Efficient Advert Assignment

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Ad Efficient Advert Talk @google.com Neil Walton, University of Amsterdam.

Ads (1) Joint work with

Search

Frank Kelly university of cambridge

Peter Key microsoft research

This Talk

We give a simple allocation and pricing mechanism

whose Nash equilibrium solves a very large optimization problem

This Talk

We give a simple allocation and pricing mechanism

whose Nash equilibrium solves a very large optimization problem

Very Large = over the infinite results of a search engine.

Outline

- A Introduction to Sponsored Search
 - Bids, Impressions, Click-Through Rate, Advertizers, Platform
- Auction or Optimize?
 - Our Mechanism, Generalized 2nd Price, VCG Mechanism, Decomposition.
- Our Results
 - Main Theorem, Implementations
- Further Results and Extensions
 - Dynamics, Multivariate Utilities, General Page Layouts, Budgets.

Introduction to Sponsored Search



Zurich Hotels

Q



Zurich Hotels

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•	Ad Group #1	Eligible	€0.15	0	0	0.00%	€0.00	€0.00	0.0

	Ad group	Status ?	Default Max. CPC	Clicks 🔝 🕌	Impr. 📍	CTR 📱	Avg. CPC	Cost ?	Avg. Pos. 👔
•	Ad Group #1	Eligible	€0.15	0	0	0.00%	€0.00	€0.00	0.0
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A man so truly in love with golf risks the dangers of a crocodile to get his ball back.

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Variability in Sponsored Search

A **mixture** of auctions:

(phrase match)







The Searcher



The Searcher

The Platform



Google bing YAHOO!

The Searcher

The Platform

The Advertiser







The Searcher

The Platform

The Advertiser



Google bing YAHOO!



Search Distribution



The Searcher

The Platform

The Advertiser



Search Distribution

 $\mathbb{P}_{ au}$



When a search occurs Click-Through p_{il}^{τ} Assignment x_{il}^{τ}



The Searcher

The Platform

Search Distribution

 $\mathbb{P}_{ au}$



When a search occurs Click-Through p_{il}^{τ} Assignment x_{il}^{τ}

The Advertiser



Receives average information

Click-Through $\mathbb{E}_{ au} \sum_{l} p_{il}^{ au} x_{il}^{ au}$ Assignment $\mathbb{E}_{ au} x_{il}^{ au}$

The Searcher

The Platform

Google

YAHOO!

Search Distribution



When a search occurs Click-Through p_{il}^{τ} Assignment x_{il}^{τ}

The Advertiser



Receives average information

Click-Through $\mathbb{E}_{ au} \sum_{l} p_{il}^{ au} x_{il}^{ au}$ Assignment $\mathbb{E}_{ au} x_{il}^{ au}$

The Searcher

The Platform

The Advertiser



Search Distribution



Google bing YAHOO!





Receives average information

Click-Through $\mathbb{E}_{ au} \sum_{l} p_{il}^{ au} x_{il}^{ au}$ Assignment $\mathbb{E}_{ au} x_{il}^{ au}$

The Searcher

The Platform

The Advertiser



Search Distribution



Google bing YAHOO!





Receives average information Click-Through $\mathbb{E}_{ au} \sum_{l} p_{il}^{ au} x_{il}^{ au}$

Assignment

ment $\mathbb{E}_{\tau} x_{il}^{\tau}$ Platform knows Advertiser knows

Auction or Optimize?

Two Auctions

- au Search Type
- λ_i Bid of ad i

$p_{il}^{\tau}~$ – Click-Through ad i slot l $y_i^{\tau}(\lambda)$ – CTR given bids

Two Auctions



 $p_{il}^{\tau}~$ – Click-Through ad i slot l $y_i^{\tau}(\lambda)$ – CTR given bids

Two Auctions



 $p_{il}^{\tau}~$ – Click-Through ad i slot l $y_i^{\tau}(\lambda)$ – CTR given bids


 $p_{il}^{\tau}~$ – Click-Through ad i slot l $y_i^{\tau}(\lambda)$ – CTR given bids





 $p_{il}^{\tau}~$ – Click-Through ad i slot l $y_i^{\tau}(\lambda)$ – CTR given bids

Auction 2:

Assign max matching

 $\max \sum_{i} \sum_{l} \lambda_{i} p_{il}^{\tau} x_{il}$ s.t. $\sum_{i} x_{il}^{\tau} \leq 1, \quad \sum_{l} x_{il}^{\tau} \leq 1$



p_{il}^{τ} – Click-Through ad i slot l $y_i^{\tau}(\lambda)$ – CTR given bids Auction 2: Assign max matching $\max \sum_{i} \sum_{l} \lambda_{i} p_{il}^{\tau} x_{il}$ s.t. $\sum_{i} x_{il}^{\tau} \le 1, \quad \sum_{i} x_{il}^{\tau} \le 1$ **Pay**, per-click $\mu_i \sim U[0, \lambda_i]$ $\lambda_i \left(1 - \frac{y_i^{\tau}(\mu_i)}{u_i^{\tau}(\lambda_i)} \right)$



 p_{il}^{τ} – Click-Through ad i slot l $y_i^{\tau}(\lambda)$ – CTR given bids Auction 2: Assign max matching $\max \sum_{i} \sum_{l} \lambda_{i} p_{il}^{\tau} x_{il}$ s.t. $\sum_{i} x_{il}^{\tau} \le 1$, $\sum_{i} x_{il}^{\tau} \le 1$ **Pay**, per-click $\mu_i \sim U[0, \lambda_i]$ $\lambda_i \left(1 - \frac{y_i^{\tau}(\mu_i)}{u_i^{\tau}(\lambda_i)} \right)$



 p_{il}^{τ} - Click-Through ad i slot l $y_i^{\tau}(\lambda)$ - CTR given bids Auction 2: Assign max matching $\max \sum_{i} \sum_{i} \lambda_i p_{il}^{\tau} x_{il}$

s.t.
$$\sum_{i} x_{il}^{\tau} \le 1$$
, $\sum_{l} x_{il}^{\tau} \le 1$

Pay, per-click $\mu_i \sim \mathrm{U}[0,\lambda_i]$

$$\lambda_i \left(1 - \frac{y_i^\tau(\mu_i)}{y_i^\tau(\lambda_i)} \right)$$

A VCG Auction

Immediate Advantages

Immediate Advantages



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Immediate Advantages



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Google (2006) said:

"Google's unique auction model uses Nobel Prize winning economic theory ... the AdWords™ Discounter makes sure that they only pay what they need in order to stay ahead of their nearest competitor."

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Q. Is this really true?

Google (2006) said:

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Q. Is this really true? A. Not really.

Google (2006) said:

"Google's unique auction model uses Nobel Prize winning economic theory ... the AdWords™ Discounter makes sure that they only pay what they need in order to stay ahead of their nearest competitor."

Q. Is this really true?Q. What did they really mean?A. Not really.

Google (2006) said:

"Google's unique auction model uses Nobel Prize winning economic theory ... the AdWords™ Discounter makes sure that they only pay what they need in order to stay ahead of their nearest competitor."

Q. Is this really true? A. Not really. Q. What did they really mean? A. The VCG Mechanism...





Clark





• Advertiser's utilities $U_i(\cdot)$

- Advertiser's utilities
- bid utilities

 $U_i(\cdot) \\ V_i(\cdot)$

- Advertiser's utilities
- bid utilities
- Assignment constraints

$$U_i(\cdot) \ V_i(\cdot) \ \mathcal{A}$$

- Advertiser's utilities
- bid utilities
- Assignment constraints

Platform Assigns:

$$U_i(\cdot)$$
$$V_i(\cdot)$$
$$\mathcal{A}$$

$$y_i^* \in \operatorname{argmax}_{y \in \mathcal{A}} \sum_i V_i(y_i)$$

- Advertiser's utilities
- bid utilities
- Assignment constraints

Platform Assigns:



- Advertiser's utilities
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Platform Assigns:

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$$U_i(\cdot)$$

 $V_i(\cdot)$
 \mathcal{A}

$$y_i^* \in \operatorname{argmax}_{y \in \mathcal{A}} \sum_i V_i(y_i)$$

$$\pi_i = \max_{y \in \mathcal{A}} \sum_{j \neq i} V_j(y_j) - \sum_{j \neq i} V_j(y_j^*)$$

- Advertiser's utilities
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 - Platform Assigns:

$$U_i(\cdot)$$

 $V_i(\cdot)$

$$y_i^* \in \operatorname{argmax}_{y \in \mathcal{A}} \sum_i V_i(y_i)$$

$$\pi_{i} = \max_{y \in \mathcal{A}} \sum_{j \neq i} V_{j}(y_{j}) - \sum_{j \neq i} V_{j}(y_{j}^{*})$$

Everyone else's
value

- Advertiser's utilities
- bid utilities
- Assignment constraints
 - Platform Assigns:

$$U_i(\cdot)$$

 $V_i(\cdot)$
 A

$$y_i^* \in \operatorname{argmax}_{y \in \mathcal{A}} \sum_i V_i(y_i)$$

$$\pi_{i} = \max_{y \in \mathcal{A}} \sum_{j \neq i} V_{j}(y_{j}) - \sum_{j \neq i} V_{j}(y_{j}^{*})$$

Everyone else's value Value without you there

- Advertiser's utilities
- bid utilities
- Assignment constraints
 - Platform Assigns:

$$U_i(\cdot)$$

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 \mathcal{A}

$$y_i^* \in \operatorname{argmax}_{y \in \mathcal{A}} \sum_i V_i(y_i)$$

$$\pi_i = \max_{y \in \mathcal{A}} \sum_{j \neq i} V_j(y_j) - \sum_{j \neq i} V_j(y_j^*)$$

- Advertiser's utilities
- bid utilities
- Assignment constraints

Platform Assigns:

$$y_i^* \in \operatorname{argmax}_{y \in \mathcal{A}} \sum_i V_i(y_i)$$

Platform Prices:

$$\pi_i = \max_{y \in \mathcal{A}} \sum_{j \neq i} V_j(y_j) - \sum_{j \neq i} V_j(y_j^*)$$

Equilibrium Advertizer:

$$\max_{V_i} \left\{ U_i(y_i^*) - \pi_i \right\}$$

$$U_i(\cdot)$$
$$V_i(\cdot)$$
$$\mathcal{A}$$

Theorem

The VCG mechanism has a dominate strategies equilibrium that is:

Theorem

The VCG mechanism has a dominate strategies equilibrium that is:

- Incentive compatible

bids are truthful: $V_i(\cdot) = U_i(\cdot)$

Theorem

The VCG mechanism has a dominate strategies equilibrium that is:

- Incentive compatible

bids are truthful: $V_i(\cdot) = U_i(\cdot)$

– Efficient

allocation is optimal: $y^* \in \operatorname{argmax}_{y \in \mathcal{A}} \sum_i U_i(y_i)$

Ргоз

- **1.** Result applies in very **general** settings
- 2. Allocation of Adverts is provably **optimal**

Ρгος

- **1.** Result applies in very **general** settings
- 2. Allocation of Adverts is provably **optimal**

Cons

1. Advertisers submit their **entire utility** function $V_i(\cdot) = U_i(\cdot)$

Ргоз

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- **1.** Advertisers submit their **entire utility** function $V_i(\cdot) = U_i(\cdot)$
- **2.** Utility $U_i(y_i)$ isn't for a single adauction but **for all adauctions**

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- 3. Platform needs to solve a massive optimization

$$y^* \in \operatorname{argmax}_{y \in \mathcal{A}} \sum_i U_i(y_i)$$

Ргоз

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l

This talk: We deal with these issue by appropriately decomposing this massive optimization.
Maximize $\sum_{i} U_i(y_i)$ subject to $y_i = \mathbb{E}_{\tau} \left[\sum_{i} p_{il}^{\tau} x_{il}^{\tau} \right], \quad i,$ $\sum_{i} x_{il}^{\tau} \le 1, \qquad l, \tau,$ $\sum_{l} x_{il}^{\tau} \leq 1, \qquad i, \tau,$ over $x_{il}^{\tau} \geq 0, y_i \geq 0 \qquad l, \qquad i, \quad \tau.$

Maximize $\sum_{i} U_{i}(y_{i})$
subject to $y_{i} = \mathbb{E}_{\tau} \left[\sum_{l} p_{il}^{\tau} x_{il}^{\tau} \right], \quad i,$
 $\sum_{i} x_{il}^{\tau} \leq 1, \qquad l, \tau,$

$$\sum_{l} x_{il}^{\tau} \le 1, \qquad i, \tau,$$

 $x_{il}^{\tau} \ge 0, y_i \ge 0 \qquad l, \qquad i, \qquad \tau.$

over













$$\begin{split} \text{Maximize} & \sum_{i} U_{i}(y_{i}) \\ \text{subject to} & y_{i} = \mathbb{E}_{\tau} \Big[\sum_{l} p_{il}^{\tau} x_{il}^{\tau} \Big], \quad i, \\ & \sum_{i} x_{il}^{\tau} \leq 1, \qquad l, \tau, \\ & \sum_{l} x_{il}^{\tau} \leq 1, \qquad i, \tau, \\ \text{over} & x_{il}^{\tau} \geq 0, y_{i} \geq 0 \qquad l, \qquad i, \quad \tau. \end{split}$$

• even if we knew all the parameters, it's impossible to solve this optimization off-line

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- even if we knew all the parameters, it's impossible to solve this optimization off-line
- Still ... maybe we can solve a lot of small optimizations...

When a search $\, au\,$ occurs, solve:

 $\begin{array}{ll} \text{Maximize} & \sum_{i} \sum_{l} \lambda_{i} p_{il}^{\tau} x_{il} \\ \text{subject to} & \sum_{i} x_{il}^{\tau} \leq 1, \\ & \sum_{l} x_{il}^{\tau} \leq 1, \\ \text{over} & x_{il}^{\tau} \geq 0, y_{i} \geq 0. \end{array}$







Lots of polynomial time algorithms:



Lots of polynomial time algorithms: Hungarian ; Hopcroft-Karp ; Bertsekas' Auction ...

Solve the big optimization with many little optimizations

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1. Substitution:

 $\begin{aligned} x^* &\in \arg\max_{x\in\mathcal{X}} U(x) \\ \Longleftrightarrow \ x^* &\in \arg\max_{x\in\mathcal{X}} V(x;\lambda^*), \\ \lambda^* &\in \arg\max_{\lambda\in\Lambda} A(\lambda;x^*) \end{aligned}$

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2. Separability:

$$\max_{\substack{x \in \mathcal{X}, y \in \mathcal{Y} \\ x \in \mathcal{X}}} \left\{ f(x) + g(y) \right\}$$
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Substitute utility $U_i(y_i)$ for $\lambda_i y_i$ MAIN IDEA: **Separate** out the resulting optimization

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MAIN IDEA:Substitute utility $U_i(y_i)$ for $\lambda_i y_i$ Separate out the resulting optimizationTHE RESULT:A massively distributed VCG Mechanism

Our Results

$$\max_{x,y} L(x,y;\lambda) = \max_{x,y} \left[\sum_{i} U_i(y_i) + \sum_{i} \lambda_i \left(y_i - \mathbb{E}_{\tau} \left[\sum_{l} p_{il}^{\tau} x_{il}^{\tau} \right] \right) \right]$$

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$$\textbf{LF-Transform}$$
$$\textbf{Assignment}$$
$$\textbf{Problem}$$

PROPOSITION 2 (Decomposition). Variables \tilde{y} , $\tilde{x}^{\tau}, \tau \in \mathcal{T}$, satisfying the feasibility conditions (7b-7e) are optimal for $SYSTEM(U, \mathcal{I}, \mathbb{P}_{\tau})$ if and only if there exist $\tilde{\lambda}_i$, $i \in \mathcal{I}$, such that A. $\tilde{\lambda}_i$ minimizes $U_i^*(\lambda_i) + \lambda_i \tilde{y}_i$ over $\lambda_i \geq 0$, for each $i \in \mathcal{I}$,

B. \tilde{x}^{τ} solves ASSIGNMENT $(\tau, \tilde{\lambda})$, with probability one under the distribution \mathbb{P}_{τ} over $\tau \in \mathcal{T}$.

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Advertiser's must signal average prices

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Platform solves Assignment when each search occurs

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- Decomposed on the timescales of Platform and Advertisers.
- Search distribution is not required.
- But it's an optimization result, we must incentivize this behaviour.

Advertizers maximizes rewards:

$$r_i(\lambda) = U_i(y_i(\lambda)) - \pi_i(\lambda)y_i(\lambda).$$

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THEOREM 1. If prices are charged according to the price function

$$\pi_i(\lambda) = \frac{1}{y_i(\lambda)} \int_0^{\lambda_i} \left(y_i(\lambda) - y_i(\mu_i, \lambda_{-i}) \right) d\mu_i$$
(18)

then there exists a unique Nash equilibrium, and it is given by the vector of optimal prices identified in Proposition 3. Thus the assignment achieved at the Nash equilibrium, $(x(\lambda^*), y(\lambda^*))$, is a solution to the SYSTEM optimization.
Main Theorem and Mechanism Design

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solve the Massive Optimization

Proof of Main Theorem

Optimality condition for the dual:

$$\begin{split} \min_{\lambda \ge 0} \quad \sum_{i} \left[U_{i}^{*}(\lambda_{i}) + \lambda_{i} y_{i}(\lambda) \right] & \underset{\text{envelope}}{\longleftrightarrow} \quad \frac{d}{d\lambda_{i}} U_{i}^{*}(\lambda_{i}) + y_{i}(\lambda) = 0, \quad \forall i \\ & \underset{\text{integrate}}{\longleftrightarrow} \quad \min_{\lambda_{i} \ge 0} \quad U_{i}^{*}(\lambda_{i}) + \int_{0}^{\lambda_{i}} y_{i}(\mu_{i}, \lambda_{-i}) d\mu_{i}, \quad \forall i \\ & \underset{\text{Moreau thrm}}{\longleftrightarrow} \quad \max_{y_{i} \ge 0} \quad U_{i}(y_{i}) - \int_{0}^{\infty} [y_{i} - y_{i}(\mu_{i}, \lambda_{-i})]_{+} d\mu_{i}, \quad \forall i \\ & \underset{\text{Substitute}}{\longleftrightarrow} \quad \max_{\lambda_{i} \ge 0} \quad U_{i}(y_{i}(\lambda)) - \int_{0}^{\lambda_{i}} y_{i}(\lambda) - y_{i}(\mu_{i}, \lambda_{-i}) d\mu_{i}, \quad \forall i \end{split}$$

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1. Let $\mu_i \sim \text{Uniform}[0, \lambda_i]$ 2. A discounted-VCG price and price

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A massively distributed VCG mechanism

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A very simple pay-per click mechanism:

A massively distributed VCG mechanism A very simple pay-per click mechanism: Assignment Pricing

$$\max \sum_{i} \sum_{l} \lambda_{i} p_{il}^{\tau} x_{il}^{\tau}$$

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A massively distributed VCG mechanismA very simple pay-per click mechanism:AssignmentPricing $\max \sum_{i} \sum_{l} \lambda_i p_{il}^{\tau} x_{il}^{\tau}$ $\lambda_i \left(1 - \frac{y_i^{\tau}(\mu_i)}{y_i^{\tau}(\lambda_i)}\right)$

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Further Results and Extensions

A natural dynamic: $\frac{d}{dt}\lambda_i(t) \ge 0$ according as $\lambda_i(t) \le U'_i(y_i(\lambda(t)))$.

A natural dynamic: $\frac{d}{dt}\lambda_i(t) \ge 0$ according as $\lambda_i(t) \le U'_i(y_i(\lambda(t)))$. Lyapunov function: $\mathcal{V}(\lambda) = \sum_i \left[U_i^*(\lambda_i) + \lambda_i y_i(\lambda)\right]$.





Further Extensions

Controlling number of slots:



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Further Extensions

Multivariate utilities:

 $U_i(y_{ik}:k\in\mathcal{K}_i)$

+ KE	WORDS	Edit 🔻 Details 👻 Bid s	trategy - Automate -	Labels 🔻							
	•	Keyword	Status ?	Max. CPC 📳 🔸	Clicks ?	Impr. 🕄	CTR 🔋	Avg. CPC	Cost ?	Avg. Pos. 👔	Labels ?
	н	ifor eth	Paused	€5.25 <i>⊠</i>	0	0	0.00%	€0.00	€0.00	0.0	
	н	eth ifor	Paused	€5.00 <i>⊠</i>	0	0	0.00%	€0.00	€0.00	0.0	
	•	Advert Zurich	⊊ Eligible	€5.00 <i>⊠</i>	0	20	0.00%	€0.00	€0.00	1.2	-
	•	optimization zurich	⊊ Eligible	€5.00 <i>⊵</i>	0	1	0.00%	€0.00	€0.00	1.0	-
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	•	optimization eth	⊊ Eligible	€2.25 ⊘	0	0	0.00%	€0.00	€0.00	0.0	
	н	eth zentrum	Paused	€0.10	0	3	0.00%	€0.00	€0.00	1.0	

Further Extensions

Budget constraints:
$$U_i(y_{ik} : k \in \mathcal{K}_i) = \frac{b_i}{q} \log \sum_{k \in \mathcal{K}_i} (w_{ik} y_{ik})^q$$

Bid strategy	?	Basic opti	ons Advanced options			
		O I'll man	ually set my bids for clicks			
		AdWord	Is will set my bids to help maximize clicks within my target budget			
Budget	?	CA\$	per day			
		Actual daily spend may vary. ?				

- Massively decomposed VCG implementation
 - Simple
 - Flexible

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- Provably solves an Infinitely Large Optimization.

Thank you for listening!

NEIL WALTON'S WEBSITE RESEARCH NEIL'S NOTES LINKS

Efficient Advert Assignment F.P. Kelly, P. Key, N.S. Walton. (2014). (Preprint) (an earlier version was presented at EC'14 see below) [pdf] [arxiv]

Incentivized Optimal Advert Assignment via Utility Decomposition F.P. Kelly, P. Key, N.S. Walton (2014). *ACM Conference on Economics and Computation*. [pdf] [arxiv] [proceedings] [bibtex]