**

# From quadratic optimization on $B(F)$ to submodular function minimization

- **Proposition:** If  $w$  is  $\varepsilon$ -optimal for  $\min_{w \in \mathbb{R}^p} \frac{1}{2} \|w\|_2^2 + f(w)$ , then at least a level set  $A$  of  $w$  is  $(\frac{\sqrt{\varepsilon p}}{2})$ -optimal for submodular function minimization
- If  $\varepsilon = \frac{2D^2}{t}$ ,  $\frac{\sqrt{\varepsilon p}}{2} = \frac{Dp^{1/2}}{\sqrt{2t}} \Rightarrow$  **no provable gains**, but:
  - Bound on the iterates  $A_t$  (with additional assumptions)
  - Possible thresholding for acceleration



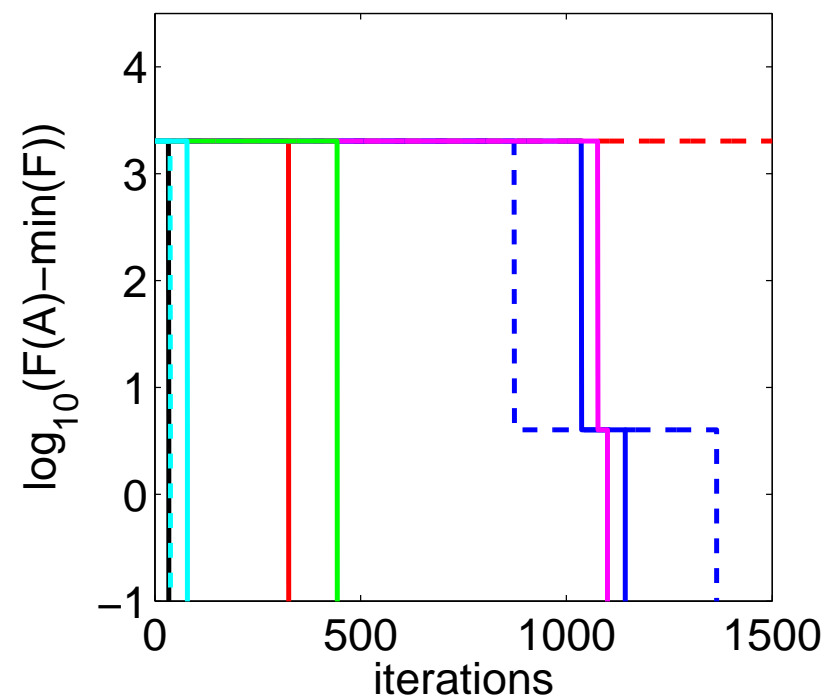
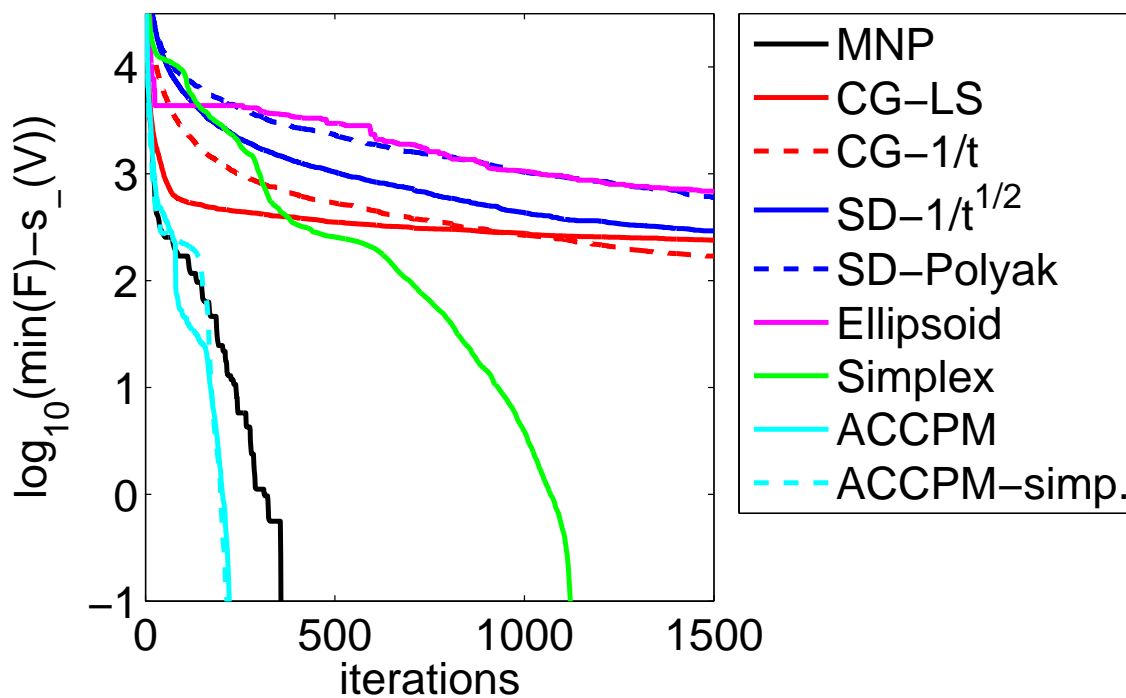
# From quadratic optimization on $B(F)$ to submodular function minimization

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  - Bound on the iterates  $A_t$  (with additional assumptions)
  - Possible thresholding for acceleration
- **Lower complexity bound for SFM**
  - **Conjecture:** no algorithm that is based **only** on a sequence of greedy algorithms obtained from linear combinations of bases can improve on the subgradient bound (after  $p/2$  iterations).

# Simulations on standard benchmark “DIMACS Genrmf-wide”, $p = 430$

- **Submodular function minimization**

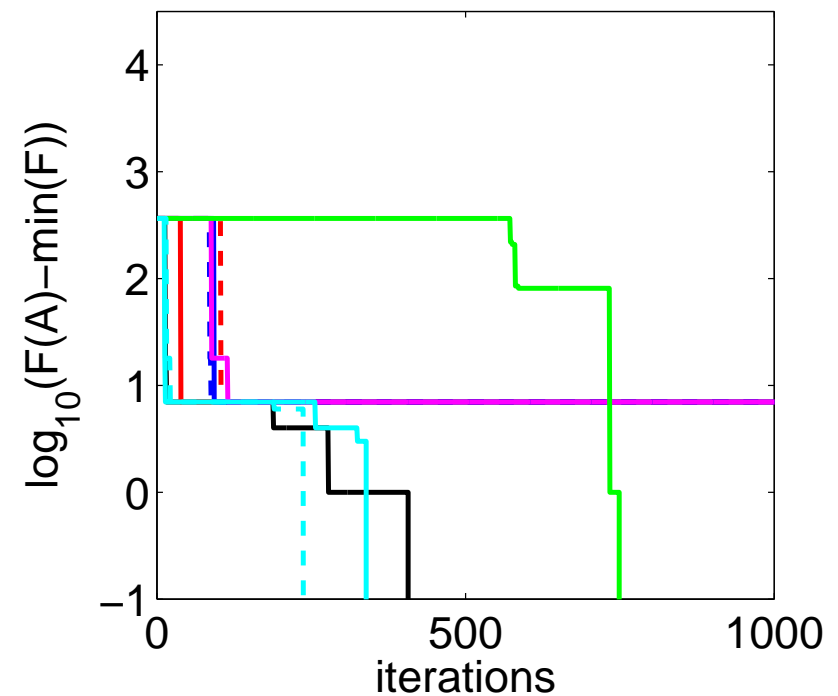
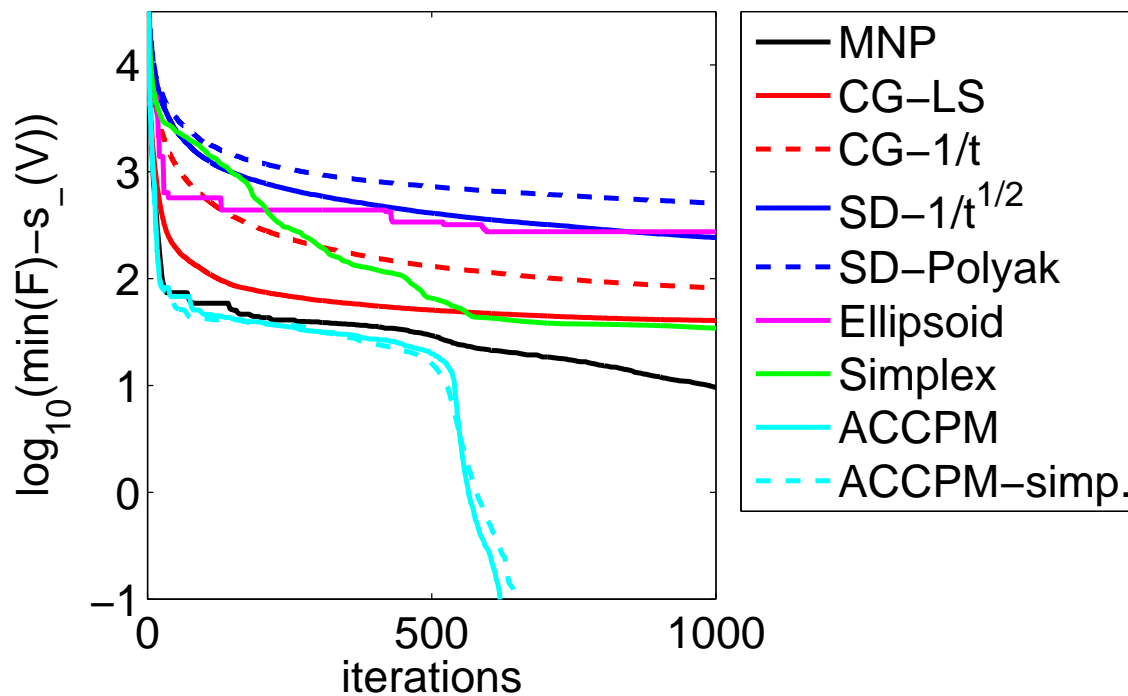
- (Left) **dual** suboptimality
- (Right) **primal** suboptimality



# Simulations on standard benchmark “DIMACS Genrmf-long”, $p = 575$

- **Submodular function minimization**

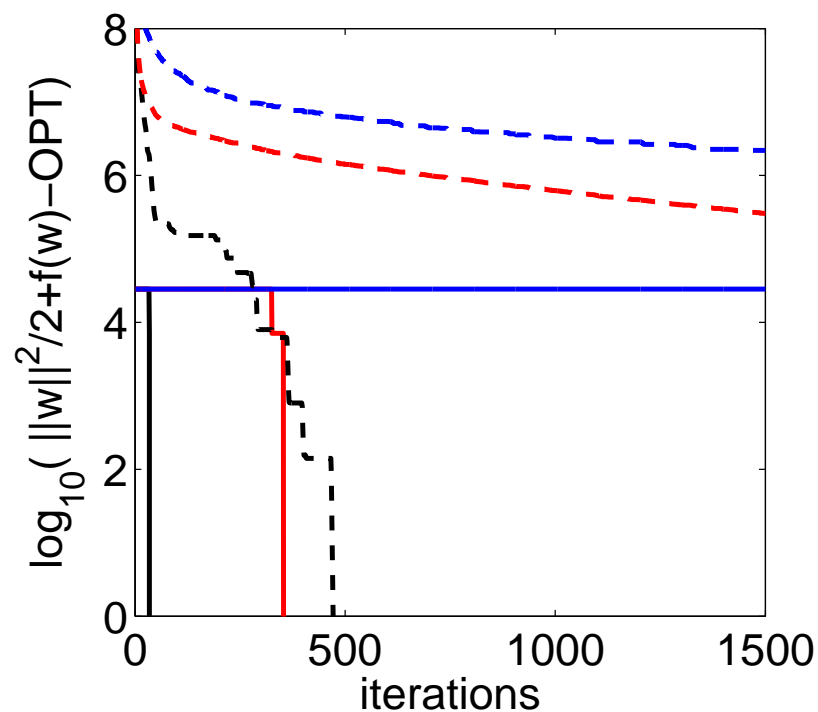
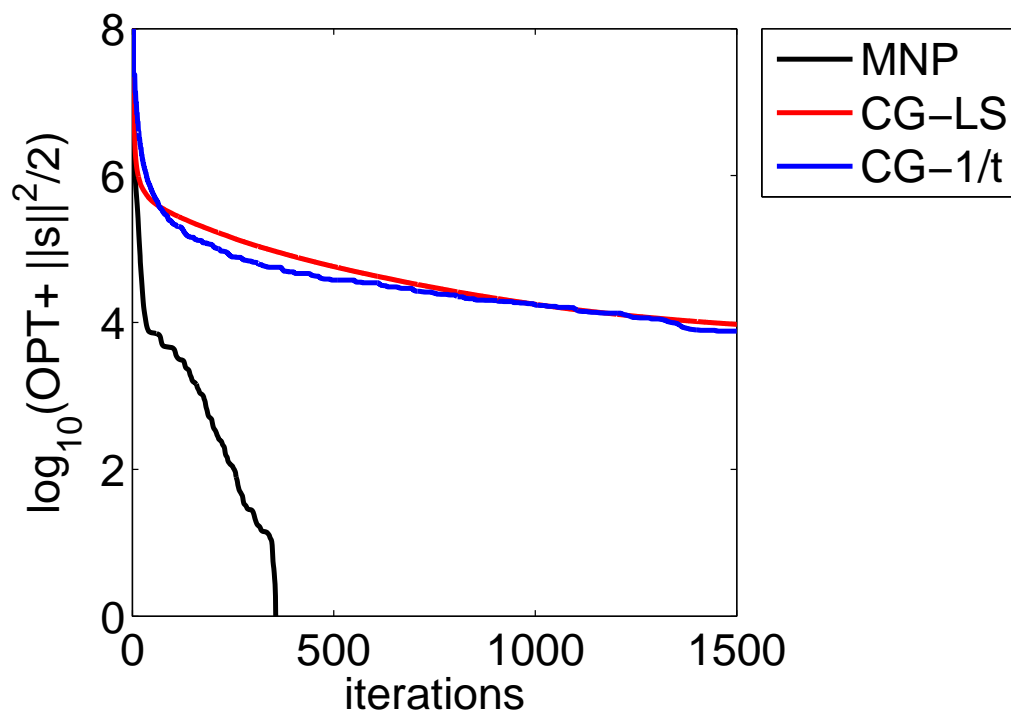
- (Left) **dual** suboptimality
- (Right) **primal** suboptimality



# Simulations on standard benchmark

- **Separable quadratic optimization**

- (Left) **dual** suboptimality
- (Right) **primal** suboptimality  
(in dashed, before the pool-adjacent-violator correction)



# Outline

## 1. Submodular functions

- Review and examples of submodular functions
- Links with convexity through Lovász extension

## 2. Submodular minimization

- Non-smooth convex optimization
- Parallel algorithm for special case

## 3. Structured sparsity-inducing norms

- Relaxation of the penalization of supports by submodular functions
- Extensions (symmetric,  $\ell_q$ -relaxation)

# From submodular minimization to proximal problems

- **Summary:** several optimization problems

- Discrete problem:  $\min_{A \subseteq V} F(A) = \min_{w \in \{0,1\}^p} f(w)$

- Continuous problem:  $\min_{w \in [0,1]^p} f(w)$

- Proximal problem (P):  $\min_{w \in \mathbb{R}^p} \frac{1}{2} \|w\|_2^2 + f(w)$

- **Solving (P) is equivalent to minimizing  $F(A) + \lambda|A|$  for all  $\lambda$**

- $\arg \min_{A \subseteq V} F(A) + \lambda|A| = \{k, w_k \geq -\lambda\}$

- Much simpler problem but no gains in terms of (provable) complexity
  - See Bach (2011a)

# Decomposable functions

- $F$  may often be decomposed as the sum of  $r$  “simple” functions:

$$F(A) = \sum_{j=1}^r F_j(A)$$

- Each  $F_j$  may be minimized efficiently
- Example: 2D grid = vertical chains + horizontal chains
- Komodakis et al. (2011); Kolmogorov (2012); Stobbe and Krause (2010); Savchynskyy et al. (2011)
  - Dual decomposition approach but slow non-smooth problem

# Decomposable functions and proximal problems (Jegelka, Bach, and Sra, 2013)

- Dual problem

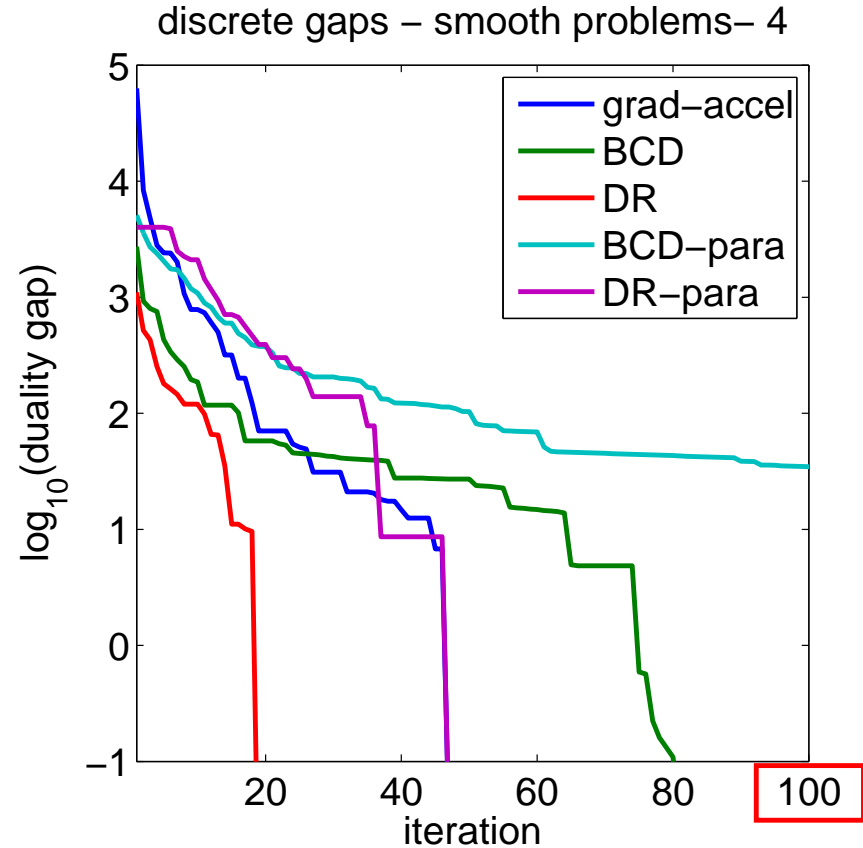
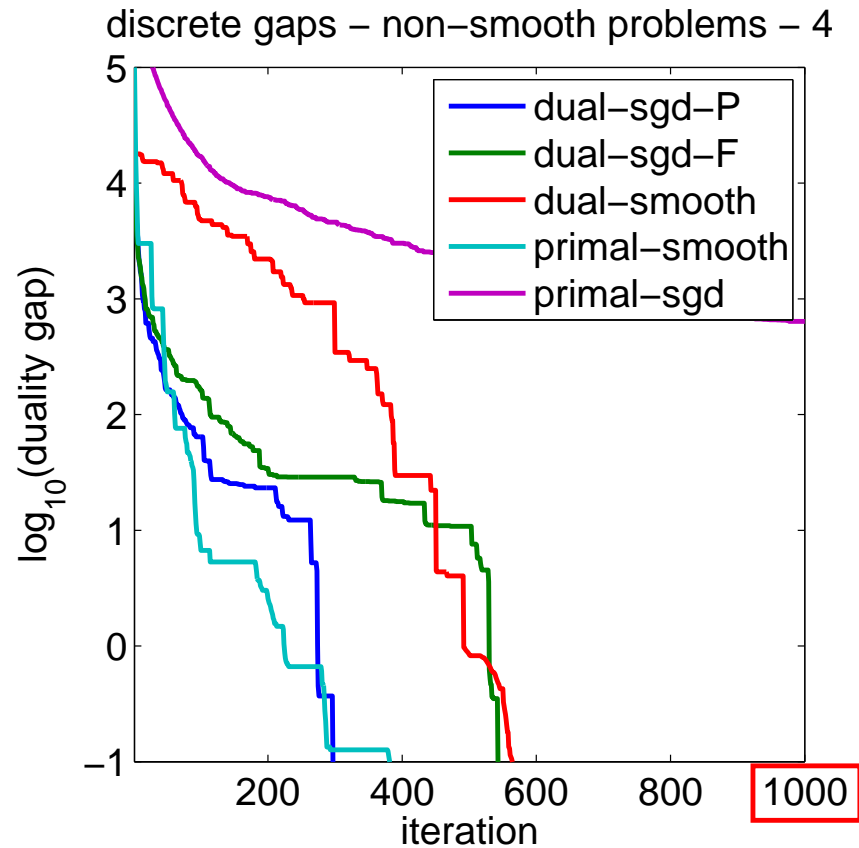
$$\begin{aligned} & \min_{w \in \mathbb{R}^p} f_1(w) + f_2(w) + \frac{1}{2} \|w\|_2^2 \\ &= \min_{w \in \mathbb{R}^p} \max_{s_1 \in B(F_1)} s_1^\top w + \max_{s_2 \in B(F_2)} s_2^\top w + \frac{1}{2} \|w\|_2^2 \\ &= \max_{s_1 \in B(F_1), s_2 \in B(F_2)} -\frac{1}{2} \|s_1 + s_2\|^2 \end{aligned}$$

- **Finding the closest point between two polytopes**
  - Several alternatives: Block coordinate ascent, Douglas Rachford splitting (Bauschke et al., 2004)
  - (a) no parameters, (b) parallelizable



# Experiments

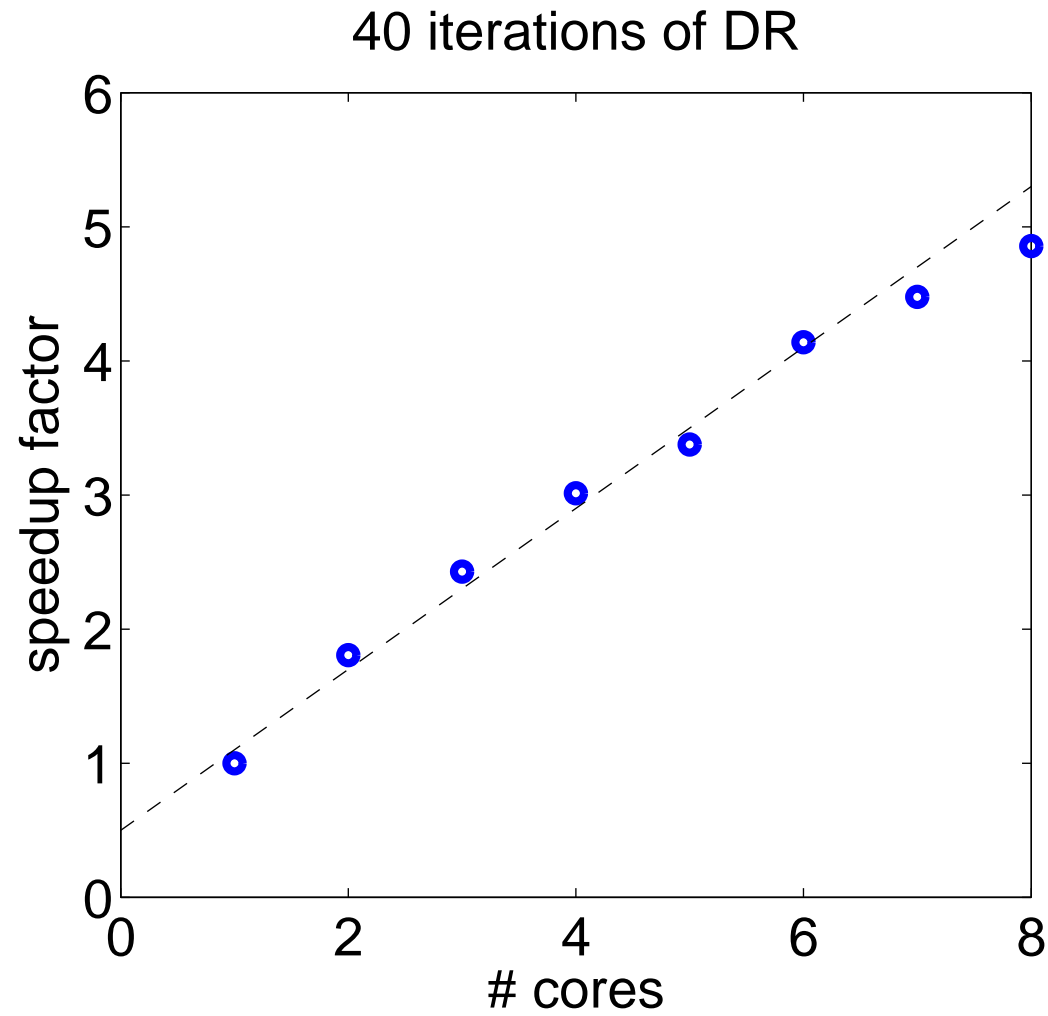
- Graph cuts on a  $500 \times 500$  image



- Matlab/C implementation 10 times slower than C-code for graph cut
  - Easy to code and parallelizable

# Parallelization

- Multiple cores



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# Structured sparsity through submodular functions

## References and Links

- **References on submodular functions**

- *Submodular Functions and Optimization* (Fujishige, 2005)
- Tutorial paper based on convex optimization (Bach, 2011b)

[www.di.ens.fr/~fbach/submodular\\_fot.pdf](http://www.di.ens.fr/~fbach/submodular_fot.pdf)

- **Structured sparsity through convex optimization**

- Algorithms (Bach, Jenatton, Mairal, and Obozinski, 2011)
- Theory/applications (Bach, Jenatton, Mairal, and Obozinski, 2012)

[www.di.ens.fr/~fbach/bach\\_jenatton\\_mairal\\_obozinski\\_FOT.pdf](http://www.di.ens.fr/~fbach/bach_jenatton_mairal_obozinski_FOT.pdf)

[www.di.ens.fr/~fbach/stat\\_science\\_structured\\_sparsity.pdf](http://www.di.ens.fr/~fbach/stat_science_structured_sparsity.pdf)

- Matlab/R/Python codes: <http://www.di.ens.fr/willow/SPAMS/>

- **Slides:** [www.di.ens.fr/~fbach/fbach\\_cargese\\_2013.pdf](http://www.di.ens.fr/~fbach/fbach_cargese_2013.pdf)

# Sparsity in supervised machine learning

- Observed data  $(x_i, y_i) \in \mathbb{R}^p \times \mathbb{R}$ ,  $i = 1, \dots, n$ 
  - Response vector  $y = (y_1, \dots, y_n)^\top \in \mathbb{R}^n$
  - Design matrix  $X = (x_1, \dots, x_n)^\top \in \mathbb{R}^{n \times p}$
- Regularized empirical risk minimization:

$$\min_{w \in \mathbb{R}^p} \frac{1}{n} \sum_{i=1}^n \ell(y_i, w^\top x_i) + \lambda \Omega(w) = \boxed{\min_{w \in \mathbb{R}^p} L(y, Xw) + \lambda \Omega(w)}$$

- Norm  $\Omega$  to promote sparsity
  - square loss +  $\ell_1$ -norm  $\Rightarrow$  **basis pursuit** in signal processing (Chen et al., 2001), **Lasso** in statistics/machine learning (Tibshirani, 1996)
  - Proxy for **interpretability**
  - Allow **high-dimensional inference**:  $\boxed{\log p = O(n)}$

# Sparsity in **unsupervised** machine learning

- **Multiple** responses/signals  $y = (y^1, \dots, y^k) \in \mathbb{R}^{n \times k}$

$$\min_{w^1, \dots, w^k \in \mathbb{R}^p} \sum_{j=1}^k \left\{ L(y^j, Xw^j) + \lambda \Omega(w^j) \right\}$$

# Sparsity in **unsupervised** machine learning

- **Multiple** responses/signals  $y = (y^1, \dots, y^k) \in \mathbb{R}^{n \times k}$

$$\min_{w^1, \dots, w^k \in \mathbb{R}^p} \sum_{j=1}^k \left\{ L(y^j, X w^j) + \lambda \Omega(w^j) \right\}$$

- **Only responses are observed**  $\Rightarrow$  **Dictionary learning**

– Learn  $X = (x^1, \dots, x^p) \in \mathbb{R}^{n \times p}$  such that  $\forall j, \|x^j\|_2 \leq 1$

$$\min_{X=(x^1, \dots, x^p)} \min_{w^1, \dots, w^k \in \mathbb{R}^p} \sum_{j=1}^k \left\{ L(y^j, X w^j) + \lambda \Omega(w^j) \right\}$$

– Olshausen and Field (1997); Elad and Aharon (2006); Mairal et al. (2009a)

- **sparse PCA**: replace  $\|x^j\|_2 \leq 1$  by  $\Theta(x^j) \leq 1$

# Sparsity in signal processing

- **Multiple** responses/signals  $x = (x^1, \dots, x^k) \in \mathbb{R}^{n \times k}$

$$\min_{\alpha^1, \dots, \alpha^k \in \mathbb{R}^p} \sum_{j=1}^k \left\{ L(x^j, D\alpha^j) + \lambda \Omega(\alpha^j) \right\}$$

- **Only responses are observed**  $\Rightarrow$  **Dictionary learning**

– Learn  $D = (d^1, \dots, d^p) \in \mathbb{R}^{n \times p}$  such that  $\forall j, \|d^j\|_2 \leq 1$

$$\min_{D=(d^1, \dots, d^p)} \min_{\alpha^1, \dots, \alpha^k \in \mathbb{R}^p} \sum_{j=1}^k \left\{ L(x^j, D\alpha^j) + \lambda \Omega(\alpha^j) \right\}$$

– Olshausen and Field (1997); Elad and Aharon (2006); Mairal et al. (2009a)

- **sparse PCA**: replace  $\|d^j\|_2 \leq 1$  by  $\Theta(d^j) \leq 1$



# Why structured sparsity?

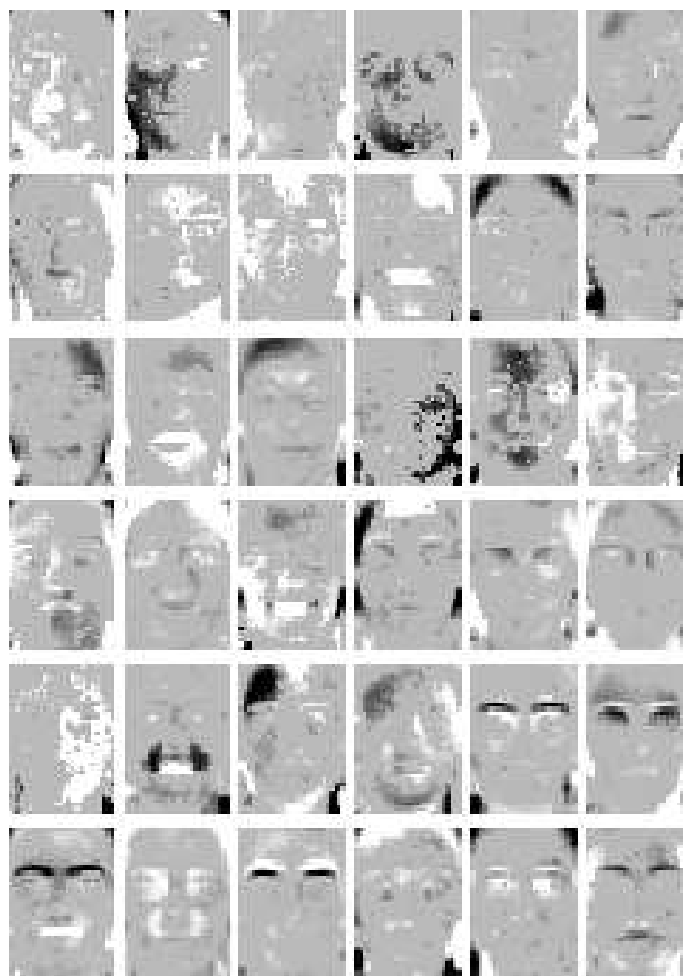
- **Interpretability**

- Structured dictionary elements (Jenatton et al., 2009b)
- Dictionary elements “organized” in a **tree** or a **grid** (Kavukcuoglu et al., 2009; Jenatton et al., 2010; Mairal et al., 2010)

# Structured sparse PCA (Jenatton et al., 2009b)



raw data



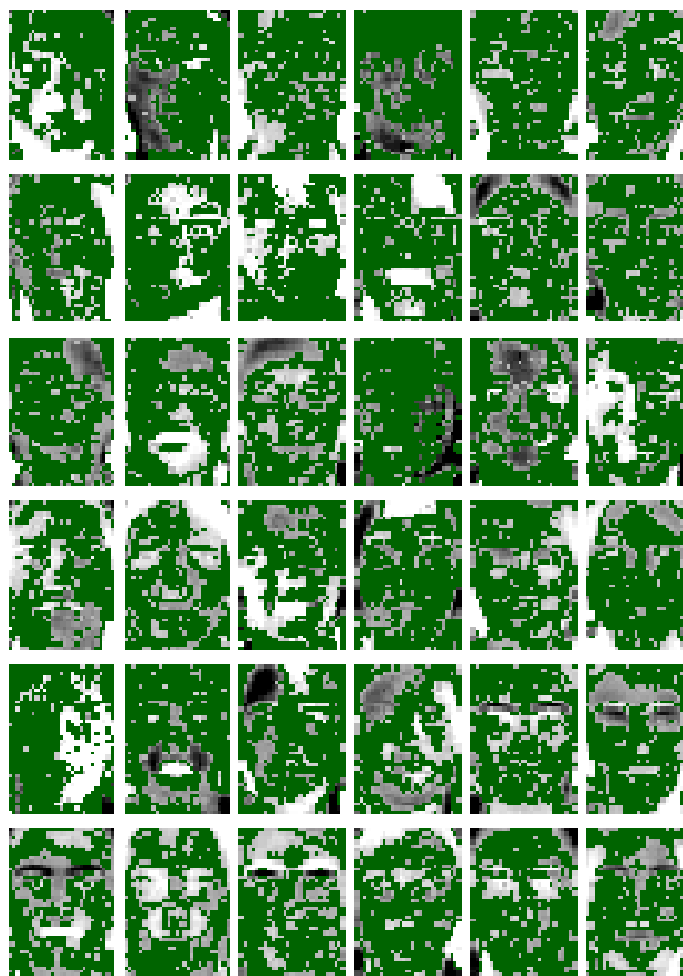
sparse PCA

- Unstructured sparse PCA  $\Rightarrow$  many zeros do not lead to better interpretability

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raw data



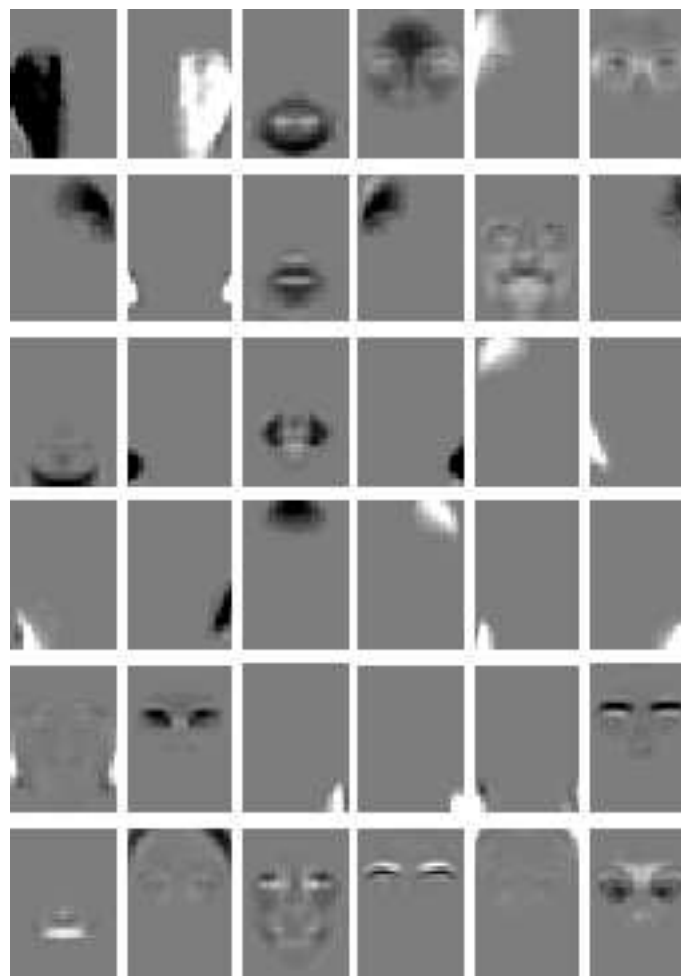
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raw data



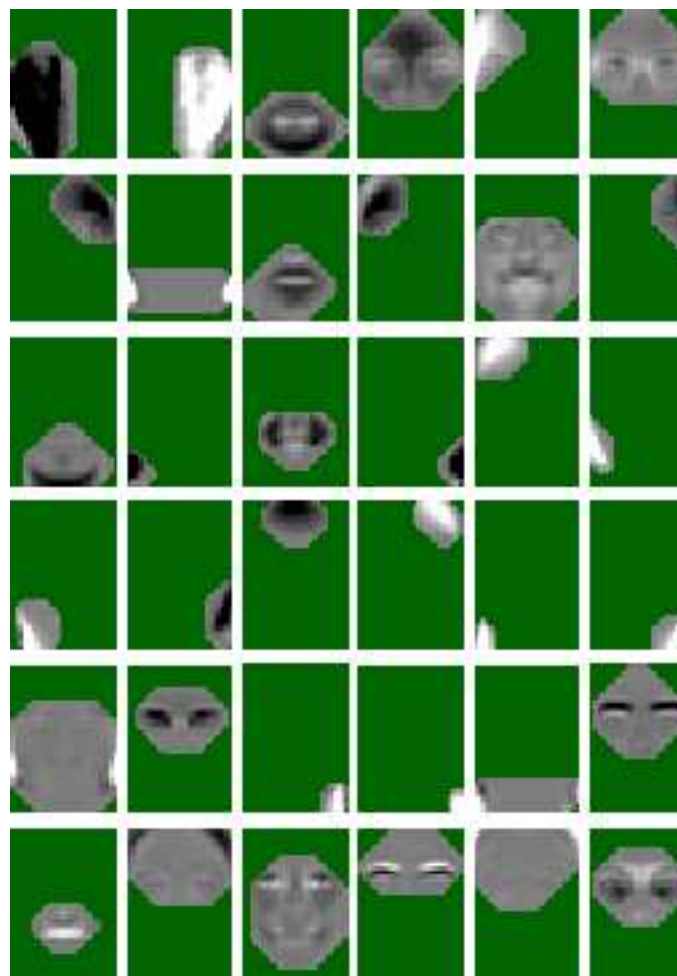
Structured sparse PCA

- Enforce selection of **convex** nonzero patterns  $\Rightarrow$  robustness to occlusion in face identification

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raw data



Structured sparse PCA

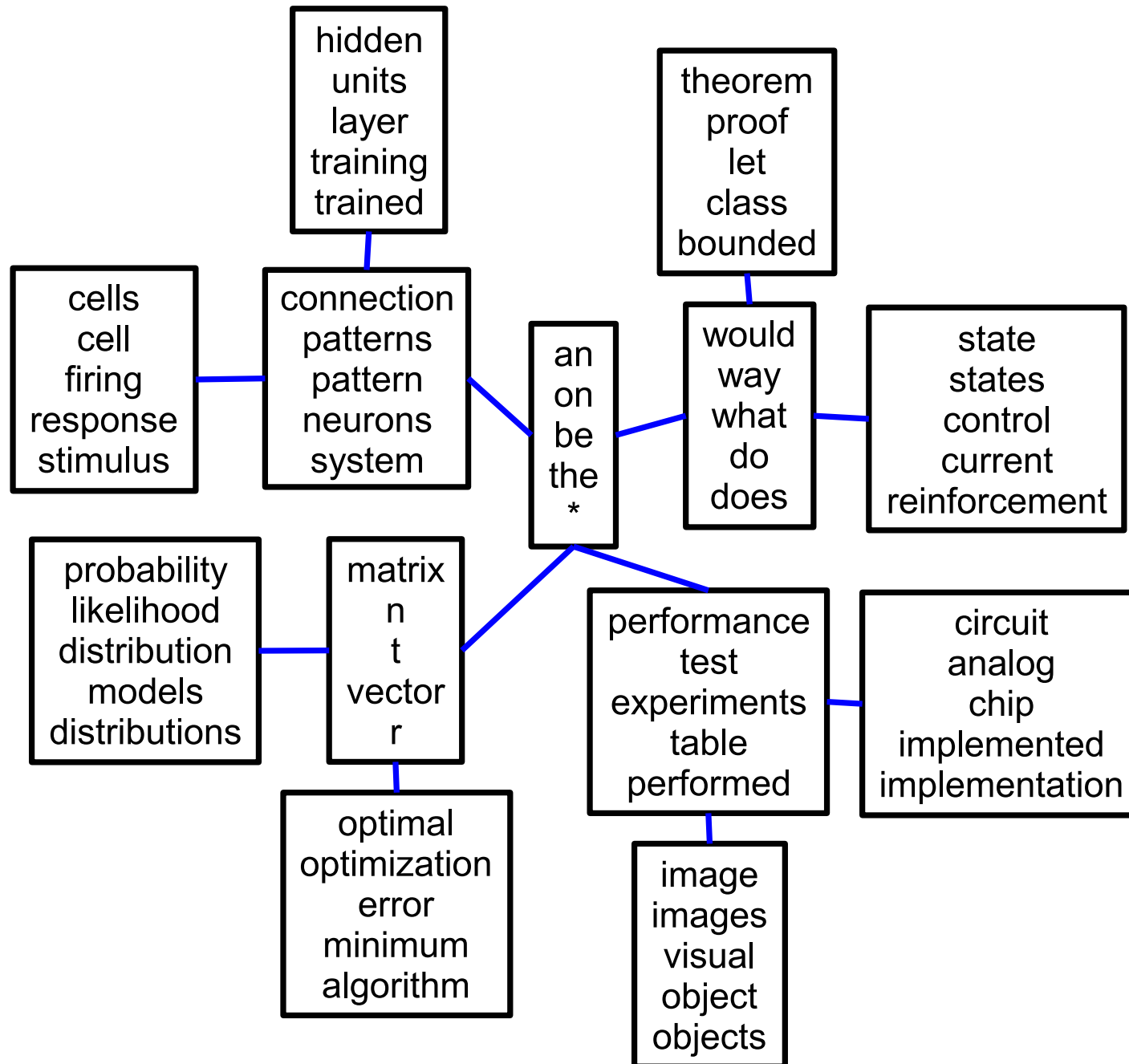
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# Modelling of text corpora (Jenatton et al., 2010)



# Why structured sparsity?

- **Interpretability**

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# Why structured sparsity?

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- Structured dictionary elements (Jenatton et al., 2009b)
- Dictionary elements “organized” in a **tree** or a **grid** (Kavukcuoglu et al., 2009; Jenatton et al., 2010; Mairal et al., 2010)

- **Stability and identifiability**

- **Prediction or estimation performance**

- When prior knowledge matches data (Haupt and Nowak, 2006; Baraniuk et al., 2008; Jenatton et al., 2009a; Huang et al., 2009)

- **Numerical efficiency**

- Non-linear variable selection with  $2^p$  subsets (Bach, 2008)

# Classical approaches to structured sparsity

- **Many application domains**

- Computer vision (Cevher et al., 2008; Mairal et al., 2009b)
- Neuro-imaging (Gramfort and Kowalski, 2009; Jenatton et al., 2011)
- Bio-informatics (Rapaport et al., 2008; Kim and Xing, 2010)

- **Non-convex approaches**

- Haupt and Nowak (2006); Baraniuk et al. (2008); Huang et al. (2009)

- **Convex approaches**

- Design of sparsity-inducing norms

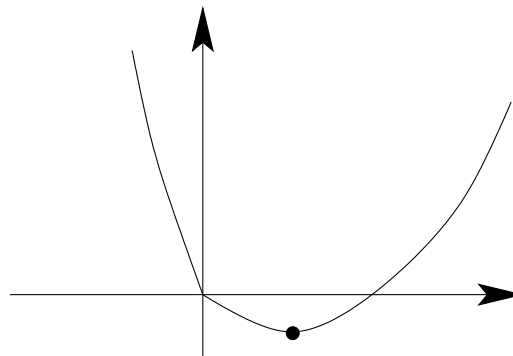
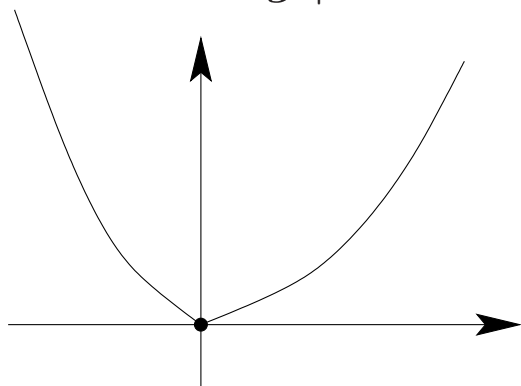
# Why $\ell_1$ -norms lead to sparsity?

- **Example 1:** quadratic problem in 1D, i.e.,

$$\min_{x \in \mathbb{R}} \frac{1}{2}x^2 - xy + \lambda|x|$$

- Piecewise quadratic function with a kink at zero

- Derivative at  $0+$ :  $g_+ = \lambda - y$  and  $0-$ :  $g_- = -\lambda - y$

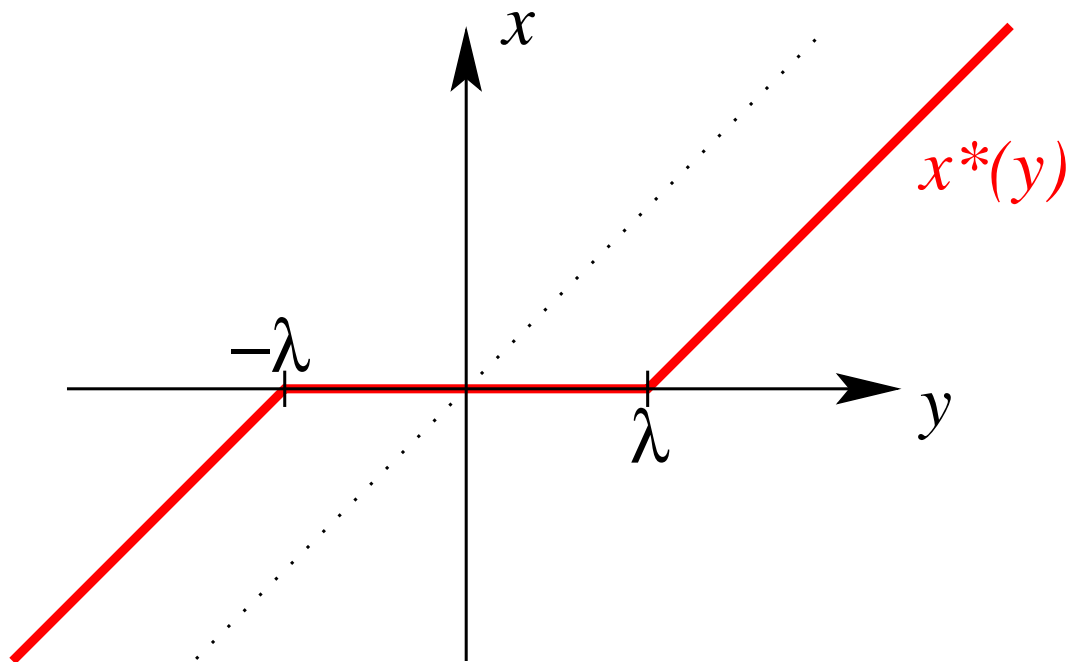


- $x = 0$  is the solution iff  $g_+ \geq 0$  and  $g_- \leq 0$  (i.e.,  $|y| \leq \lambda$ )
- $x \geq 0$  is the solution iff  $g_+ \leq 0$  (i.e.,  $y \geq \lambda$ )  $\Rightarrow x^* = y - \lambda$
- $x \leq 0$  is the solution iff  $g_- \leq 0$  (i.e.,  $y \leq -\lambda$ )  $\Rightarrow x^* = y + \lambda$

- Solution  $x^* = \text{sign}(y)(|y| - \lambda)_+$  = **soft thresholding**

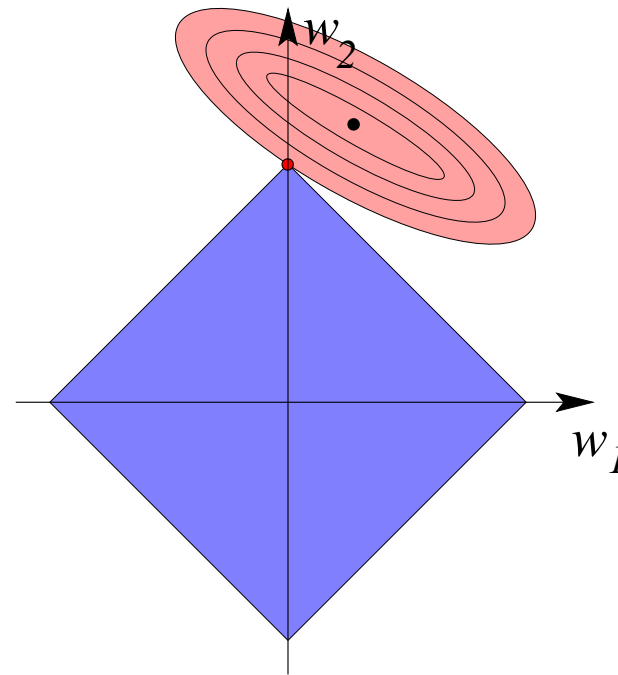
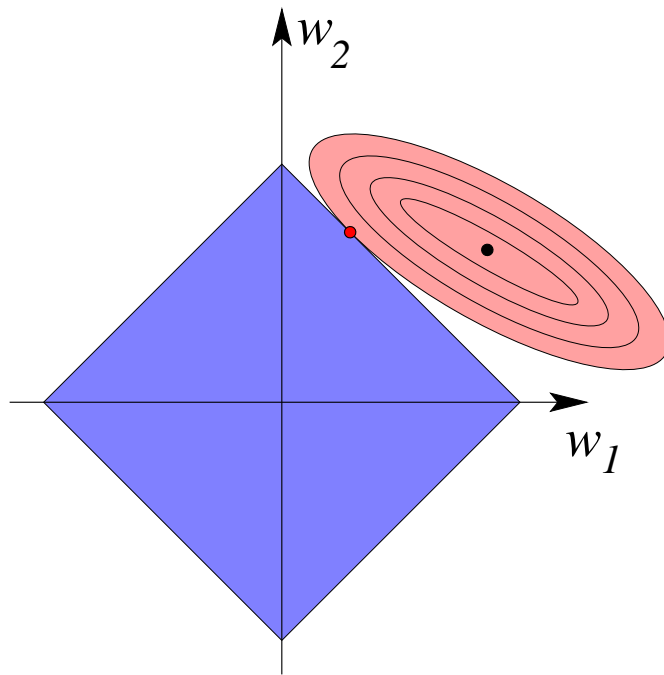
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# Why $\ell_1$ -norms lead to sparsity?

- **Example 2:** minimize quadratic function  $Q(w)$  subject to  $\|w\|_1 \leq T$ .
  - **coupled soft** thresholding
- Geometric interpretation
  - NB : penalizing is “equivalent” to constraining



- **Non-smooth optimization!**

# Gaussian hare ( $\ell_2$ ) vs. Laplacian tortoise ( $\ell_1$ )



- Smooth vs. non-smooth optimization
- See Bach, Jenatton, Mairal, and Obozinski (2011)

# Sparsity-inducing norms

- Popular choice for  $\Omega$

- The  $\ell_1$ - $\ell_2$  norm,

$$\sum_{G \in \mathbf{H}} \|w_G\|_2 = \sum_{G \in \mathbf{H}} \left( \sum_{j \in G} w_j^2 \right)^{1/2}$$

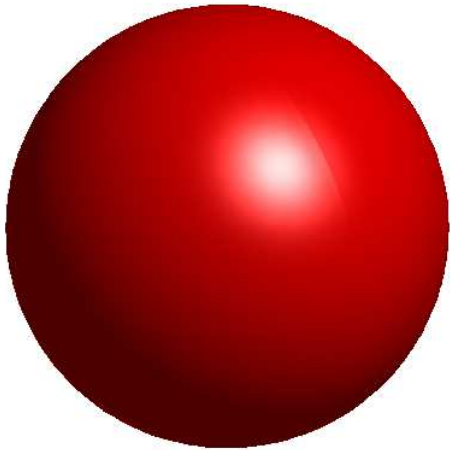
- with  $\mathbf{H}$  a **partition** of  $\{1, \dots, p\}$
- The  $\ell_1$ - $\ell_2$  norm sets to zero **groups of non-overlapping variables** (as opposed to single variables for the  $\ell_1$ -norm)
- For the square loss, group Lasso (Yuan and Lin, 2006)



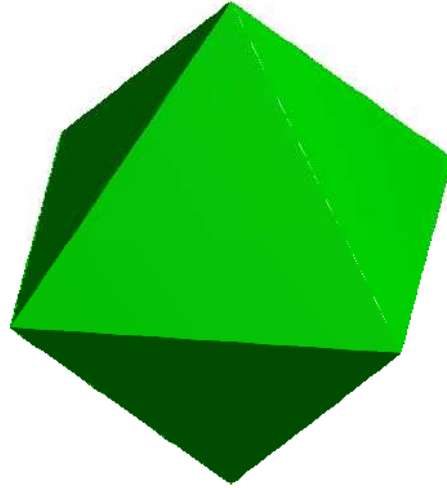


# Unit norm balls

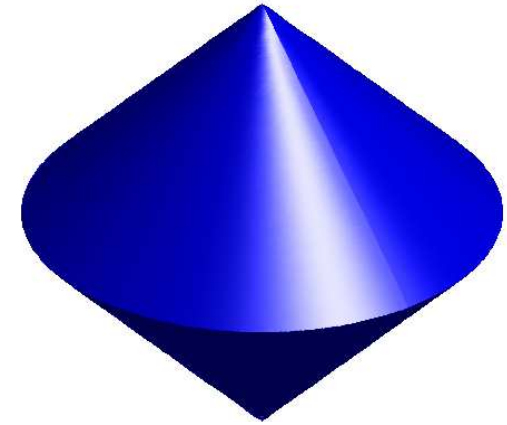
## Geometric interpretation



$$\|w\|_2$$



$$\|w\|_1$$



$$\sqrt{w_1^2 + w_2^2} + |w_3|$$



# Sparsity-inducing norms

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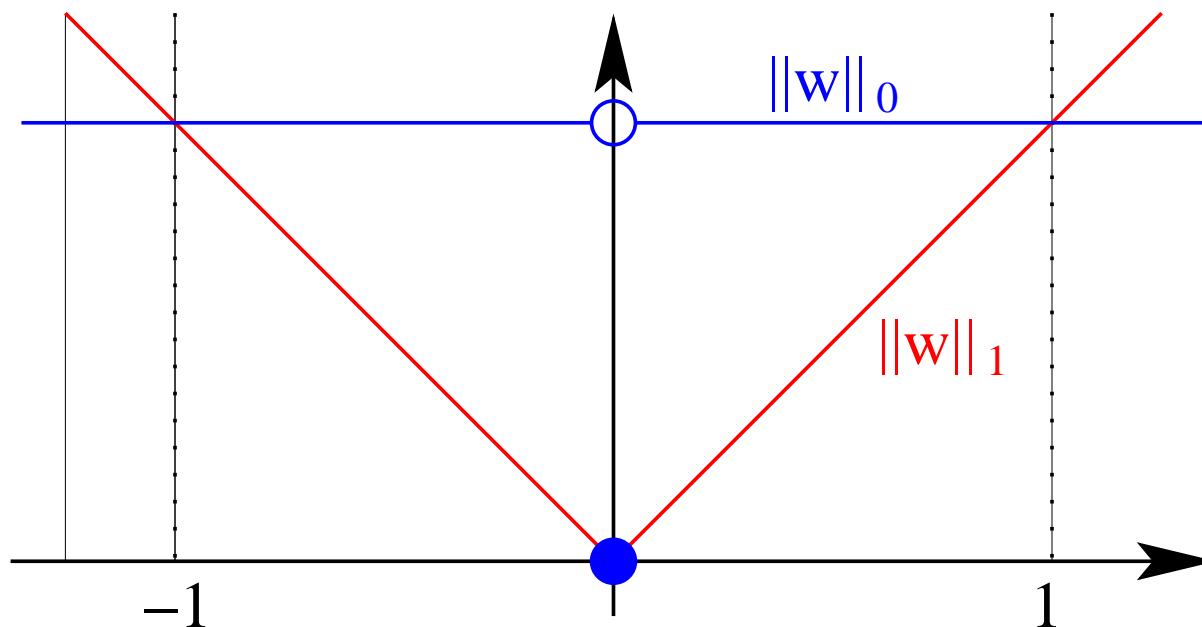
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- What if the set of groups  $\mathbf{H}$  is not a partition anymore?
- **Is there any systematic way?**

## $\ell_1$ -norm = convex envelope of cardinality of support

- Let  $w \in \mathbb{R}^p$ . Let  $V = \{1, \dots, p\}$  and  $\text{Supp}(w) = \{j \in V, w_j \neq 0\}$
- **Cardinality of support:**  $\|w\|_0 = \text{Card}(\text{Supp}(w))$
- Convex envelope = largest convex lower bound (see, e.g., Boyd and Vandenberghe, 2004)



- $\ell_1$ -norm = convex envelope of  $\ell_0$ -quasi-norm on the  $\ell_\infty$ -ball  $[-1, 1]^p$

# Convex envelopes of general functions of the support (Bach, 2010)

- Let  $F : 2^V \rightarrow \mathbb{R}$  be a **set-function**
  - Assume  $F$  is **non-decreasing** (i.e.,  $A \subset B \Rightarrow F(A) \leq F(B)$ )
  - Explicit prior knowledge on supports (Haupt and Nowak, 2006; Baraniuk et al., 2008; Huang et al., 2009)
- Define  $\Theta(w) = F(\text{Supp}(w))$ : **How to get its convex envelope?**
  1. Possible if  $F$  is also **submodular**
  2. Allows **unified** theory and algorithm
  3. Provides **new** regularizers

# Submodular functions and structured sparsity

- Let  $F : 2^V \rightarrow \mathbb{R}$  be a **non-decreasing submodular set-function**
- **Proposition:** the convex envelope of  $\Theta : w \mapsto F(\text{Supp}(w))$  on the  $\ell_\infty$ -ball is  $\Omega : w \mapsto f(|w|)$  where  $f$  is the Lovász extension of  $F$

# Proof - I

- Notation:  $g : w \mapsto F(\text{supp}(w))$  defined on  $[-1, 1]^p$
- Computation of the **Fenchel dual**

$$\begin{aligned} g^*(s) &= \max_{\|w\|_\infty \leq 1} w^\top s - g(w) \\ &= \max_{\delta \in \{0,1\}^p} \max_{\|w\|_\infty \leq 1} (\delta \circ w)^\top s - f(\delta) \text{ by definition of } g \\ &= \max_{\delta \in \{0,1\}^p} \delta^\top |s| - f(\delta) \text{ by maximizing out } w \\ &= \max_{\delta \in [0,1]^p} \delta^\top |s| - f(\delta) \text{ because } F - |s| \text{ is submodular} \end{aligned}$$

## Proof - II

- Notation:  $g : w \mapsto F(\text{supp}(w))$  defined on  $[-1, 1]^p$
- Fenchel dual:  $g^*(s) = \max_{\delta \in [0, 1]^p} \delta^\top |s| - f(\delta)$

## Proof - II

- Notation:  $g : w \mapsto F(\text{supp}(w))$  defined on  $[-1, 1]^p$
- Fenchel dual:  $g^*(s) = \max_{\delta \in [0, 1]^p} \delta^\top |s| - f(\delta)$
- Computation of the **Fenchel bi-dual**, for all  $w$  such that  $\|w\|_\infty \leq 1$ :

$$\begin{aligned} g^{**}(w) &= \max_{s \in \mathbb{R}^p} s^\top w - g^*(s) \\ &= \max_{s \in \mathbb{R}^p} \min_{\delta \in [0, 1]^p} s^\top w - \delta^\top |s| + f(\delta) \\ &= \min_{\delta \in [0, 1]^p} \max_{s \in \mathbb{R}^p} s^\top w - \delta^\top |s| + f(\delta) \text{ by strong duality} \\ &= \min_{\delta \in [0, 1]^p, \delta \geq |w|} f(\delta) = f(|w|) \text{ because } F \text{ is nonincreasing} \end{aligned}$$

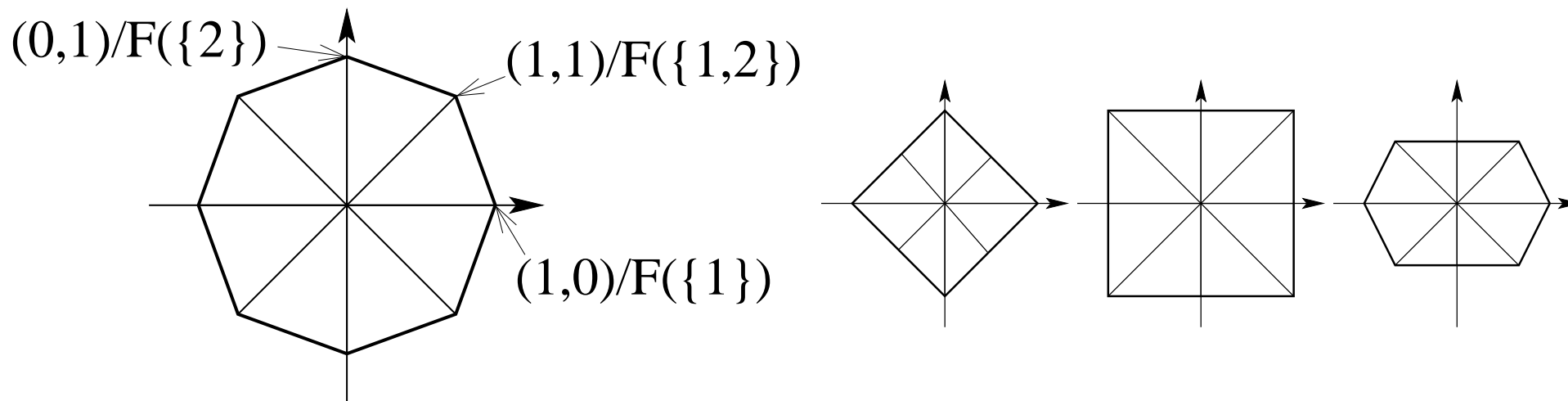
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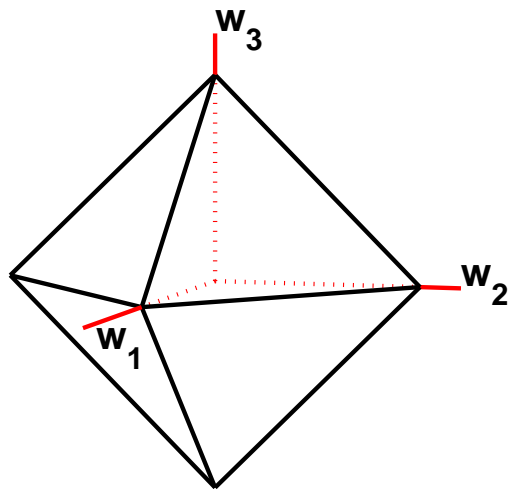
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- **Sparsity-inducing properties:**  $\Omega$  is a **polyhedral** norm



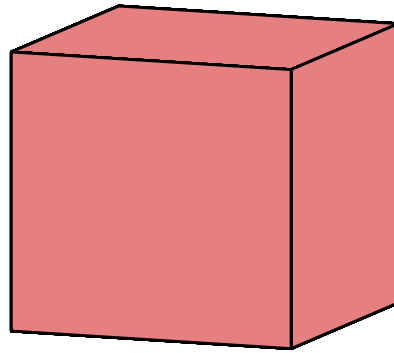
- $A$  is stable if for all  $B \supset A$ ,  $B \neq A \Rightarrow F(B) > F(A)$
- With probability one, stable sets are the only allowed active sets

# Polyhedral unit balls



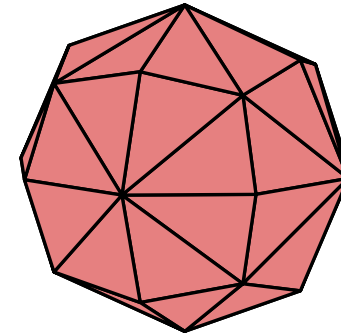
$$F(A) = |A|$$

$$\Omega(w) = \|w\|_1$$



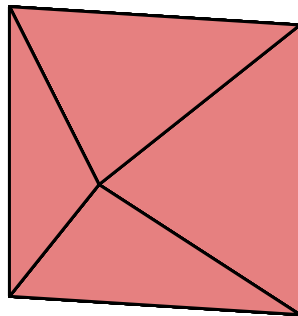
$$F(A) = \min\{|A|, 1\}$$

$$\Omega(w) = \|w\|_\infty$$



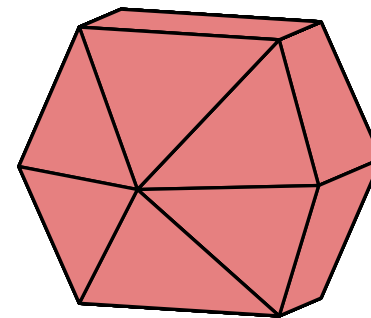
$$F(A) = |A|^{1/2}$$

all possible extreme points



$$F(A) = 1_{\{A \cap \{1\} \neq \emptyset\}} + 1_{\{A \cap \{2,3\} \neq \emptyset\}}$$

$$\Omega(w) = |w_1| + \|w_{\{2,3\}}\|_\infty$$



$$F(A) = 1_{\{A \cap \{1,2,3\} \neq \emptyset\}}$$

$$+ 1_{\{A \cap \{2,3\} \neq \emptyset\}} + 1_{\{A \cap \{3\} \neq \emptyset\}}$$

$$\Omega(w) = \|w\|_\infty + \|w_{\{2,3\}}\|_\infty + |w_3|$$

# Submodular functions and structured sparsity

## Examples

- **From  $\Omega(w)$  to  $F(A)$ :** provides new insights into existing norms

- Grouped norms with **overlapping** groups (Jenatton et al., 2009a)

$$\Omega(w) = \sum_{G \in \mathbf{H}} \|w_G\|_{\infty}$$

- $\ell_1$ - $\ell_{\infty}$  norm  $\Rightarrow$  sparsity at the group level
- Some  $w_G$ 's are set to zero for some groups  $G$

$$(\text{Supp}(w))^c = \bigcup_{G \in \mathbf{H}'} G \text{ for some } \mathbf{H}' \subseteq \mathbf{H}$$

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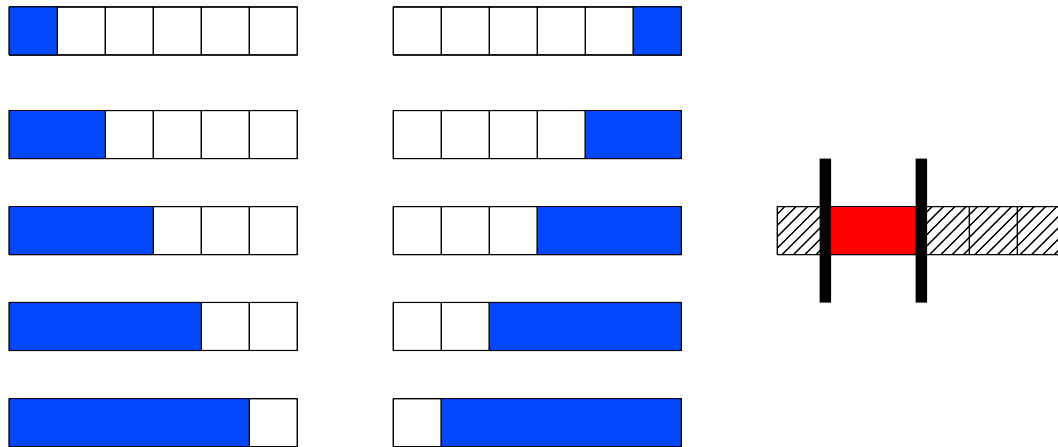
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- Justification not only limited to allowed sparsity patterns

# Selection of contiguous patterns in a sequence

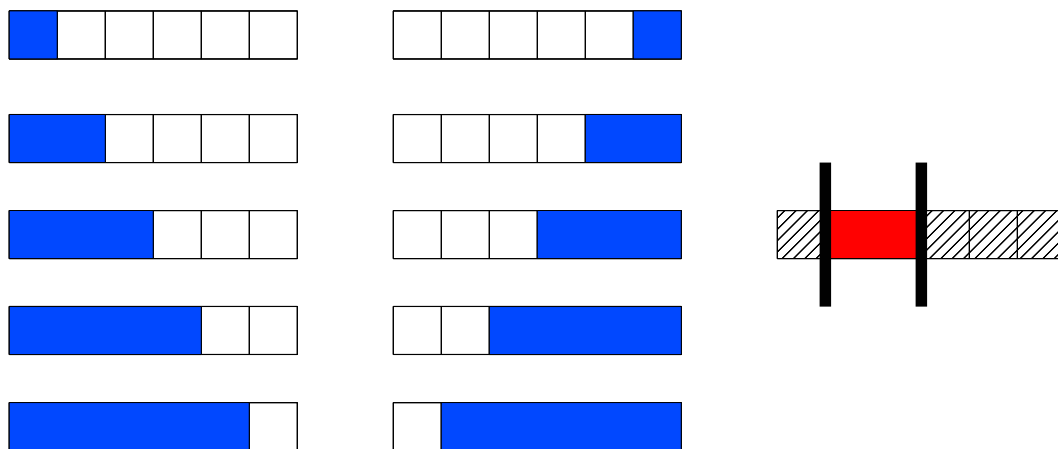
- Selection of contiguous patterns in a sequence



- $\mathbf{H}$  is the set of blue groups: any union of blue groups set to zero leads to the selection of a **contiguous pattern**

# Selection of contiguous patterns in a sequence

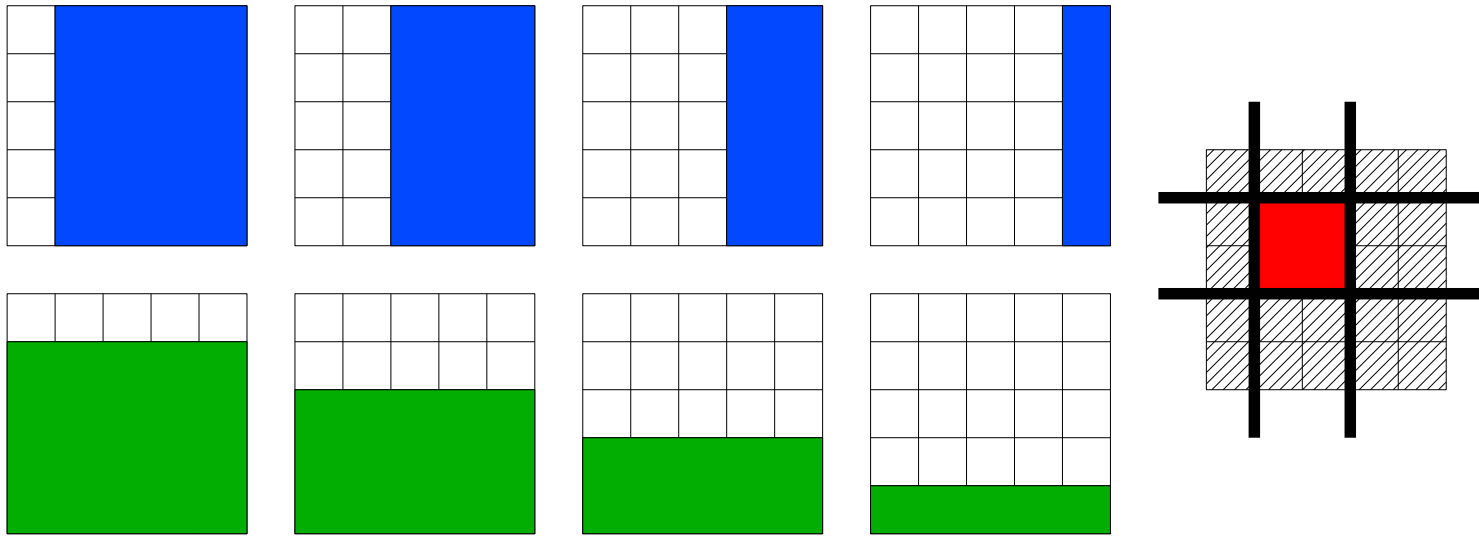
- Selection of contiguous patterns in a sequence



- $\mathbf{H}$  is the set of blue groups: any union of blue groups set to zero leads to the selection of a **contiguous pattern**
- $\sum_{G \in \mathbf{H}} \|w_G\|_{\infty} \Rightarrow F(A) = p - 2 + \text{Range}(A)$  if  $A \neq \emptyset$

# Other examples of set of groups $\mathbf{H}$

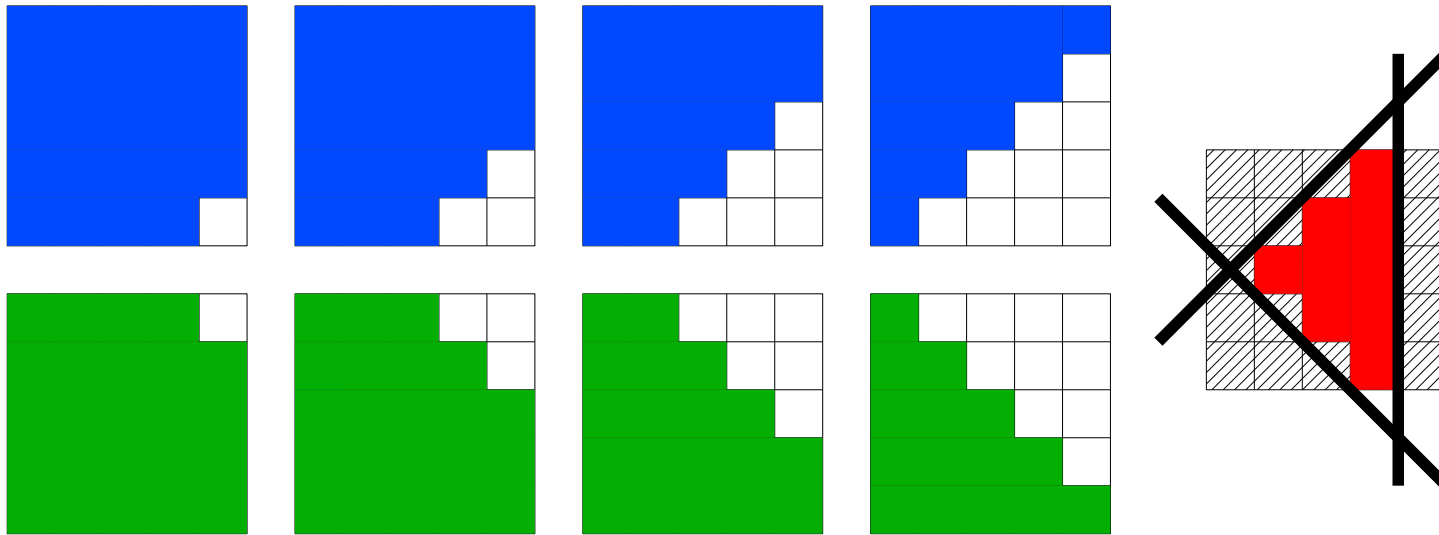
- Selection of rectangles on a 2-D grids,  $p = 25$



- $\mathbf{H}$  is the set of blue/green groups (with their not displayed complements)
- Any union of blue/green groups set to zero leads to the selection of a rectangle

# Other examples of set of groups $H$

- Selection of diamond-shaped patterns on a 2-D grids,  $p = 25$ .

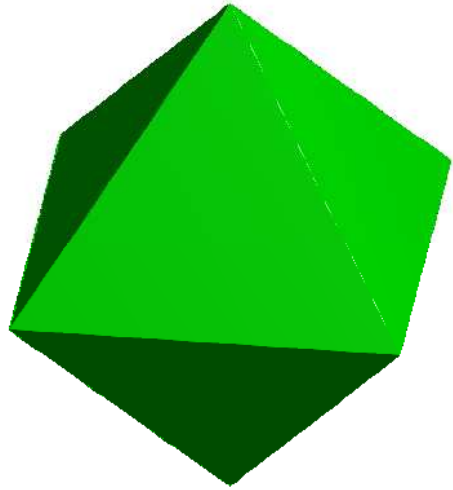


- It is possible to extend such settings to 3-D space, or more complex topologies

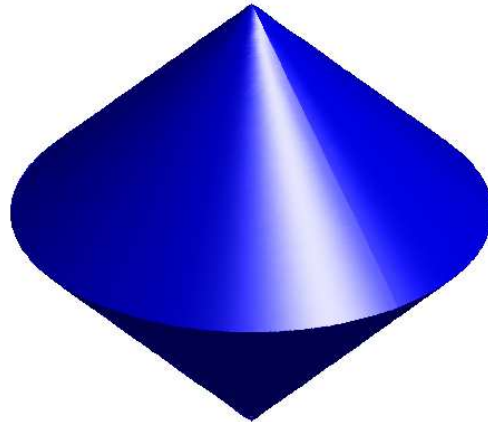


# Unit norm balls

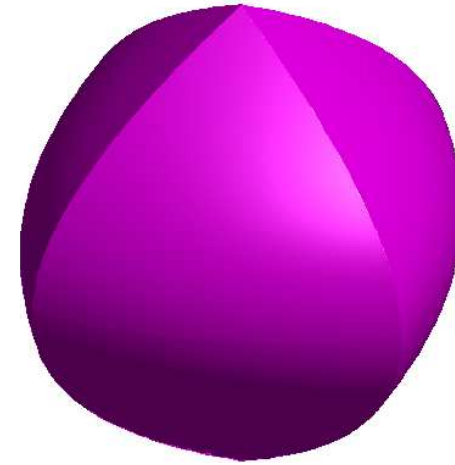
## Geometric interpretation



$$\|w\|_1$$



$$\sqrt{w_1^2 + w_2^2} + |w_3|$$



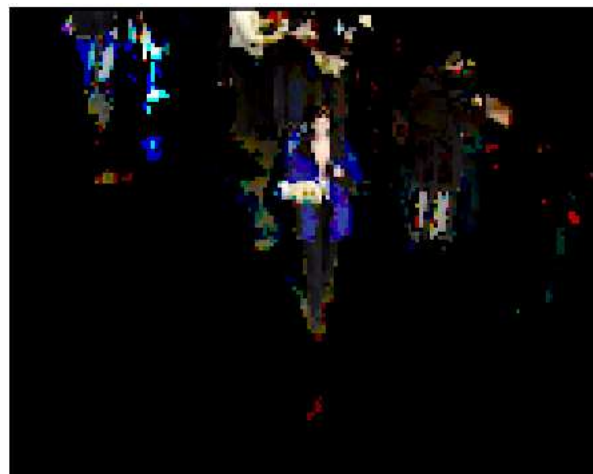
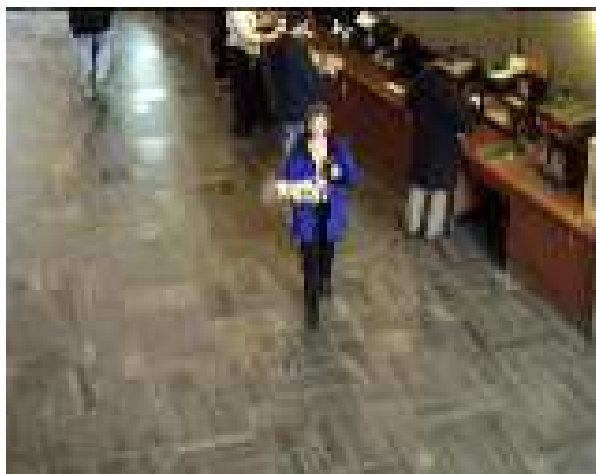
$$\|w\|_2 + |w_1| + |w_2|$$

# Application to background subtraction (Mairal, Jenatton, Obozinski, and Bach, 2010)

Input

$\ell_1$ -norm

Structured norm

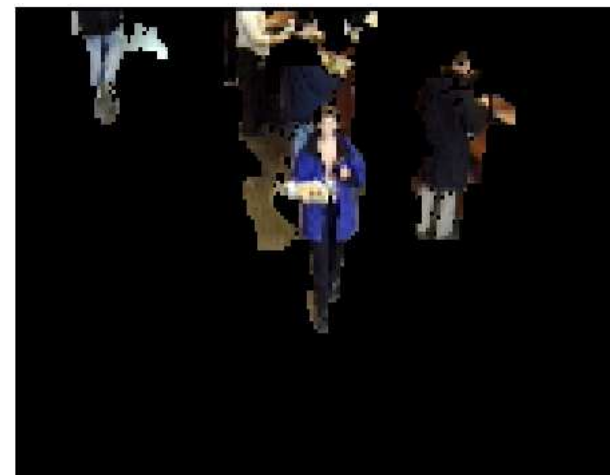
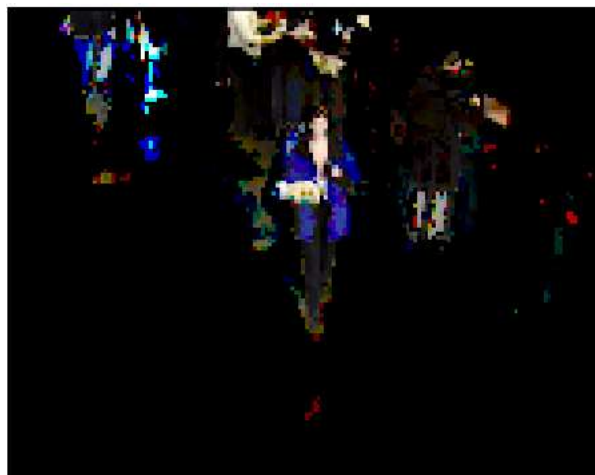


# Application to background subtraction (Mairal, Jenatton, Obozinski, and Bach, 2010)

Background

$\ell_1$ -norm

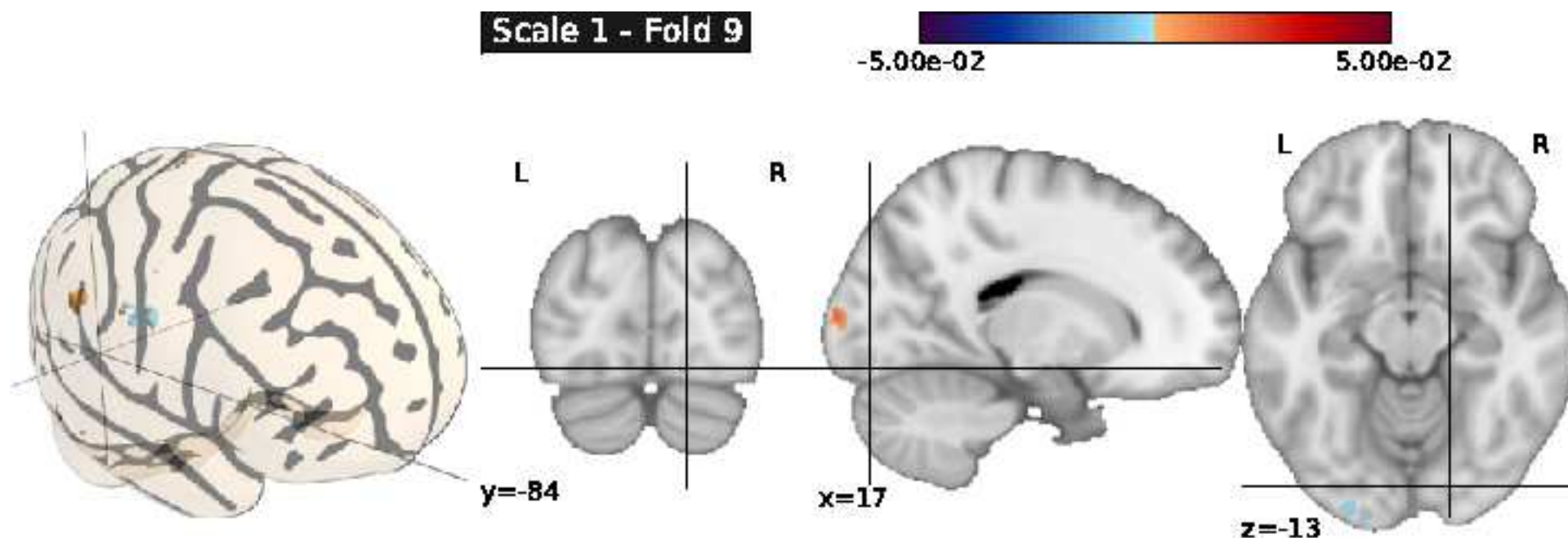
Structured norm



# Application to neuro-imaging

## Structured sparsity for fMRI (Jenatton et al., 2011)

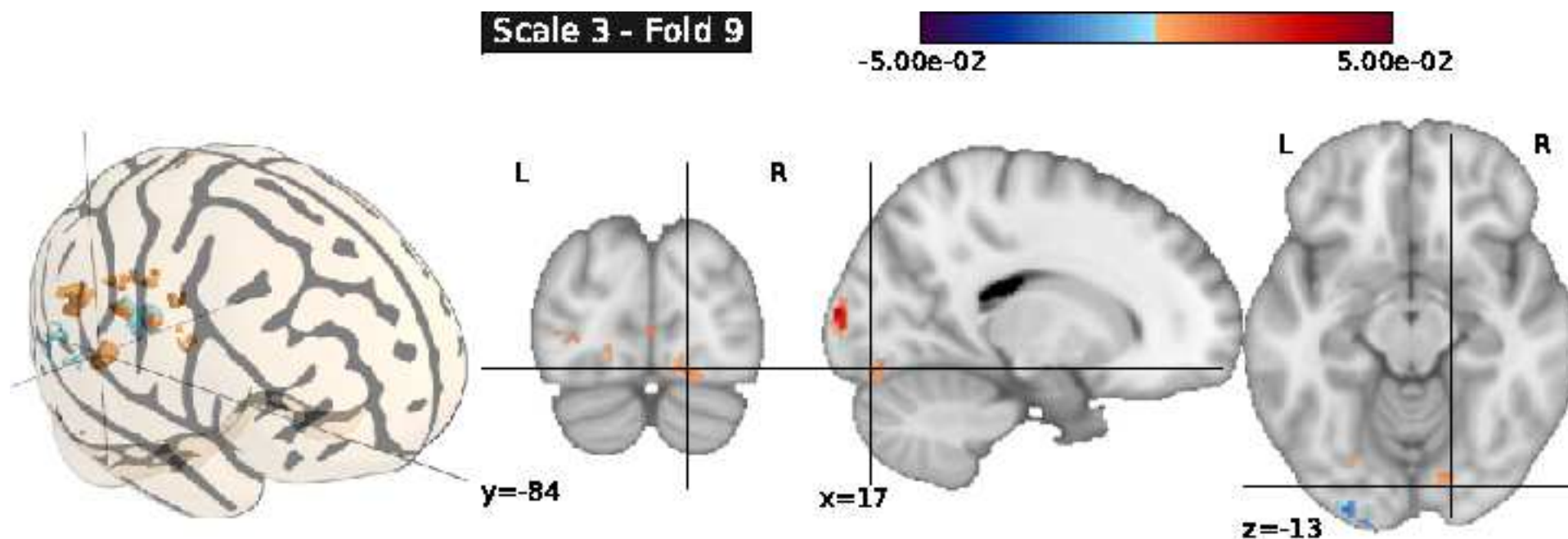
- “Brain reading”: prediction of (seen) object size
- Multi-scale activity levels through hierarchical penalization



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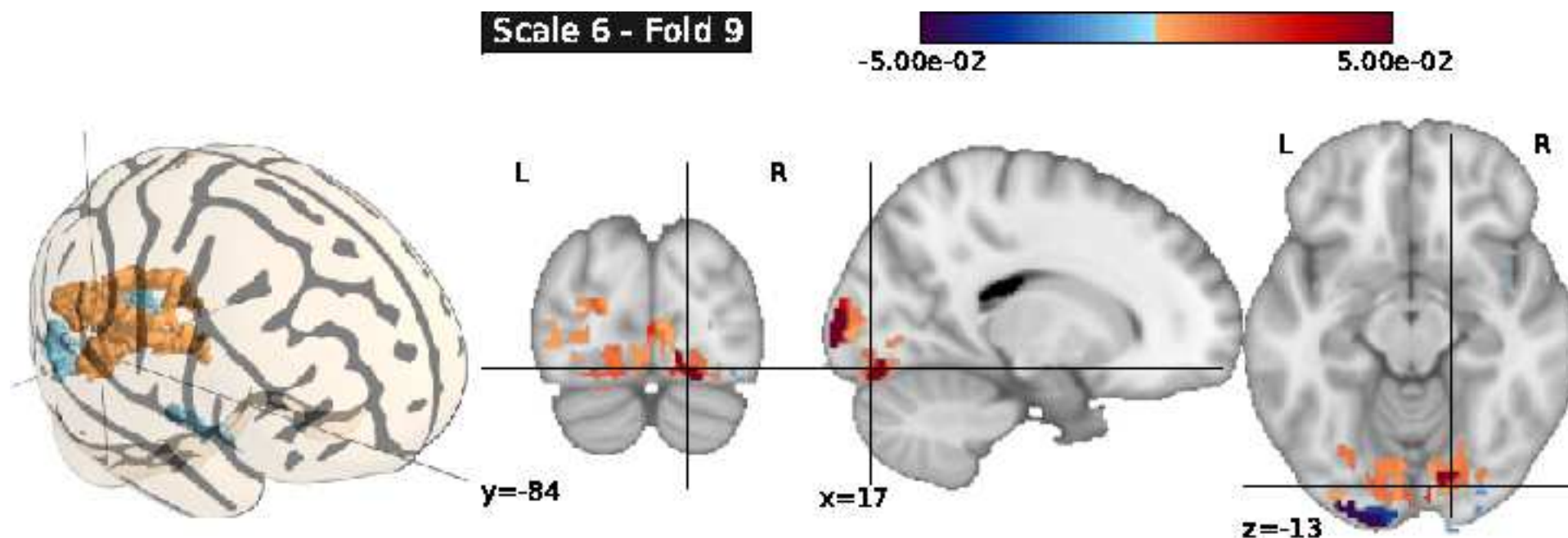




# Application to neuro-imaging

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# Sparse Structured PCA

(Jenatton, Obozinski, and Bach, 2009b)

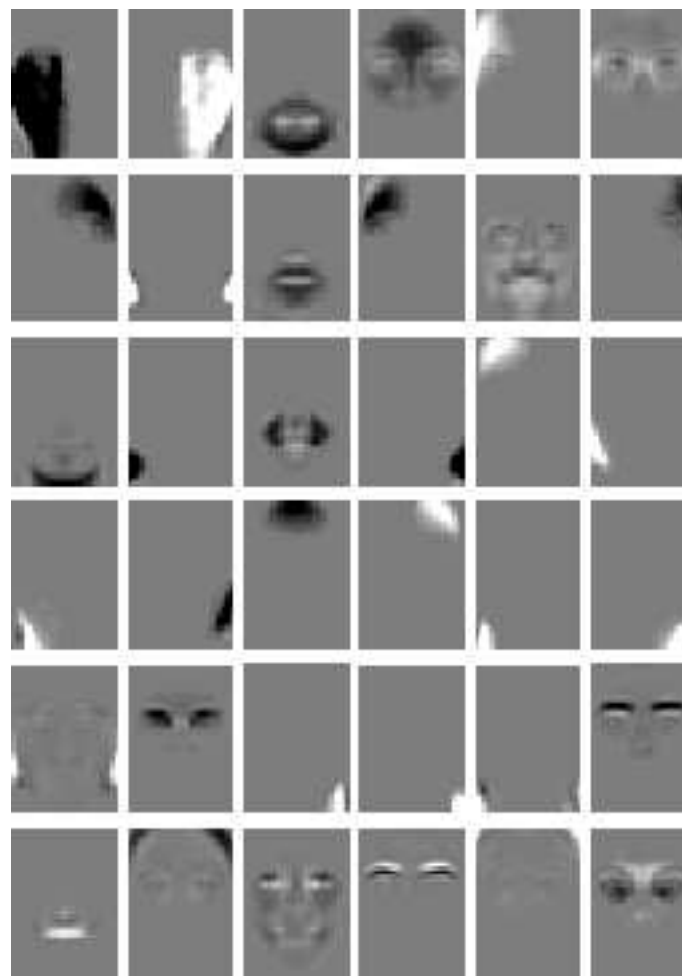
- Learning **sparse and structured dictionary elements**:

$$\min_{W \in \mathbb{R}^{k \times n}, X \in \mathbb{R}^{p \times k}} \frac{1}{n} \sum_{i=1}^n \|y^i - X w^i\|_2^2 + \lambda \sum_{j=1}^p \Omega(x^j) \text{ s.t. } \forall i, \|w^i\|_2 \leq 1$$

## Application to face databases (2/3)



(unstructured) sparse PCA

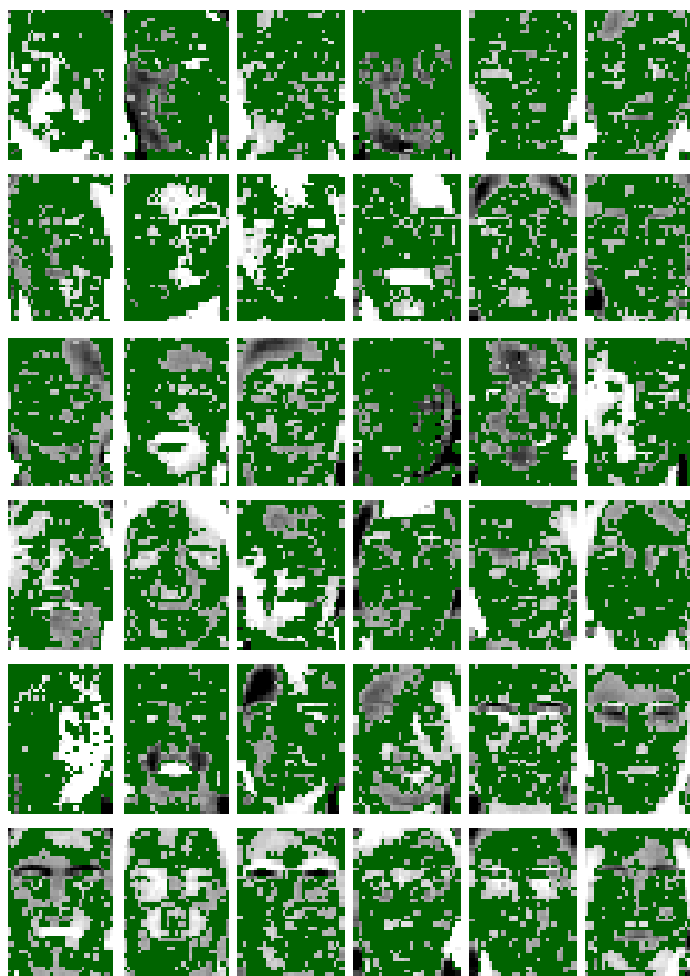


Structured sparse PCA

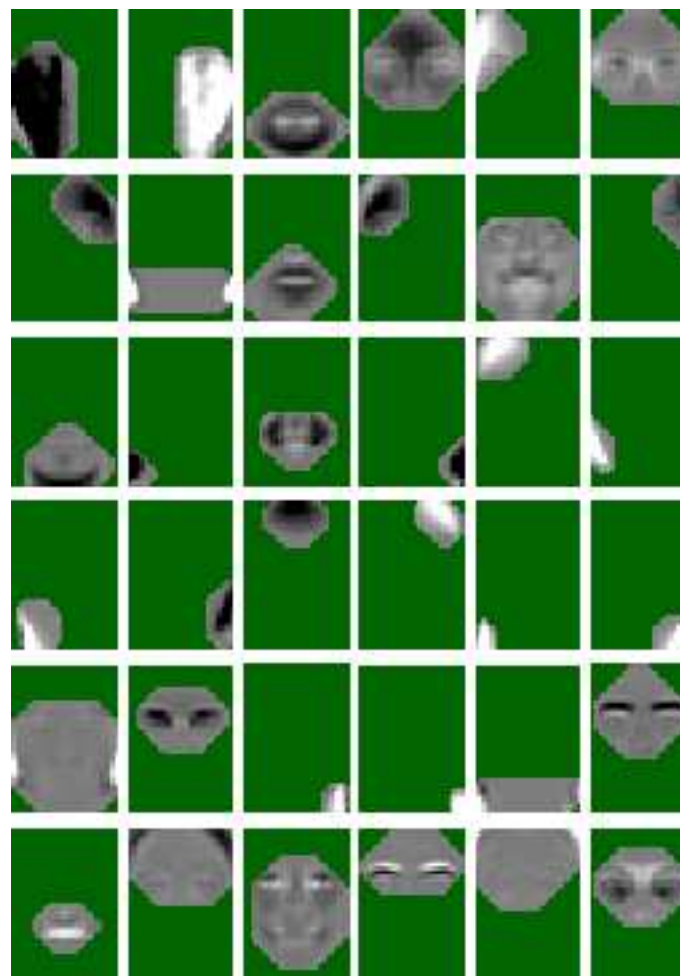
- Enforce selection of **convex** nonzero patterns  $\Rightarrow$  robustness to occlusion



## Application to face databases (2/3)



(unstructured) sparse PCA

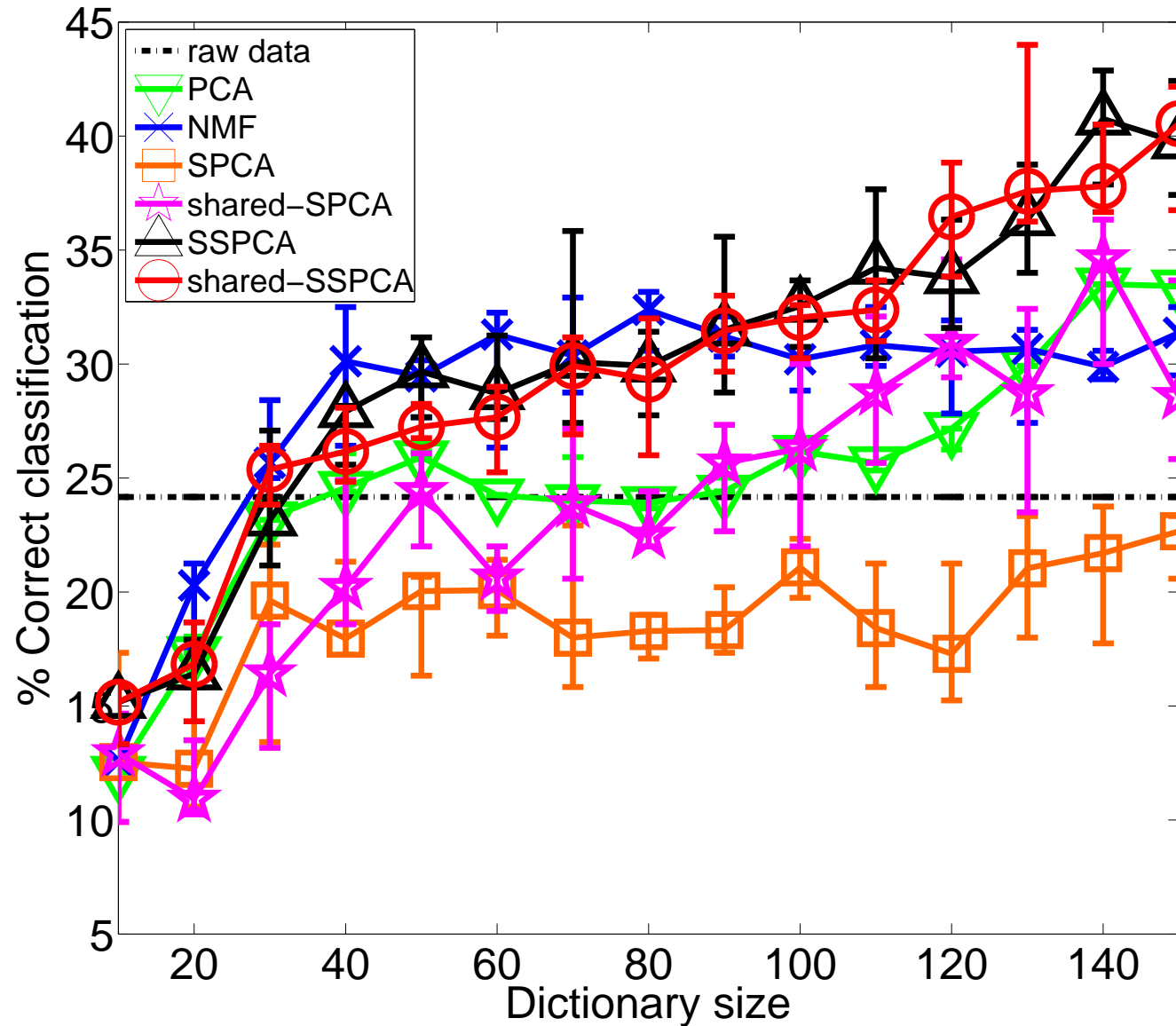


Structured sparse PCA

- Enforce selection of **convex** nonzero patterns  $\Rightarrow$  robustness to occlusion

# Application to face databases (3/3)

- Quantitative performance evaluation on classification task



# Dictionary learning vs. sparse structured PCA

## Exchange roles of $X$ and $w$

- Sparse structured PCA (**structured dictionary elements**):

$$\min_{W \in \mathbb{R}^{k \times n}, X \in \mathbb{R}^{p \times k}} \frac{1}{n} \sum_{i=1}^n \|y^i - X w^i\|_2^2 + \lambda \sum_{j=1}^k \Omega(x^j) \text{ s.t. } \forall i, \|w^i\|_2 \leq 1.$$

- Dictionary learning with **structured sparsity for codes**  $w$ :

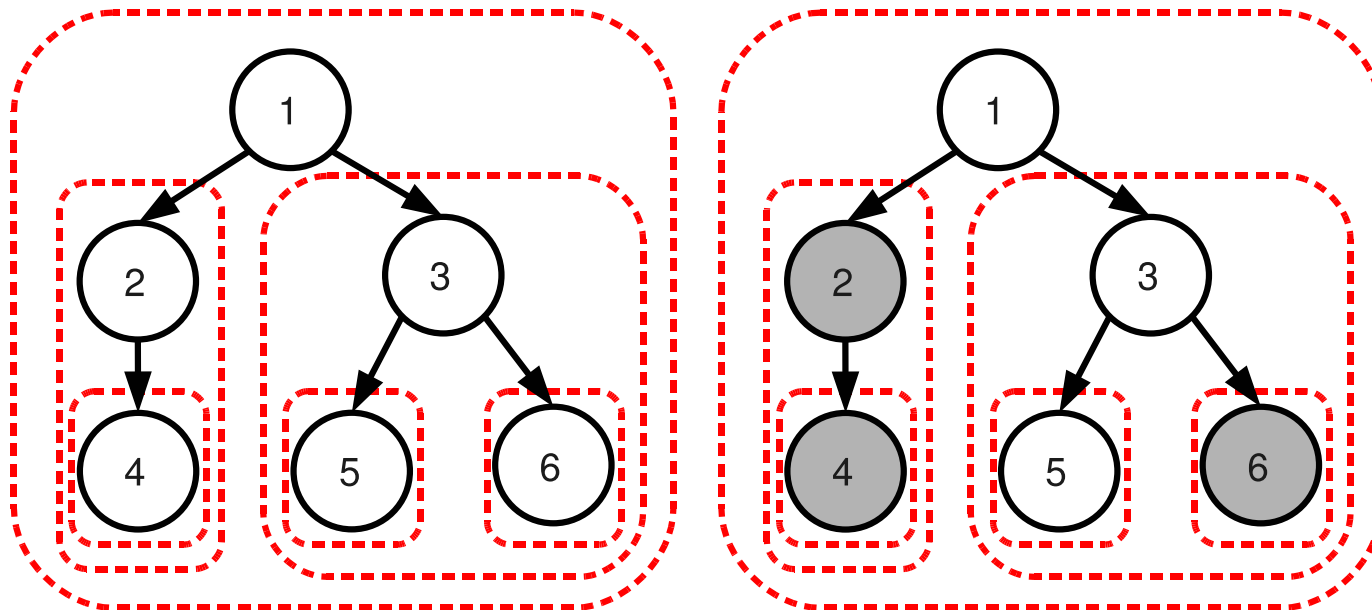
$$\min_{W \in \mathbb{R}^{k \times n}, X \in \mathbb{R}^{p \times k}} \frac{1}{n} \sum_{i=1}^n \|y^i - X w^i\|_2^2 + \lambda \Omega(w^i) \text{ s.t. } \forall j, \|x^j\|_2 \leq 1.$$

- **Optimization: proximal methods**

- Requires solving many times  $\min_{w \in \mathbb{R}^p} \frac{1}{2} \|y - w\|_2^2 + \lambda \Omega(w)$
- **Modularity of implementation** if proximal step is efficient (Jenatton et al., 2010; Mairal et al., 2010)

# Hierarchical dictionary learning (Jenatton, Mairal, Obozinski, and Bach, 2010)

- Structure on codes  $w$  (not on dictionary  $X$ )
- Hierarchical penalization:  $\Omega(w) = \sum_{G \in \mathbf{H}} \|w_G\|_\infty$  where groups  $G$  in  $\mathbf{H}$  are equal to **set of descendants** of some nodes in a tree



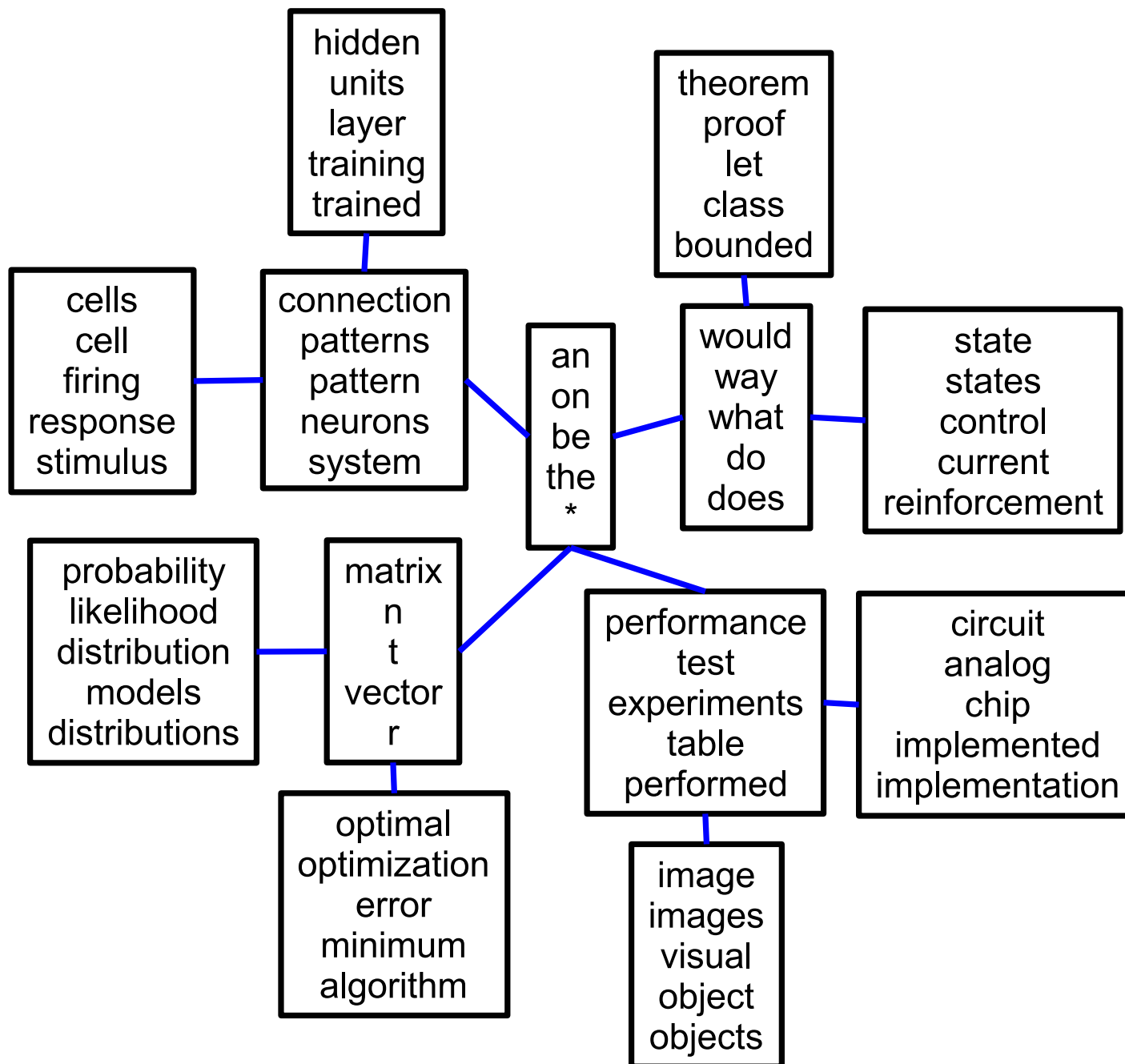
- Variable selected after its ancestors (Zhao et al., 2009; Bach, 2008)

# Hierarchical dictionary learning

## Modelling of text corpora

- Each document is modelled through word counts
- Low-rank matrix factorization of word-document matrix
- Probabilistic topic models (Blei et al., 2003)
  - Similar structures based on non parametric Bayesian methods (Blei et al., 2004)
  - **Can we achieve similar performance with simple matrix factorization formulation?**

# Modelling of text corpora - Dictionary tree



# Submodular functions and structured sparsity

## Examples

- **From  $\Omega(w)$  to  $F(A)$ :** provides new insights into existing norms

- Grouped norms with **overlapping** groups (Jenatton et al., 2009a)

$$\Omega(w) = \sum_{G \in \mathbf{H}} \|w_G\|_{\infty} \quad \Rightarrow \quad F(A) = \text{Card}(\{G \in \mathbf{H}, G \cap A \neq \emptyset\})$$

- Justification not only limited to allowed sparsity patterns

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- Justification not only limited to allowed sparsity patterns

- **From  $F(A)$  to  $\Omega(w)$ :** provides new sparsity-inducing norms

- $F(A) = g(\text{Card}(A)) \Rightarrow \Omega$  is a combination of **order statistics**

- **Non-factorial priors** for supervised learning:  $\Omega$  depends on the eigenvalues of  $X_A^{\top} X_A$  and not simply on the cardinality of  $A$



# Unified optimization algorithms

- **Polyhedral norm** with  $O(3^p)$  faces and extreme points
  - Not suitable to linear programming toolboxes
- **Subgradient** ( $w \mapsto \Omega(w)$  non-differentiable)
  - subgradient may be obtained in polynomial time  $\Rightarrow$  too slow

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  - subgradient may be obtained in polynomial time  $\Rightarrow$  too slow
- **Proximal methods** (e.g., Beck and Teboulle, 2009)
  - $\min_{w \in \mathbb{R}^p} L(y, Xw) + \lambda\Omega(w)$ : differentiable + non-differentiable
  - Efficient when  $(P) : \min_{w \in \mathbb{R}^p} \frac{1}{2}\|w - v\|_2^2 + \lambda\Omega(w)$  is “easy”
  - **Fact:**  $(P)$  is equivalent to submodular function minimization

# Optimization for sparsity-inducing norms

(see Bach, Jenatton, Mairal, and Obozinski, 2011)

- Gradient descent as a **proximal method** (differentiable functions)

$$- w_{t+1} = \arg \min_{w \in \mathbb{R}^p} L(w_t) + (w - w_t)^\top \nabla L(w_t) + \frac{B}{2} \|w - w_t\|_2^2$$

$$- w_{t+1} = w_t - \frac{1}{B} \nabla L(w_t)$$

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- Problems of the form:

$$\min_{w \in \mathbb{R}^p} L(w) + \lambda \Omega(w)$$

$$- w_{t+1} = \arg \min_{w \in \mathbb{R}^p} L(w_t) + (w - w_t)^\top \nabla L(w_t) + \lambda \Omega(w) + \frac{B}{2} \|w - w_t\|_2^2$$

$$- \Omega(w) = \|w\|_1 \Rightarrow \text{Thresholded gradient descent}$$

- Similar convergence rates than smooth optimization

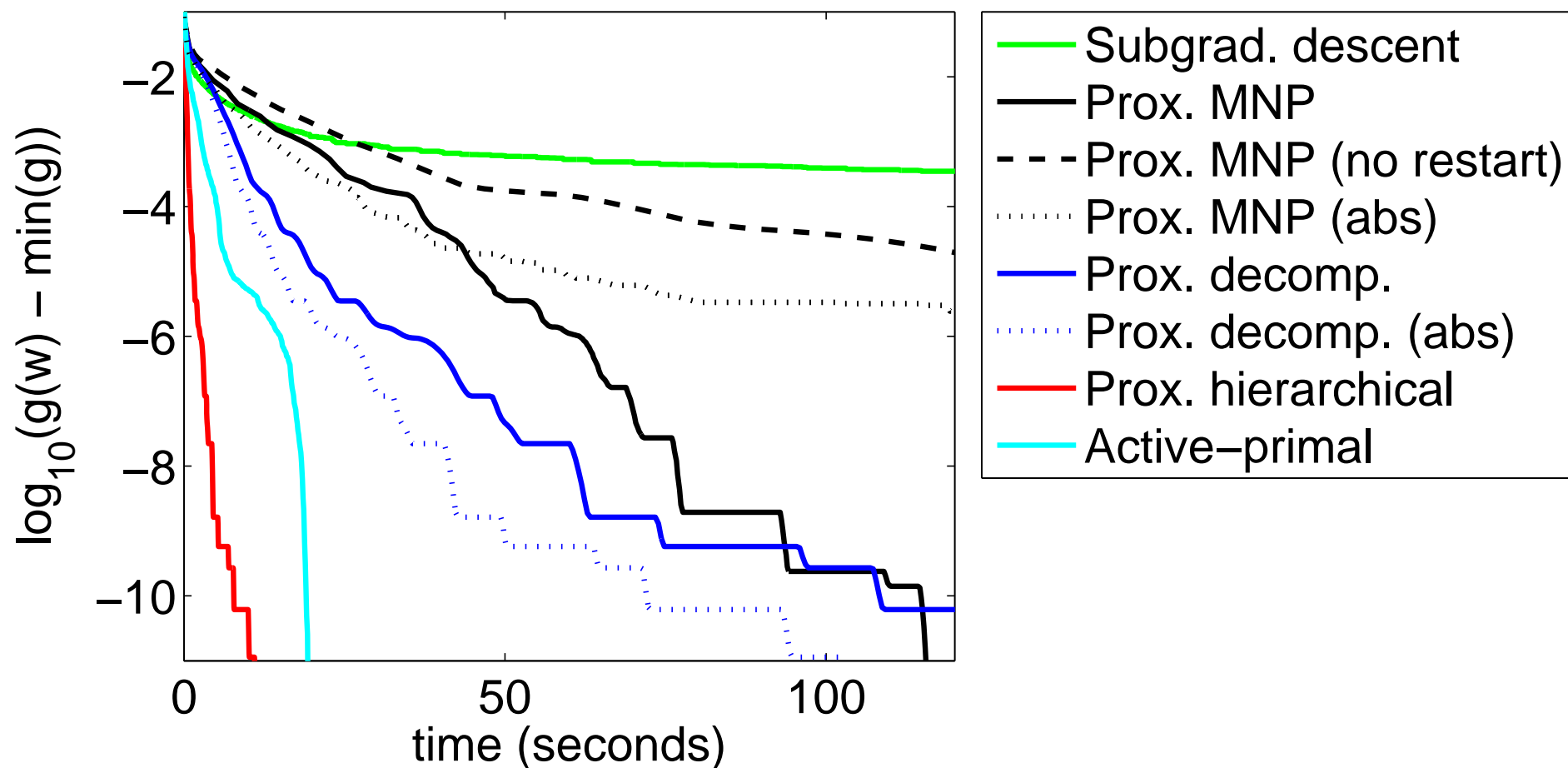
- Acceleration methods (Nesterov, 2007; Beck and Teboulle, 2009)

# Unified optimization algorithms

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  - **Fact:**  $(P)$  is equivalent to submodular function minimization
- **Active-set methods**

# Comparison of optimization algorithms

- Tree-based regularization ( $p = 511$ )
- See Bach et al. (2011) for larger-scale problems



# Unified theoretical analysis

- **Decomposability**

- Key to theoretical analysis (Negahban et al., 2009)
- **Property:**  $\forall w \in \mathbb{R}^p$ , and  $\forall J \subset V$ , if  $\min_{j \in J} |w_j| \geq \max_{j \in J^c} |w_j|$ , then  $\Omega(w) = \Omega_J(w_J) + \Omega^{J^c}(w_{J^c})$

- **Support recovery**

- Extension of known sufficient condition (Zhao and Yu, 2006; Negahban and Wainwright, 2008)

- **High-dimensional inference**

- Extension of known sufficient condition (Bickel et al., 2009)
- Matches with analysis of Negahban et al. (2009) for common cases

# Support recovery - $\min_{w \in \mathbb{R}^p} \frac{1}{2n} \|y - Xw\|_2^2 + \lambda \Omega(w)$

## • Notation

- $\rho(J) = \min_{B \subset J^c} \frac{F(B \cup J) - F(J)}{F(B)} \in (0, 1]$  (for  $J$  stable)
- $c(J) = \sup_{w \in \mathbb{R}^p} \Omega_J(w_J) / \|w_J\|_2 \leq |J|^{1/2} \max_{k \in V} F(\{k\})$

## • Proposition

- Assume  $y = Xw^* + \sigma\varepsilon$ , with  $\varepsilon \sim \mathcal{N}(0, I)$
- $J =$  smallest stable set containing the support of  $w^*$
- Assume  $\nu = \min_{j, w_j^* \neq 0} |w_j^*| > 0$
- Let  $Q = \frac{1}{n} X^\top X \in \mathbb{R}^{p \times p}$ . Assume  $\kappa = \lambda_{\min}(Q_{JJ}) > 0$
- Assume that for  $\eta > 0$ , 
$$(\Omega^J)^* [(\Omega_J(Q_{JJ}^{-1} Q_{Jj}))_{j \in J^c}] \leq 1 - \eta$$
- If  $\lambda \leq \frac{\kappa\nu}{2c(J)}$ ,  $\hat{w}$  has support equal to  $J$ , with probability larger than 
$$1 - 3P\left(\Omega^*(z) > \frac{\lambda\eta\rho(J)\sqrt{n}}{2\sigma}\right)$$
- $z$  is a multivariate normal with covariance matrix  $Q$



# Consistency - $\min_{w \in \mathbb{R}^p} \frac{1}{2n} \|y - Xw\|_2^2 + \lambda \Omega(w)$

## • Proposition

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- $J =$  smallest stable set containing the support of  $w^*$
- Let  $Q = \frac{1}{n} X^\top X \in \mathbb{R}^{p \times p}$ .
- Assume that  $\forall \Delta$  s.t.  $\Omega^J(\Delta_{J^c}) \leq 3\Omega_J(\Delta_J)$ ,  $\Delta^\top Q \Delta \geq \kappa \|\Delta_J\|_2^2$

– Then  $\Omega(\hat{w} - w^*) \leq \frac{24c(J)^2 \lambda}{\kappa \rho(J)^2}$  and  $\frac{1}{n} \|X\hat{w} - Xw^*\|_2^2 \leq \frac{36c(J)^2 \lambda^2}{\kappa \rho(J)^2}$

with probability larger than  $1 - P(\Omega^*(z) > \frac{\lambda \rho(J) \sqrt{n}}{2\sigma})$

- $z$  is a multivariate normal with covariance matrix  $Q$

## • Concentration inequality ( $z$ normal with covariance matrix $Q$ ):

- $\mathcal{T}$  set of stable inseparable sets
- Then  $P(\Omega^*(z) > t) \leq \sum_{A \in \mathcal{T}} 2^{|A|} \exp\left(-\frac{t^2 F(A)^2 / 2}{1^\top Q_{AA} 1}\right)$

# Symmetric submodular functions (Bach, 2011)

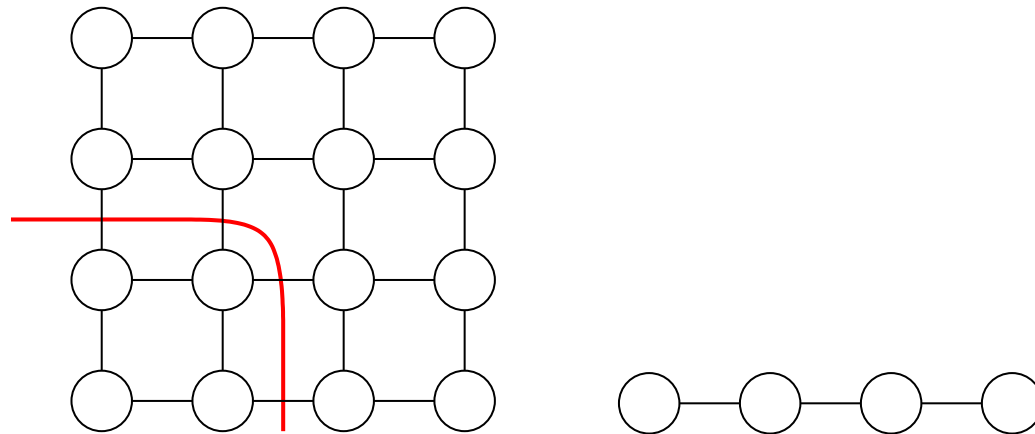
- Let  $F : 2^V \rightarrow \mathbb{R}$  be a symmetric submodular set-function
- **Proposition:** The Lovász extension  $f(w)$  is the convex envelope of the function  $w \mapsto \max_{\alpha \in \mathbb{R}} F(\{w \geq \alpha\})$  on the set  $[0, 1]^p + \mathbb{R}1_V = \{w \in \mathbb{R}^p, \max_{k \in V} w_k - \min_{k \in V} w_k \leq 1\}$ .
- Shaping all level sets

# Symmetric submodular functions - Examples

- From  $\Omega(w)$  to  $F(A)$ : provides new insights into existing norms

- Cuts - total variation

$$F(A) = \sum_{k \in A, j \in V \setminus A} d(k, j) \Rightarrow f(w) = \sum_{k, j \in V} d(k, j)(w_k - w_j)_+$$

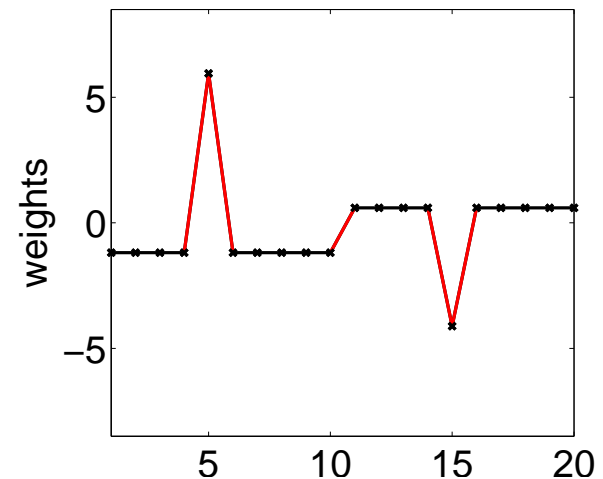
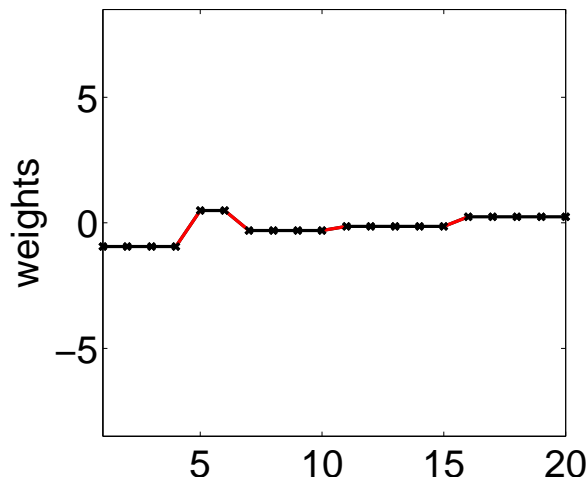
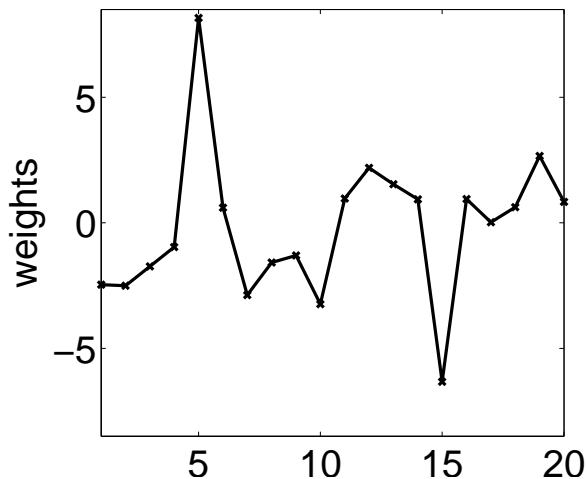
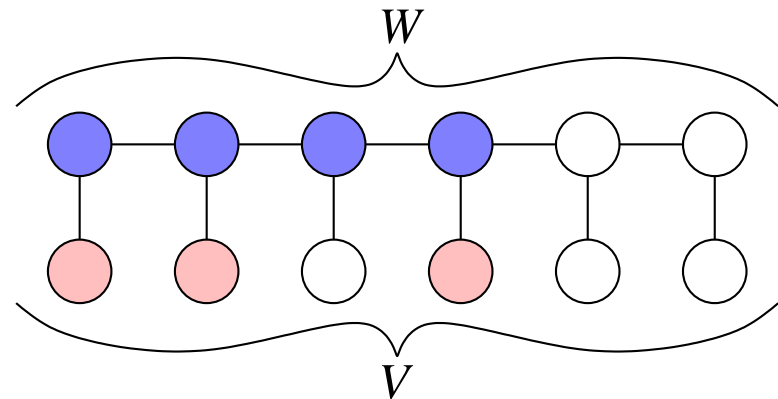


- NB: graph may be directed
- Application to change-point detection (Tibshirani et al., 2005; Harchaoui and Lévy-Leduc, 2008)

# Symmetric submodular functions - Examples

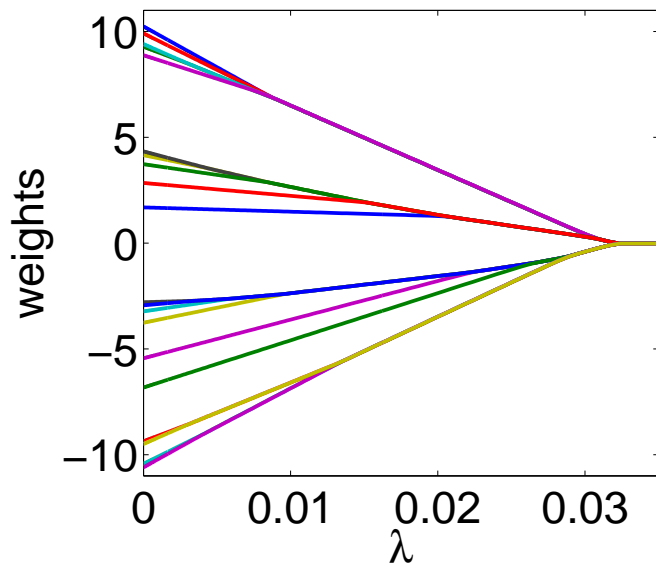
- From  $F(A)$  to  $\Omega(w)$ : provides new sparsity-inducing norms
  - Regular functions (Boykov et al., 2001; Chambolle and Darbon, 2009)

$$F(A) = \min_{B \subset W} \sum_{k \in B, j \in W \setminus B} d(k, j) + \lambda |A \Delta B|$$

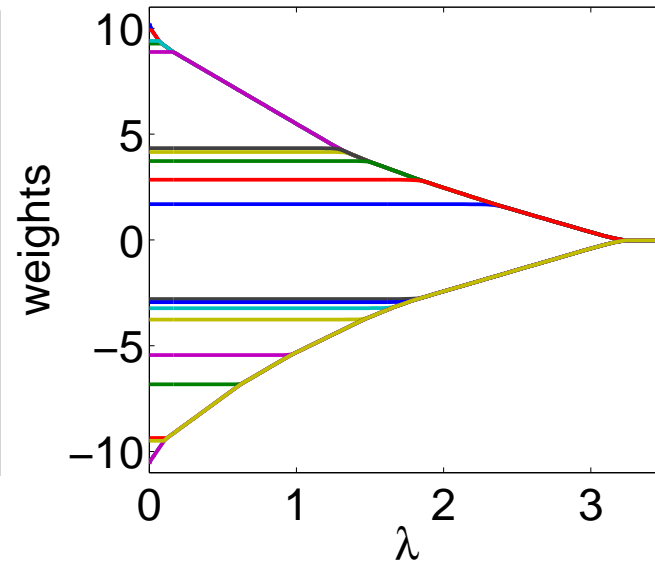


# Symmetric submodular functions - Examples

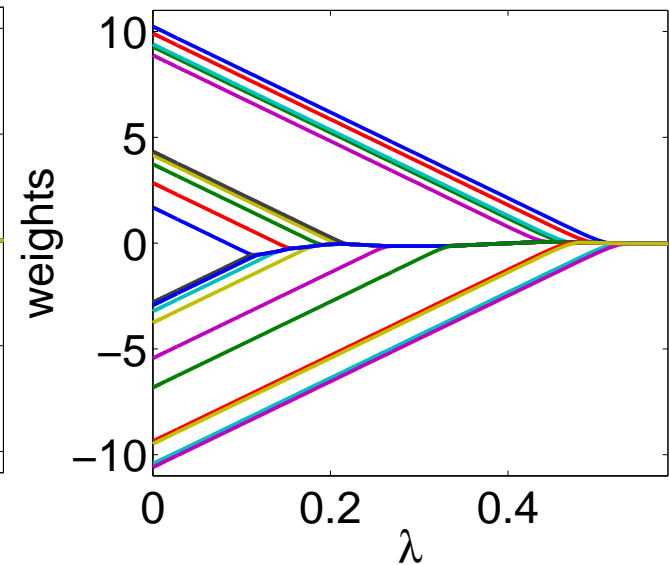
- From  $F(A)$  to  $\Omega(w)$ : provides new sparsity-inducing norms
  - $F(A) = g(\text{Card}(A)) \Rightarrow$  priors on the size and numbers of clusters



$$|A|(p - |A|)$$



$$1_{|A| \in (0, p)}$$



$$\max\{|A|, p - |A|\}$$

- Convex formulations for clustering (Hocking, Joulin, Bach, and Vert, 2011)

# $\ell_2$ -relaxation of combinatorial penalties (Obozinski and Bach, 2012)

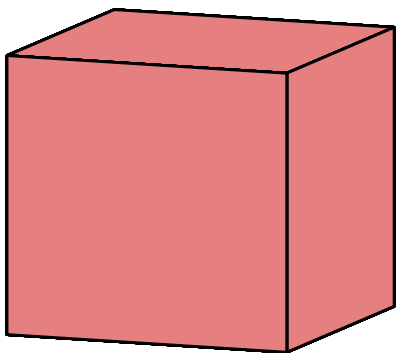
- **Main result** of Bach (2010):

- $f(|w|)$  is the convex envelope of  $F(\text{Supp}(w))$  on  $[-1, 1]^p$

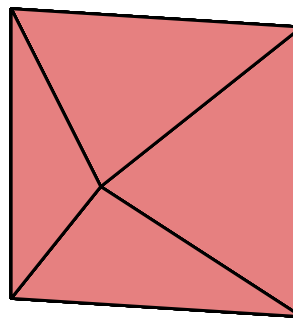
- **Problems:**

- Limited to submodular functions

- Limited to  $\ell_\infty$ -relaxation: undesired artefacts



$$F(A) = \min\{|A|, 1\}$$
$$\Omega(w) = \|w\|_\infty$$



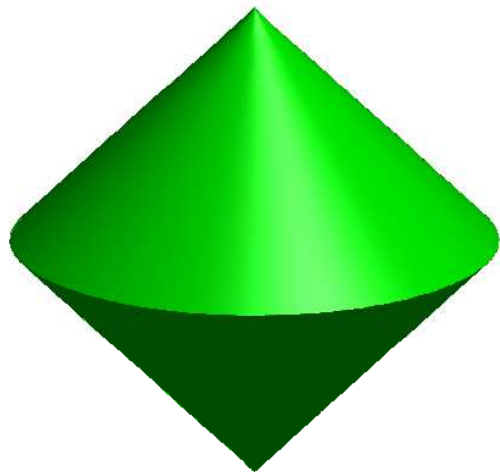
$$F(A) = 1_{\{A \cap \{1\} \neq \emptyset\}} + 1_{\{A \cap \{2,3\} \neq \emptyset\}}$$
$$\Omega(w) = |w_1| + \|w_{\{2,3\}}\|_\infty$$

# $\ell_2$ -relaxation of **submodular** penalties (Obozinski and Bach, 2012)

- $F$  a nondecreasing submodular function with Lovász extension  $f$
- Define  $\Omega_2(w) = \min_{\eta \in \mathbb{R}_+^p} \frac{1}{2} \sum_{i \in V} \frac{|w_i|^2}{\eta_i} + \frac{1}{2} f(\eta)$ 
  - NB: general formulation (Micchelli et al., 2011; Bach et al., 2011)
- **Proposition 1:**  $\Omega_2$  is the convex envelope of  $w \mapsto F(\text{Supp}(w)) \|w\|_2$
- **Proposition 2:**  $\Omega_2$  is the *homogeneous* convex envelope of  $w \mapsto \frac{1}{2} F(\text{Supp}(w)) + \frac{1}{2} \|w\|_2^2$
- **Jointly penalizing and regularizing**
  - Extension possible to  $\ell_q$ ,  $q > 1$

## From $l_\infty$ to $l_2$

### Removal of undesired artefacts



$$F(A) = 1_{\{A \cap \{3\} \neq \emptyset\}} + 1_{\{A \cap \{1,2\} \neq \emptyset\}}$$

$$\Omega_2(w) = |w_3| + \|w_{\{1,2\}}\|_2$$



$$F(A) = 1_{\{A \cap \{1,2,3\} \neq \emptyset\}} \\ + 1_{\{A \cap \{2,3\} \neq \emptyset\}} + 1_{\{A \cap \{2\} \neq \emptyset\}}$$

- Extension to non-submodular functions + tightness study: see Obozinski and Bach (2012)



# Beyond submodular functions?

- Let  $F$  be **any** set-function
- **“Edmonds extension”**: the convex envelope of  $w \mapsto F(\text{Supp}(w))$  on  $[0, 1]^p$  is equal to

$$f(w) = \sup_{\forall A \subseteq V, s(A) \leq F(A)} w^\top s = \sup_{s \in P(F)} w^\top s$$

- When is it an extension of  $F$ ?
- **Lower combinatorial envelope**:  $G(B) = f(1_B) = \sup_{s \in P(F)} s(B)$ 
  - $G \leq F$
  - Property: idempotent operation
- **A new class of set-functions**: functions for which  $G = F$

# Conclusion

- **Structured sparsity for machine learning and statistics**
    - Many applications (image, audio, text, etc.)
    - May be achieved through structured sparsity-inducing norms
    - Link with submodular functions: unified analysis and algorithms
- Submodular functions to encode discrete structures**

# Conclusion

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- Many applications (image, audio, text, etc.)
- May be achieved through structured sparsity-inducing norms
- Link with submodular functions: unified analysis and algorithms

- **Submodular functions to encode discrete structures**

- **On-going work on machine learning and submodularity**

- Improved complexity bounds for submodular function minimization
- Submodular function maximization
- Importing concepts from machine learning (e.g., graphical models)
- Multi-way partitions for computer vision
- Online learning
- Going beyond linear programming duality?

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