A Market Test for the Positivity of Arrow-Debreu Prices

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Introduction

- Classic Black & Scholes (1973) option pricing based on:
 - a *dynamic hedging* argument
 - a *model* for the asset dynamics (geometric BM)
- Sensitive to liquidity, transaction costs, model risk ...
- What can we say about derivative prices with much weaker assumptions?

Static Arbitrage

Here, we rely on a *minimal set of assumptions*:

- no assumption on the asset distribution
- one period model

An arbitrage in this simple setting is a *buy and hold* strategy:

- form a portfolio at no cost today with a strictly positive payoff at maturity
- no trading involved between today and the option's maturity

What for?

- Data validation (e.g. before calibration), static arbitrage means market data is incompatible with *any* dynamic model. . .
- Test extrapolation formulas
- In illiquid markets, find optimal static hedge

Outline

• Static Arbitrage

- Harmonic Analysis on Semigroups
- No Arbitrage Conditions

Simplest Example: Put Call Parity



Static Arbitrage: Calls

Also, necessary and sufficient conditions on call prices:

Suppose we have a set of market prices for calls $C(K_i) = p_i$, then there is no arbitrage iff there is a function C(K):

- C(K) positive
- C(K) decreasing
- C(K) convex
- $C(K_i) = p_i \text{ and } C(0) = S$

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This is very easy to test. . .
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Dow Jones index call option prices on Mar. 17 2004, maturity Apr. 16 2004





A. d'Aspremont, Third Oxford-Princeton Conference on Financial Mathematics, Nov. 12 2005.

Why?

Data quality...

- All the prices are last quotes (not simultaneous)
- Low volume
- Some transaction costs

Problem: this data is used to calibrate models and price other derivatives...

Dimension n: Basket Options

• A basket call payoff is given by:

$$\left(\sum_{i=1}^{k} w_i S_i - K\right)^+$$

where w_1, \ldots, w_k are the basket's weights and K is the option's strike price

- Examples include: Index options, spread options, swaptions...
- Basket option prices are used to gather information on *correlation*

We denote by C(w, K) the price of such an option, can we get conditions to test basket price data?

Necessary Conditions

Similar to dimension one...

Suppose we have a set of market prices for calls $C(w_i, K_i) = p_i$, and there is no arbitrage, then the function C(w, K) satisfies:

- C(w, K) positive
- C(w, K) decreasing in K, increasing in w
- C(w, K) jointly convex in (w, K)
- $C(w_i, K_i) = p_i$ and C(0) = S

This is still *tractable* in dimension n as a *linear program*.

Sufficient?

A key difference with dimension one: Bertsimas & Popescu (2002) show that the exact problem is NP-Hard.

- These conditions are *only necessary*...
- Numerical cost is minimal (small LP)
- We can show *sufficiency* in some particular cases (see d'Aspremont & El Ghaoui (2005) and Davis & Hobson (2005) for details)

In practice: these conditions are far from being tight, how can we *refine* them?

Arrow-Debreu prices

• Arrow-Debreu: There is no arbitrage in the static market iff there is a probability measure π such that:

$$C(w,K) = \mathbf{E}_{\pi}(w^T x - K)^+$$

- $\pi(x)$ represents Arrow-Debreu state prices.
- Discretize on a uniform grid: This turns this into a *linear program* with m^n variables, where n is the number of assets x_i and m is the number of bins.
- Numerically: hopeless. . .
- Explicit conditions derived by Henkin & Shananin (1990) (link with Radon transform), but intractable. . .

Tractable Conditions

• Bochner's theorem on the Fourier transform of positive measures:

$$f(s) = \int e^{-i < s, x >} d\lambda(x)$$
 with λ positive

 $f(\boldsymbol{s})$ positive semidefinite

 \uparrow

which means testing if the matrices $f(s_i s_j)$ are positive semidefinite

• Can we generalize this result to other transforms? In particular:

$$\int_{\mathbf{R}^n_+} (w^T x - K)^+ d\pi(x)$$

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Harmonic Analysis on Semigroups

Some quick definitions...

• A pair (\mathbb{S}, \cdot) is called a *semigroup* iff:

o if s, t ∈ S then s · t is also in S
o there is a neutral element e ∈ S such that e · s = s for all s ∈ S

- The *dual* S^* of S is the set of *semicharacters*, *i.e.* applications $\chi : S \to \mathbf{R}$ such that
 - $\chi(s)\chi(t) = \chi(s \cdot t)$ for all s, t ∈ S $\chi(e) = 1$, where e is the neutral element in S
- A function f : S → R is positive semidefinite iff for every family {s_i} ⊂ S the matrix with elements f(s_i · s_j) is positive semidefinite

Harmonic Analysis on Semigroups

Last definitions (honest)...

• A function α is called an *absolute value* on $\mathbb S$ iff

$$\circ \ \alpha(e) = 1 \circ \ \alpha(s \cdot t) \le \alpha(s)\alpha(t), \text{ for all } s, t \in \mathbb{S}$$

• A function f is *bounded* with respect to the absolute value α iff there is a constant C > 0 such that

$$|f(s)| \le C\alpha(s), \quad s \in \mathbb{S}$$

• *f* is *exponentially bounded* iff it is bounded with respect to an absolute value

Carleman type conditions on growth for moment determinacy, etc. . .

Harmonic Analysis on Semigroups: Central Result

The central result, see Berg, Christensen & Ressel (1984) based on Choquet's theorem:

- the set of exponentially bounded *positive definite functions* is a *Bauer simplex* whose extreme points are the bounded semicharacters...
- this means that we have the following representation for positive definite functions on \mathbb{S} :

$$f(s) = \int_{\mathbb{S}^*} \chi(s) d\mu(\chi)$$

where μ is a Radon measure on \mathbb{S}^*

Harmonic Analysis on Semigroups: Simple Examples

• *Berstein's theorem* for the Laplace transform

$$\mathbb{S}=(\mathbf{R}_+,+)\text{, }\chi_x(t)=e^{-xt} \ \text{ and } \ f(t)=\int_{\mathbf{R}_+}e^{-xt}d\mu(x)$$

• with involution, *Bochner's theorem* for the Fourier transform

$$\mathbb{S} = (\mathbf{R}, +), \ \chi_x(t) = e^{2\pi i x t} \text{ and } f(t) = \int_{\mathbf{R}} e^{2\pi i x t} d\mu(x)$$

• *Hamburger's solution* to the unidimensional moment problem

$$\mathbb{S} = (\mathbf{N}, +), \ \chi_x(k) = x^k \text{ and } f(k) = \int_{\mathbf{R}} x^k d\mu(x)$$

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The Option Pricing Problem Revisited

What is the appropriate semigroup here?

- Basket option payoffs $(w^T x K)^+$ are not ideal in this setting.
- Solution: use *straddles*: $|w^T x K|$
- Straddles are just the *sum of a call and a put*, their price can be computed from that of the corresponding call and forward by call-put parity.
- The fact that $|w^T x K|^2$ is a polynomial keeps the complexity low.

Payoff Semigroup

 The fundamental semigroup S here is the multiplicative payoff semigroup generated by the cash, the forwards and the straddles:

$$\mathbb{S} = \{1, x_1, \dots, x_n, |w_1^T x - K_1|, \dots, |w_m^T x - K_m|, x_1^2, x_1 x_2, \dots\}$$

• The *semicharacters* are the functions $\chi_x : \mathbb{S} \to \mathbb{R}$ which evaluate the payoffs at a certain point x

$$\chi_x(s) = s(x), \text{ for all } s \in \mathbb{S}$$

The Option Pricing Problem Revisited

• The original static arbitrage problem can be reformulated as

find
$$f$$

subject to $f(|w_i^T x - K_i|) = p_i, \quad i = 1, ..., m$
 $f(s) = \mathbf{E}_{\pi}[s], \quad s \in \mathbb{S}$ (f moment function)

- The variable is now $f: \mathbb{S} \to \mathbf{R}$, a function that associates to each payoff s in \mathbb{S} , its price f(s)
- The *representation result* in Berg et al. (1984) shows when a (price) function f : S → R can be represented as

$$f(s) = \mathbf{E}_{\pi}[s]$$

Option Pricing: Main Theorem

If we assume that the asset distribution has a compact support included in \mathbf{R}_{+}^{n} , and note e_{i} for i = 1, ..., n + m the forward and option payoff functions we get:

A function $f(s): \mathbb{S} \to \mathbf{R}$ can be represented as

$$f(s) = \mathbf{E}_{\nu}[s(x)], \text{ for all } s \in \mathbb{S},$$

for some measure ν with compact support, iff for some $\beta > 0$:

(i) f(s) is positive semidefinite

(ii) $f(e_i s)$ is positive semidefinite for i = 1, ..., n + m

(iii) $\left(\beta f(s) - \sum_{i=1}^{n+m} f(e_i s)\right)$ is positive semidefinite

this turns the basket arbitrage problem into a *semidefinite program*

Semidefinite Programming

A *semidefinite program* is written:

minimize
$$\operatorname{Tr} CX$$

subject to $\operatorname{Tr} A_i X = b_i, \quad i = 1, \dots, m$
 $X \succeq 0,$

in the variable $X \in \mathbf{S}^n$, with parameters $C, A_i \in \mathbf{S}^n$ and $b_i \in \mathbf{R}$ for i = 1, ..., m. Its *dual* is given by:

maximize
$$b^T \lambda$$

subject to $C - \sum_{i=1}^m \lambda_i A_i \succeq 0$,

in the variable $\lambda \in \mathbf{R}^m$.

Extension of interior point techniques for linear programming show how to solve these convex programs *efficiently* (see Nesterov & Nemirovskii (1994), Sturm (1999) and Boyd & Vandenberghe (2004)).

Option Pricing: a Semidefinite Program

We get a relaxation by only sampling the elements of S up to a certain degree, the variable is then the vector f(s) with

$$e = (1, x_1, \dots, x_n, |w_1^T x - K_1|, \dots, |w_m^T x - K_m|, x_1^2, x_1 x_2, \dots, |w_m^T x - K_m|^N)$$

testing for the absence of arbitrage is then a *semidefinite program*:

find
$$f$$

subject to $M_N(f(s)) \succeq 0$
 $M_N(f(e_j s)) \succeq 0$, for $j = 1, ..., n$,
 $M_N\left(f((\beta - \sum_{k=1}^{n+m} e_k)s)\right) \succeq 0$
 $f(e_j) = p_j$, for $j = 1, ..., n+m$ and $s \in \mathbb{S}$

where $M_N(f(s))_{ij} = f(s_i s_j)$ and $M_N(f(e_k s))_{ij} = f(e_k s_i s_j)$

Conic Duality

Let $\Sigma \subset \mathcal{A}(\mathbb{S})$ be the set of polynomials that are sums of squares of polynomials in $\mathcal{A}(\mathbb{S})$, and \mathcal{P} the set of positive semidefinite sequences on \mathbb{S}

instead of the conic duality between probability measures and positive portfolios

$$p(x) \ge 0 \Leftrightarrow \int p(x) d\nu \ge 0$$
, for all measures ν

• we use the duality between positive semidefinite sequences $\mathcal P$ and sums of squares polynomials Σ

$$p \in \Sigma \Leftrightarrow \langle f, p \rangle \ge 0 \text{ for all } f \in \mathcal{P}$$

with $p = \sum_{i} q_i \chi_{s_i}$ and $f : \mathbb{S} \to \mathbf{R}$, where $\langle f, p \rangle = \sum_{i} q_i f(s_i)$

Option Pricing: Caveats

- *Size*: grows exponentially with the number of assets: no free lunch...
- In dimension 2, for spread options, this is:

$$\binom{2+d}{2}(k+1)$$

where d is the degree of the relaxation and k the number of assets.

• Conditioning issues. . .

Conclusion

- Testing for static arbitrage in option price data is easy in dimension one
- The extension on basket options (swaptions, etc) is NP-hard but good relaxations can be found
- We get a computationally friendly set of conditions for the absence of arbitrage
- Small scale problems are tractable in practice as semidefinite programs

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