Optimal solutions

for Sparse Principal Component Analysis

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Introduction

Principal Component Analysis

- Classic dimensionality reduction tool.
- Numerically cheap: $O(n^2)$ as it only requires computing a few dominant eigenvectors.

Sparse PCA

- Get **sparse** factors capturing a maximum of variance.
- Numerically hard: combinatorial problem.
- Controlling the sparsity of the solution is also hard in practice.

Introduction



Clustering of the gene expression data in the PCA versus sparse PCA basis with 500 genes. The factors f on the left are dense and each use all 500 genes while the sparse factors g_1 , g_2 and g_3 on the right involve 6, 4 and 4 genes respectively. (Data: Iconix Pharmaceuticals)

Introduction

Principal Component Analysis. Given a (centered) data set $A \in \mathbb{R}^{n \times m}$ composed of m observations on n variables, we form the covariance matrix $C = A^T A / (m - 1)$ and solve:

 $\begin{array}{ll} \text{maximize} & x^T C x\\ \text{subject to} & \|x\| = 1, \end{array}$

in the variable $x \in \mathbf{R}^n$, i.e. we maximize the variance explained by the factor x.

Sparse Principal Component Analysis. We constrain the cardinality of the factor x and solve:

maximize $x^T C x$ subject to Card(x) = k||x|| = 1,

in the variable $x \in \mathbf{R}^n$, where $\mathbf{Card}(x)$ is the number of nonzero coefficients in the vector x and k > 0 is a parameter controlling sparsity.

Outline

- Introduction
- Algorithms
- Optimality
- Numerical Results

Existing methods. . .

- Cadima & Jolliffe (1995): the loadings with small absolute value are thresholded to zero.
- SPCA Zou, Hastie & Tibshirani (2006), non-convex algo. based on a l_1 penalized representation of PCA as a regression problem.
- A convex relaxation in d'Aspremont, El Ghaoui, Jordan & Lanckriet (2007).
- Non-convex optimization methods: SCoTLASS by Jolliffe, Trendafilov & Uddin (2003) or Sriperumbudur, Torres & Lanckriet (2007).
- A greedy algorithm by Moghaddam, Weiss & Avidan (2006b).

Simplest solution: just sort variables according to variance, keep the k variables with highest variance. **Schur-Horn theorem**: the diagonal of a matrix majorizes its eigenvalues.



Other simple solution: **Thresholding**, compute the first factor x from regular PCA and keep the k variables corresponding to the k largest coefficients.

Greedy search (see Moghaddam et al. (2006b)). Written on the square root here.

- 1. Preprocessing. Permute elements of Σ accordingly so that its diagonal is decreasing. Compute the Cholesky decomposition $\Sigma = A^T A$. Initializate $I_1 = \{1\}$ and $x_1 = a_1/||a_1||$.
- 2. Compute

$$i_k = \operatorname*{argmax}_{i \notin I_k} \lambda_{max} \left(\sum_{j \in I_k \cup \{i\}} a_j a_j^T \right)$$

3. Set $I_{k+1} = I_k \cup \{i_k\}$.

4. Compute x_{k+1} as the dominant eigenvector of $\sum_{j \in I_{k+1}} a_j a_j^T$.

5. Set k = k + 1. If k < n go back to step 2.

Algorithms: complexity

Greedy Search

- Iteration k of the greedy search requires computing (n k) maximum eigenvalues, hence has complexity $O((n k)k^2)$ if we exploit the Gram structure.
- This means that computing a full path of solutions has complexity $O(n^4)$.

Approximate Greedy Search

• We can exploit the following first-order inequality:

$$\lambda_{max}\left(\sum_{j\in I_k\cup\{i\}}a_ja_j^T\right) \ge \lambda_{max}\left(\sum_{j\in I_k}a_ja_j^T\right) + (a_i^Tx_k)^2$$

where x_k is the dominant eigenvector of $\sum_{j \in I_k} a_j a_j^T$.

• We only need to solve one maximum eigenvalue problem per iteration, with cost $O(k^2)$. The complexity of computing a full path of solution is now $O(n^3)$.

Approximate greedy search.

- 1. Preprocessing. Permute elements of Σ accordingly so that its diagonal is decreasing. Compute the Cholesky decomposition $\Sigma = A^T A$. Initializate $I_1 = \{1\}$ and $x_1 = a_1/||a_1||$.
- 2. Compute $i_k = \operatorname{argmax}_{i \notin I_k} (x_k^T a_i)^2$
- 3. Set $I_{k+1} = I_k \cup \{i_k\}$.
- 4. Compute x_{k+1} as the dominant eigenvector of $\sum_{j \in I_{k+1}} a_j a_j^T$.
- 5. Set k = k + 1. If k < n go back to step 2.

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• We can write the sparse PCA problem in penalized form:

$$\max_{\|x\| \le 1} x^T C x - \rho \operatorname{Card}(x)$$

in the variable $x \in \mathbf{R}^n$, where $\rho > 0$ is a parameter controlling sparsity.

• This problem is equivalent to solving:

$$\max_{\|x\|=1} \sum_{i=1}^{n} ((a_i^T x)^2 - \rho)_+$$

in the variable $x \in \mathbf{R}^n$, where the matrix A is the Cholesky decomposition of C, with $C = A^T A$. We only keep variables for which $(a_i^T x)^2 \ge \rho$.

• Sparse PCA equivalent to solving:

$$\max_{\|x\|=1} \sum_{i=1}^{n} ((a_i^T x)^2 - \rho)_+$$

in the variable $x \in \mathbf{R}^n$, where the matrix A is the Cholesky decomposition of C, with $C = A^T A$.

• This problem is also equivalent to solving:

$$\max_{X \succeq 0, \text{ Tr } X=1, \text{ Rank}(X)=1} \sum_{i=1}^{n} (a_i^T X a_i - \rho)_+$$

in the variables $X \in \mathbf{S}_n$, where $X = xx^T$. Note that the rank constraint can be dropped.

The problem

$$\max_{X \succeq 0, \text{ Tr } X=1} \sum_{i=1}^n (a_i^T X a_i - \rho)_+$$

is a convex maximization problem, hence is still hard. We can formulate a semidefinite relaxation by writing it in the equivalent form:

maximize
$$\sum_{i=1}^{n} \operatorname{Tr}(X^{1/2}a_{i}a_{i}^{T}X^{1/2} - \rho X)_{+}$$

subject to $\operatorname{Tr}(X) = 1, X \succeq 0, \operatorname{Rank}(X) = 1,$

in the variable $X \in \mathbf{S}_n$. If we drop the rank constraint, this becomes a convex problem and using

$$\mathbf{Tr}(X^{1/2}BX^{1/2})_{+} = \max_{\{0 \le P \le X\}} \mathbf{Tr}(PB) (= \min_{\{Y \ge B, Y \ge 0\}} \mathbf{Tr}(YX)).$$

we can get the following equivalent SDP:

max.
$$\sum_{i=1}^{n} \operatorname{Tr}(P_i B_i)$$

s.t. $\operatorname{Tr}(X) = 1, X \succeq 0, X \succeq P_i \succeq 0,$

which is a semidefinite program in the variables $X \in \mathbf{S}_n, P_i \in \mathbf{S}_n$.

Algorithms: optimality - Primal/dual formulation

• Primal problem:

max.
$$\sum_{i=1}^{n} \operatorname{Tr}(P_i B_i)$$

s.t. $\operatorname{Tr}(X) = 1, X \succeq 0, X \succeq P_i \succeq 0,$

which is a semidefinite program in the variables $X \in \mathbf{S}_n, P_i \in \mathbf{S}_n$.

• Dual problem:

min.
$$\lambda_{\max}(\sum_{i=1}^{n} Y_i)$$

s.t. $Y_i \succeq B_i, \ Y_i \succeq 0,$

• KKT conditions...

- When the solution of this last SDP has rank one, it also produces a globally optimal solution for the sparse PCA problem.
- In practice, this semidefinite program but we can use it to test the optimality of the solutions computed by the approximate greedy method.
- When the SDP has a rank one, the KKT optimality conditions for the semidefinite relaxation are given by:

$$\begin{cases} \left(\sum_{i=1}^{n} Y_{i}\right) X = \lambda_{\max}\left(\sum_{i=1}^{n} Y_{i}\right) X \\ x^{T}Y_{i}x = \begin{cases} \left(a_{i}^{T}x\right)^{2} - \rho \text{ if } i \in I \\ 0 \text{ if } i \in I^{c} \end{cases} \\ Y_{i} \succeq B_{i}, Y_{i} \succeq 0. \end{cases} \end{cases}$$

• This is a (large) semidefinite feasibility problem, but a **good guess** for Y_i often turns out to be sufficient.

Optimality: sufficient conditions. Given a sparsity pattern I, setting x to be the largest eigenvector of $\sum_{i \in I} a_i a_i^T$. If there is a parameter ρ_I such that:

$$\max_{i \notin I} (a_i^T x)^2 \le \rho_I \le \min_{i \in I} (a_i^T x)^2.$$

and

$$\lambda_{\max}\left(\sum_{i\in I}\frac{B_i x x^T B_i}{x^T B_i x} + \sum_{i\in I^c} Y_i\right) \le \sigma$$

where

$$Y_{i} = \max\left\{0, \rho \frac{(a_{i}^{T} a_{i} - \rho)}{(\rho - (a_{i}^{T} x)^{2})}\right\} \frac{(\mathbf{I} - xx^{T}) a_{i} a_{i}^{T} (\mathbf{I} - xx^{T})}{\|(\mathbf{I} - xx^{T}) a_{i}\|^{2}}, \quad i \in I^{c}.$$

Then the vector z such that $z = \operatorname{argmax}_{\{z_{I^c}=0, \|z\|=1\}} z^T \Sigma z$, which is formed by padding zeros to the dominant eigenvector of the submatrix $\Sigma_{I,I}$ is a global solution to the sparse PCA problem for $\rho = \rho_I$.

Optimality: why bother?

Compressed sensing. Following Candès & Tao (2005) (see also Donoho & Tanner (2005)), recover a signal $f \in \mathbf{R}^n$ from corrupted measurements:

$$y = Af + e,$$

where $A \in \mathbb{R}^{m \times n}$ is a coding matrix and $e \in \mathbb{R}^m$ is an unknown vector of errors with **low cardinality**.

This is equivalent to solving the following (combinatorial) problem:

minimize $||x||_0$ subject to Fx = Fy

where $||x||_0 = \mathbf{Card}(x)$ and $F \in \mathbf{R}^{p \times m}$ is a matrix such that FA = 0.

Compressed sensing: restricted isometry

Candès & Tao (2005): given a matrix $F \in \mathbb{R}^{p \times m}$ and an integer S such that $0 < S \leq m$, we define its **restricted isometry** constant δ_S as the smallest number such that for any subset $I \subset [1, m]$ of cardinality at most S we have:

$$(1 - \delta_S) \|c\|^2 \le \|F_I c\|^2 \le (1 + \delta_S) \|c\|^2,$$

for all $c \in \mathbf{R}^{|I|}$, where F_I is the submatrix of F formed by keeping only the columns of F in the set I.

Compressed sensing: perfect recovery

The following result then holds.

Proposition 1. Candès & Tao (2005). Suppose that the restricted isometry constants of a matrix $F \in \mathbb{R}^{p \times m}$ satisfy :

$$\delta_S + \delta_{2S} + \delta_{3S} < 1 \tag{1}$$

for some integer S such that $0 < S \le m$, then if x is an optimal solution of the convex program:

 $\begin{array}{ll} \text{minimize} & \|x\|_1\\ \text{subject to} & Fx = Fy \end{array}$

such that $Card(x) \le S$ then x is also an optimal solution of the combinatorial problem:

minimize $||x||_0$ subject to Fx = Fy.

Compressed sensing: restricted isometry

The restricted isometry constant δ_S in condition (1) can be computed by solving the following sparse PCA problem:

$$(1 + \delta_S) = \max x^T (F^T F) x$$

s. t. $\operatorname{Card}(x) \leq S$
 $\|x\| = 1,$

in the variable $x \in \mathbf{R}^m$ and another sparse PCA problem on $\alpha \mathbf{I} - F^T F$ to get the other inequality.

- Candès & Tao (2005) obtain an asymptotic proof that some random matrices satisfy the restricted isometry condition with overwhelming probability (i.e. exponentially small probability of failure)
- When they hold, the optimality conditions and upper bounds for sparse PCA allow us to prove (**deterministically** and with **polynomial complexity**) that a finite dimensional matrix satisfies the restricted isometry condition.

Optimality: Subset selection for least-squares

We consider p data points in \mathbb{R}^n , in a data matrix $X \in \mathbb{R}^{p \times n}$, and real numbers $y \in \mathbb{R}^p$. We consider the problem:

$$s(k) = \min_{w \in \mathbf{R}^n, \text{ Card } w \le k} \|y - Xw\|^2.$$
⁽²⁾

- Given the sparsity pattern $u \in \{0,1\}^n$, solution in closed form.
- **Proposition**: $u \in \{0,1\}^n$ is optimal for subset selection if and only if u is optimal for the sparse PCA problem on the matrix

$$X^{T}yy^{T}X - (y^{T}X(u)(X(u)^{T}X(u))^{-1}X(u)^{T}y)X^{T}X$$

- Sparse PCA allows to give deterministic sufficient conditions for optimality.
- To be compared on necessary and sufficient statistical consistency condition (Zhao & Yu (2006)):

 $||X_{I^c}^T X_I (X_I^T X_I)^{-1} \operatorname{sign}(w_I)||_{\infty} \leq 1$

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Numerical Results

Artificial data. We generate a matrix U of size 150 with uniformly distributed coefficients in [0, 1]. We let $v \in \mathbf{R}^{150}$ be a sparse vector with:

$$v_i = \begin{cases} 1 & \text{if } i \le 50\\ 1/(i-50) & \text{if } 50 < i \le 100\\ 0 & \text{otherwise} \end{cases}$$

We form a test matrix

$$\Sigma = U^T U + \sigma v v^T,$$

where σ is the signal-to-noise ratio.

Gene expression data. We run the approximate greedy algorithm on two gene expression data sets, one on **colon cancer** from Alon, Barkai, Notterman, Gish, Ybarra, Mack & Levine (1999), the other on **lymphoma** from Alizadeh, Eisen, Davis, Ma, Lossos & Rosenwald (2000). We only keep the 500 genes with largest variance.

Numerical Results - Artificial data



ROC curves for sorting, thresholding, fully greedy solutions and approximate greedy solutions for $\sigma = 2$.

Numerical Results - Artificial data



Variance versus cardinality tradeoff curves for $\sigma = 10$ (bottom), $\sigma = 50$ and $\sigma = 100$ (top). Optimal points are in bold.

Numerical Results - Gene expression data



Variance versus cardinality tradeoff curve for two gene expression data sets, lymphoma (top) and colon cancer (bottom). Optimal points are in bold.

Numerical Results - Subset selection on a noisy sparse vector



Backward greedy algorithm and Lasso. Probability of achieved (red dotted line) and provable (black solid line) optimality versus noise for greedy selection against Lasso (green large dots). *Left:* Lasso consistency condition satisfied (Zhao & Yu (2006)). *Right:* consistency condition not satisfied.

Conclusion & Extensions

Sparse PCA in **practice**, if your problem has. . .

- A **million** variables: can't even form a covariance matrix. **Sort** variables according to variance and keep a few thousand.
- A few **thousand** variables (more if Gram format): **approximate greedy** method described here.
- A few hundred variables: use DSPCA, SPCA, full greedy search, etc.

Of course, these techniques can be combined.

Discussion - Extensions

- Large SDP to obtain certificated of optimality of a combinatorial problem
- Efficient solvers for the semidefinite relaxation (exploiting low rank, randomization, etc.). (We have never solved it for n > 10!)
- Find better matrices with restricted isometry property.

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