A direct formulation for sparse PCA using semidefinite programming

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Introduction

Principal Component Analysis (*PCA*): classic tool in multivariate data analysis

- **Input**: a covariance matrix A
- **Output**: a sequence of *factors* ranked by *variance*
- Each factor is a *linear* combination of the problem variables

Typical use: *reduce* the number of *dimensions* of a model while *maximizing* the *information* (variance) contained in the simplified model.

Introduction

Numerically: just an eigenvalue decomposition of the covariance matrix:

$$A = \sum_{i=1}^{n} \lambda_i x_i x_i^T$$

where...

- The factors x_i are uncorrelated
- The result of the PCA is usually not sparse, i.e. each factor is a linear combination of *all the variables* in the model.

Can we get *sparse* factors instead?

Sparse PCA: Applications

Why sparse factors?

- *Financial time series analysis*: sparse factors often mean less assets in the portfolio, hence less fixed transaction costs
- *Multiscale data processing*: get sparse structure from motion data, ...
- *Gene expression data*: each variable is a particular gene, sparse factors highlight the action of a few genes, making interpretation easier
- *Image processing*: sparse factors involve only specific zones or objects in the image.

Applications, previous works

Further details. . .

- Financial time series analysis, dimensionality reduction, hedging, etc (Rebonato (1998),...)
- Multiscale data processing (Chennubhotla & Jepson (2001),...)
- Gene expression data (survey by Wall, Rechtsteiner & Rocha (2002), ...)
- Signal & image processing, vision, OCR, ECG (Johnstone & Lu (2003))

A: rank one approximation

Problem definition:

• Here, we focus on the *first factor* x, computed as the solution of:

$$\min_{x \in \mathbf{R}} \|A - xx^T\|_F$$

where $||X||_F$ is the Frobenius norm of X, i.e. $||X||_F = \sqrt{\mathbf{Tr}(X^2)}$

• In this case, we get an *exact* solution $\lambda^{\max}(A)x_1x_1^T$ where $\lambda^{\max}(X)$ is the maximum eigenvalue and x_1 is the associated eigenvector.

Variational formulation

We can rewrite the previous problem as:

$$\begin{array}{ll} \max & x^T A x \\ \text{subject to} & \|x\|_2 = 1. \end{array} \tag{1}$$

This problem is *easy*, its solution is again $\lambda^{\max}(A)$ at x_1 .

Here however, we want a little bit more. . . We look for a *sparse* solution and solve instead:

$$\begin{array}{ll} \max & x^T A x \\ \text{subject to} & \|x\|_2 = 1 \\ & \mathbf{Card}(x) \le k, \end{array} \tag{2}$$

where Card(x) denotes the cardinality (number of non-zero elements) of x. This is non-convex and *numerically hard*.

Related literature

Previous work:

- Cadima & Jolliffe (1995): the loadings with small absolute value are thresholded to zero.
- A non-convex method called SCoTLASS by Jolliffe & Uddin (2003). (Same problem formulation)
- Zou, Hastie & Tibshirani (2004): a regression based technique called sparse PCA (S-PCA) (SPCA). Based on the fact that PCA can be written as a regression-type (non convex) optimization problem, using LASSO Tibshirani (1996) a l₁ norm penalty.

Performance:

- These methods are either very suboptimal or *nonconvex*
- Regression: works for *large scale* examples

Start from:

$$\begin{array}{ll} \max & x^T A x\\ \text{subject to} & \|x\|_2 = 1\\ & \mathbf{Card}(x) \leq k, \end{array}$$

let $X = xx^T$, and write everything in terms of the matrix X:

$$\begin{array}{ll} \max & \mathbf{Tr}(AX) \\ \text{subject to} & \mathbf{Tr}(X) = 1 \\ & \mathbf{Card}(X) \leq k^2 \\ & X = xx^T. \end{array}$$

This is *strictly equivalent!*

Why? If $X = xx^T$, then:

- in the objective: $x^T A x = \mathbf{Tr}(AX)$
- the constraint $\mathbf{Card}(x) \leq k$ becomes $\mathbf{Card}(X) \leq k^2$
- the constraint $||x||_2 = 1$ becomes $\mathbf{Tr}(X) = 1$.

We can go a little further and replace $X = xx^T$ by an equivalent $X \succeq 0$, $\operatorname{\mathbf{Rank}}(X) = 1$, to get:

max
$$\mathbf{Tr}(AX)$$

subject to $\mathbf{Tr}(X) = 1$
 $\mathbf{Card}(X) \le k^2$
 $X \succeq 0, \ \mathbf{Rank}(X) = 1,$
(3)

Again, this is the same problem!

Numerically, this is still *hard*:

- The $\mathbf{Card}(X) \leq k^2$ is still non-convex
- So is the constraint $\operatorname{\mathbf{Rank}}(X) = 1$

but, we have made *some progress*:

- The objective $\mathbf{Tr}(AX)$ is now *linear* in X
- The (non-convex) constraint $||x||_2 = 1$ became a *linear* constraint $\mathbf{Tr}(X) = 1$.

To solve this problem *efficiently*, we need to relax the two non-convex constraints above.

Easy to do here. . .

If $u \in \mathbf{R}^p$, $\mathbf{Card}(u) = q$ implies $||u||_1 \le \sqrt{q} ||u||_2$. We transform the non-convex problem into a convex relaxation:

- Replace $\mathbf{Card}(X) \leq k^2$ by the weaker (but convex) $\mathbf{1}^T | X | \mathbf{1} \leq k$
- Simply drop the rank constraint

Our problem becomes now:

$$\begin{array}{ll} \max & \mathbf{Tr}(AX) \\ \text{subject to} & \mathbf{Tr}(X) = 1 \\ & \mathbf{1}^T |X| \mathbf{1} \le k \\ & X \succeq 0, \end{array} \tag{4}$$

This is a convex program and can be solved *efficiently*.

Semidefinite programming

In fact, we get a **semidefinite program** in the variable $X \in \mathbf{S}^n$, which can be solved using *SEDUMI* by Sturm (1999) or *SDPT3* by Toh, Todd & Tutuncu (1996).

max	$\mathbf{Tr}(AX)$
subject to	$\operatorname{Tr}(X) = 1$
	$1^T X 1 \le k$
	$X \succeq 0.$

- Polynomial complexity. . .
- Problem here: the program has $O(n^2)$ dense constraints on the matrix X (sampling fails, . . .).

In practice, hard to solve problems with large n without additional work.

Singular Value Decomposition

Same technique works for Singular Value Decomposition instead of PCA.

• The variational formulation of *SVD* is here:

min
$$\|A - uv^T\|_F$$

subject to $\mathbf{Card}(u) \leq k_1$
 $\mathbf{Card}(v) \leq k_2,$

in the variables $(u, v) \in \mathbf{R}^m \times \mathbf{R}^n$ where $k_1 \leq m$, $k_2 \leq n$ are fixed.

• This can be relaxed as the following *semidefinite program*:

$$\begin{array}{ll} \max & \mathbf{Tr}(A^T X_{12}) \\ \text{subject to} & X \succeq 0, \ \mathbf{Tr}(X_{ii}) = 1 \\ & \mathbf{1}^T |X_{ii}| \mathbf{1} \le k_i, \quad i = 1, 2 \\ & \mathbf{1}^T |X_{12}| \mathbf{1} \le \sqrt{k_1 k_2}, \end{array}$$

in the variable $X \in \mathbf{S}^{m+n}$ with blocks X_{ij} for i, j = 1, 2.

Large-scale problems

IP versus first-order methods

Interior Point methods for semidefinite/cone programs

- Produce a solution up to *machine precision*
- Compute a Newton step at each iteration: *costly*

In our case:

- We are not really interested in getting a solution up to machine precision
- The problems are *too big* to compute a Newton step. . .

Solution: use *first-order techniques*...

First-order methods

Basic model for the problem: *black-box* oracle producing

- the function value f(x)
- a subgradient $g(x) \in \partial f(x)$

f is here convex, non-smooth. Using only this info, we need $O(1/\varepsilon^2)$ steps to find an ε -optimal solution.

However, if the function is convex with a *Lipschitz-continuous gradient* with constant L then

• we need only
$$O\left(\sqrt{L/arepsilon}
ight)$$
 steps to get an $arepsilon$ -optimal solution...

Smoothness brings a *massive* improvement in complexity. . .

Sparse PCA?

In our case, we look at a penalized version of the relaxed sparse PCA problem:

$$\max_{U} \operatorname{Tr}(AU) - \mathbf{1}^{T} |U| \mathbf{1} : U \succeq 0, \operatorname{Tr} U = 1.$$
(5)

Difference?

- If we can solve the dual, these two formulations are equivalent.
- Otherwise: scale A...

Problem here, the function to minimize is not smooth! Can we hope to do better than the worst case complexity of $O(1/\varepsilon^2)$?

Nesterov (2003): the answer is *yes*, exploits particular *problem structure*...

Sparse PCA?

We can rewrite our problem as a *convex-concave* game:

$$\max_{\{U \succeq 0, \operatorname{Tr} U=1\}} \operatorname{Tr}(AU) - \mathbf{1}^T |U| \mathbf{1} = \min_{X \in \mathcal{Q}_1} \max_{U \in \mathcal{Q}_2} \langle X, U \rangle + \operatorname{Tr}(AU)$$

where

•
$$Q_1 = \{ X \in S^n : |X_{ij}| \le 1, \ 1 \le i, j \le n \}$$

•
$$\mathcal{Q}_2 = \{ U \in \mathcal{S}^n : \operatorname{Tr} U = 1 \}$$

Sparse PCA: complexity

Why a *convex-concave* game?

- Recent result by Nesterov (2003) shows that this specific structure can be exploited to significantly reduce the complexity compared to the black-box case
- All the algorithm steps can be worked out explicitly in this case

Result:

• Complexity down to $O(1/\varepsilon)$ instead of $O(1/\varepsilon^2)!$

Smooth minimization of non-smooth functions

What makes the algorithm in Nesterov (2003) work:

- First use the convex-concave game structure to regularize the function. (Inf-convolution with strictly convex function, à la Moreau-Yosida. See for example Lemaréchal & Sagastizábal (1997))
- Then use the optimal first-order minimization algorithm in Nesterov (1983) to minimize the smooth approximation.

The method works particularly well if:

- All the steps in the regularization can be performed in closed-form
- All the auxiliary minimization sub-problems can be solved in closed-form

This is the case here.

Complexity

• Max number of iterations is given by

$$N = 4 \|B\|_{1,2} \sqrt{\frac{D_1 D_2}{\sigma_1 \sigma_2}} \cdot \frac{1}{\epsilon},$$

with

$$D_1 = n^2/2, \ \sigma_1 = 1, \ D_2 = \log(n), \ \sigma_2 = 1, \ \|B\|_{1,2} = 1.$$

• Since each iteration costs ${\cal O}(n^3)$ flops, the worst-case flop count to get a $\varepsilon\text{-optimal solution}$ is given by

$$O\left(\frac{n^4\sqrt{\log n}}{\epsilon}\right)$$

Robustness & sparse decomposition

Duality - robustness

We look at the penalized problem:

max.
$$\operatorname{Tr}(AU) - \rho \mathbf{1}^T |U| \mathbf{1}$$

s.t. $\operatorname{Tr} U = 1$
 $U \succeq 0$

which can be written:

$$\max_{\{\operatorname{Tr} U=1, U \succeq 0\}} \min_{\{|X_{ij}| \le \rho\}} \operatorname{Tr}((A+X)U)$$

or also:

$$\min_{\{|X_{ij}| \le \rho\}} \quad \lambda^{\max}(A+X)$$

This dual has a very natural interpretation. . .

Duality - robustness

 $\min_{\{|X_{ij}| \le \rho\}} \quad \lambda^{\max}(A+X)$

- Worst-case *robust* maximum eigenvalue problem
- Uniformly distributed noise with magnitude ρ on the coefficients of the covariance matrix A

We ask for *sparsity*, we get *robustness* at the same time. . .

Sparse PCA: stopping the decomposition

Standard PCA:

- Finite decomposition, will stop after at most n eigenvectors are found
- Orthogonal decomposition

However, use the *robustness* interpretation:

- Run the decomposition
- Test if $\max_{ij} |A_{ij}| \le \rho$.
- If yes the matrix is *undistinguishable* from the *noise*, stop. . .

Numerical results

Numerical results

Compare with existing techniques. . .

- PITPROPS data from Zou et al. (2004)
- Compare regression technique and semidefinite relaxation (DSPCA) detailed here
- Test a sparse PCA on the PITPROPS data:
- Match the explained variance for each factor
- Minimize factor cardinality using regression & DSPCA



Cumulative cardinality and cumulative explained variance for SPCA and DSPCA as a function of the number of principal components: black line for normal PCA, blue for SPCA and red for DSPCA (full for $k_1 = 5$ and dash-dot for $k_1 = 6$).

Sparse factors. . .

Example:

- Use a covariance matrix from forward rates with maturity 1Y to 10Y
- Compute first factor normally (average of rates)
- Apply the DSPCA technique to get a sparse second factor



Second Factor



The second factor is much sparser than in the PCA case, explained variance goes from 16% to 14%. . .

Cardinality versus *k***: model**

Start with a sparse vector v = (1, 0, 1, 0, 1, 0, 1, 0, 1, 0). We then define the matrix A as:

$$A = U^T U + 15 \ v v^T$$

here $U \in \mathbf{S}^{10}$ is a random matrix (uniform coefs in [0,1]).

We solve:

$$\begin{array}{ll} \max & \mathbf{Tr}(AX) \\ \text{subject to} & \mathbf{Tr}(X) = 1 \\ & \mathbf{1}^T |X| \mathbf{1} \leq k \\ & X \succeq 0, \end{array}$$

• Try
$$k = 1, ..., 10$$

- For each k, sample a 100 matrices A
- Plot average solution cardinality (and standard dev. as error bars)

Cardinality versus *k*



(k+1) is a good predictor of the cardinality. . .

Sparsity versus # iterations

Start with a sparse vector $v = (1, 0, 1, 0, 1, 0, 1, 0, 1, 0, \dots, 0) \in \mathbb{R}^{20}$. We then define the matrix A as:

$$A = U^T U + 100 \ v v^T$$

here $U \in \mathbf{S}^{20}$ is a random matrix (uniform coefs in [0, 1]).

We solve:

max
$$\operatorname{Tr}(AU) - \rho \mathbf{1}^T |U| \mathbf{1}$$

s.t. $\operatorname{Tr} U = 1$
 $U \succeq 0$

for $\rho = 5$.

Sparsity versus # iterations



Number of iterations: 10,000 to 100,000. Computing time: 12" to 110".

Conclusion

- Semidefinite relaxation for sparse PCA
- *Robustness* & *sparsity* at the same time (cf. dual)
- Can solve large-scale problems with *optimal* first-order method by Nesterov (2003)

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